

# MODEL PREDICTIVE CONTROL

## HYBRID MODELS FOR MPC

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[http://cse.lab.imtlucca.it/~bemporad/mpc\\_course.html](http://cse.lab.imtlucca.it/~bemporad/mpc_course.html)

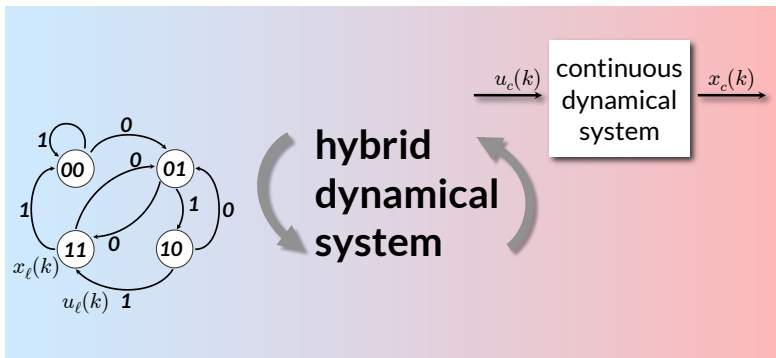


# COURSE STRUCTURE

- ✓ Basic concepts of model predictive control (MPC) and linear MPC
- ✓ Linear time-varying and nonlinear MPC
- ✓ Quadratic programming (QP) and explicit MPC
  - Hybrid MPC
  - Stochastic MPC
  - Learning-based MPC

# HYBRID MODELS

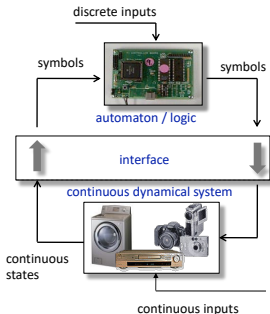
# HYBRID DYNAMICAL SYSTEMS



- Variables are **binary-valued**  
 $x_\ell \in \{0, 1\}^{n_\ell}$ ,  $u_\ell \in \{0, 1\}^{m_\ell}$
- Dynamics = **finite state machine**
- **Logic constraints**

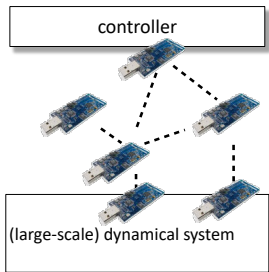
- Variables are **real-valued**  
 $x_c \in \mathbb{R}^{n_c}$ ,  $u_c \in \mathbb{R}^{m_c}$
- **Difference/differential equations**
- **Linear inequality** constraints

# TECHNOLOGICAL PUSH FOR STUDYING HYBRID SYSTEMS



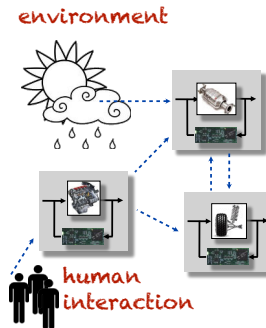
**embedded  
systems**

>1995



**networked  
control systems**

>2005



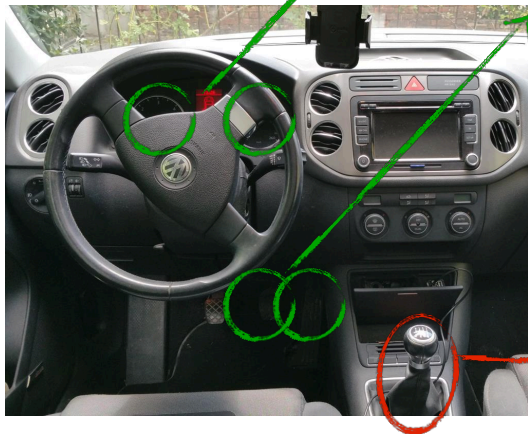
**cyber-physical  
systems**

>2010

# AN EXAMPLE OF "INTRINSICALLY HYBRID" SYSTEM

- Vehicle

continuous dynamical variables  
(speed, torque, ...)



continuous commands  
(brake & gas pedal)

+

discrete command  
(R,N,1,2,3,4,5)

# KEY REQUIREMENTS FOR HYBRID MODELS

- **Descriptive** enough to capture the behavior of the system
  - **continuous** dynamics (physical systems)
  - **logic** components (switches, automata)
  - **interconnection** between logic and dynamics
- **Simple** enough for solving analysis and synthesis problems

$$\left\{ \begin{array}{l} x' = Ax + Bu \\ y = Cx + Du \end{array} \right. \longleftrightarrow \left\{ \begin{array}{l} x' = f(x, u, t) \\ y = g(x, u, t) \end{array} \right.$$

**linear hybrid systems**

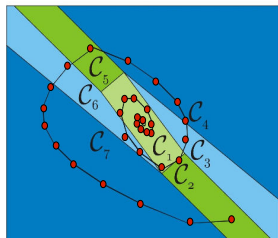
“Perfection is achieved not when there is nothing more to add,  
but when there is nothing left to take away.”



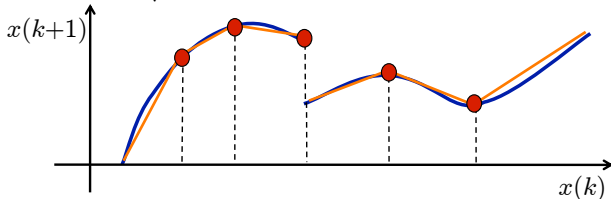
A. de Saint-Exupéry  
(1900–1944)

# PIECEWISE AFFINE SYSTEMS

$$\begin{aligned}x(k+1) &= A_{i(k)}x(k) + B_{i(k)}u(k) + f_{i(k)} \\y(k) &= C_{i(k)}x(k) + D_{i(k)}u(k) + g_{i(k)} \\i(k) &\text{ s.t. } H_{i(k)}x(k) + J_{i(k)}u(k) \leq K_{i(k)}\end{aligned}$$



- PWA systems can approximate nonlinear dynamics arbitrarily well (even discontinuous ones)

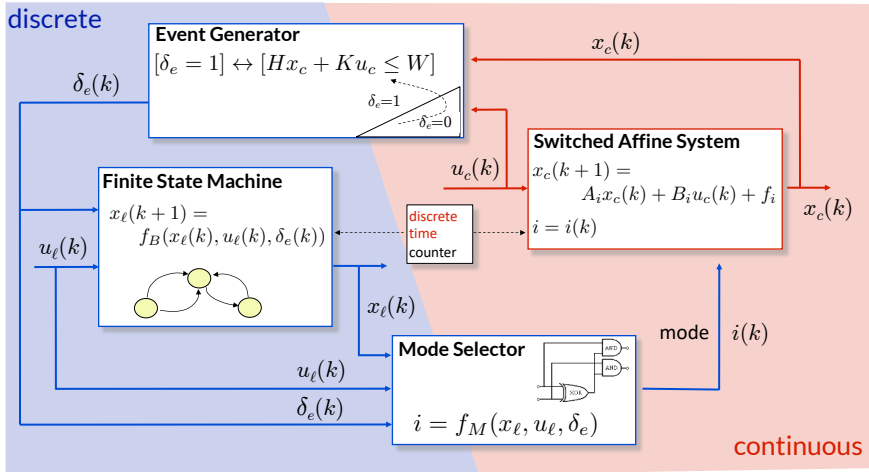


- Model **descriptiveness** vs model **complexity tradeoff**

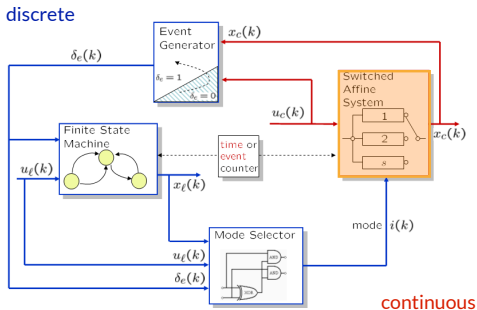


# DISCRETE HYBRID AUTOMATON (DHA)

(Torrissi, Bemporad, 2004)



# SWITCHED AFFINE SYSTEM

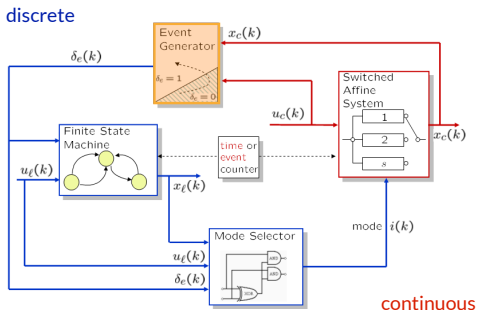


- The **affine dynamics** depend on the current mode  $i(k)$ :

$$x_c(k+1) = A_{i(k)}x_c(k) + B_{i(k)}u_c(k) + f_{i(k)}$$

$$x_c \in \mathbb{R}^{n_c}, u_c \in \mathbb{R}^{m_c}$$

# EVENT GENERATOR



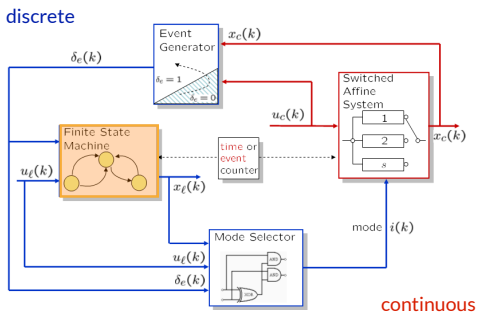
- **Event variables** are generated by **linear threshold conditions** over continuous states, continuous inputs, and time:

$$[\delta_e^i(k) = 1] \leftrightarrow [H^i x_c(k) + K^i u_c(k) \leq W^i]$$

$$x_c \in \mathbb{R}^{n_c}, \quad u_c \in \mathbb{R}^{m_c}$$
$$\delta_e \in \{0, 1\}^{n_e}$$

- Example:  $[\delta_e(k) = 1] \leftrightarrow [x_c(k) \geq 0]$

# FINITE STATE MACHINE



- The binary state of the **finite state machine** evolves according to a Boolean state update function  $f_B : \{0, 1\}^{n_\ell + m_\ell + n_e} \rightarrow \{0, 1\}^{n_\ell}$ :

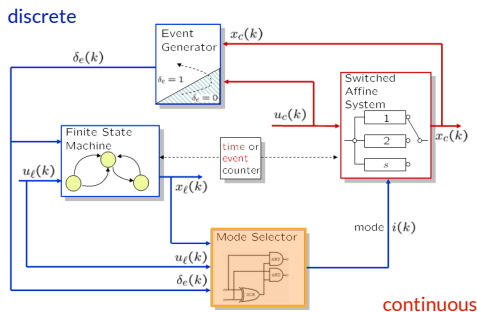
$$x_\ell(k+1) = f_B(x_\ell(k), u_\ell(k), \delta_e(k))$$

$$x_\ell \in \{0, 1\}^{n_\ell}, \quad u_\ell \in \{0, 1\}^{m_\ell}$$

$$\delta_e \in \{0, 1\}^{n_e}$$

- Example:**  $x_\ell(k+1) = \neg\delta_e(k) \vee (x_\ell(k) \wedge u_\ell(k))$

# MODE SELECTOR



The mode selector can be seen as the output function of the discrete dynamics

- The active **mode**  $i(k)$  is selected by a Boolean function of the current binary states, binary inputs, and event variables:

$$i(k) = f_M(x_\ell(k), u_\ell(k), \delta_e(k))$$

$$x_\ell \in \{0, 1\}^{n_\ell}, \quad u_\ell \in \{0, 1\}^{m_\ell}$$

$$\delta_e \in \{0, 1\}^{n_e}$$

- Example:

$$i(k) = \begin{bmatrix} \neg u_\ell(k) \vee x_\ell(k) \\ u_\ell(k) \wedge x_\ell(k) \end{bmatrix} \Rightarrow \begin{array}{c|cc} u_\ell/x_\ell & 0 & 1 \\ \hline 0 & i = \begin{bmatrix} 1 \\ 0 \end{bmatrix} & i = \begin{bmatrix} 1 \\ 0 \end{bmatrix} \\ \hline 1 & i = \begin{bmatrix} 0 \\ 0 \end{bmatrix} & i = \begin{bmatrix} 1 \\ 1 \end{bmatrix} \end{array}$$

the system has 3 modes

# CONVERSION OF LOGIC FORMULAS TO LINEAR INEQUALITIES

(Glover, 1975) (Williams, 1977) (Hooker, 2000)

- Key observation:  $X_1 \vee X_2 = \text{true}$   $\longrightarrow$   $\delta_1 + \delta_2 \geq 1, \delta_1, \delta_2 \in \{0, 1\}$
- We want to impose the Boolean statement

$$F(X_1, \dots, X_n) = \text{true}$$

- Convert the formula to **Conjunctive Normal Form (CNF)**

$$\bigwedge_{j=1}^m \left( \bigvee_{i \in P_j} X_i \bigvee_{i \in N_j} \bar{X}_i \right) = \text{true}, \quad P_j \cup N_j \subseteq \{1, \dots, n\}$$

- Transform the CNF into the equivalent linear inequalities

$$\left\{ \begin{array}{l} \sum_{i \in P_1} \delta_i + \sum_{i \in N_1} (1 - \delta_i) \geq 1 \\ \qquad \qquad \qquad \qquad \qquad \qquad \vdots \\ \sum_{i \in P_m} \delta_i + \sum_{i \in N_m} (1 - \delta_i) \geq 1 \end{array} \right. \longrightarrow \begin{array}{l} A\delta \leq b, \delta \in \{0, 1\}^n \\ \text{polyhedron} \end{array}$$

Any logic proposition can be translated into integer linear inequalities

# LOGIC → INEQUALITIES: SYMBOLIC APPROACH

- Example: we want to impose the following condition on  $X_1, X_2, X_3$

$$F(X_1, X_2, X_3) = [X_3 \leftrightarrow X_1 \wedge X_2] = \mathbf{true}$$

- Convert  $F$  to Conjunctive Normal Form (CNF):

(see e.g. <https://www.wolframalpha.com> or just google "CNF + converter" ...)

$$(X_3 \vee \neg X_1 \vee \neg X_2) \wedge (X_1 \vee \neg X_3) \wedge (X_2 \vee \neg X_3)$$

- Transform into inequalities:

$$\begin{cases} \delta_3 + (1 - \delta_1) + (1 - \delta_2) & \geq 1 \\ \delta_1 + (1 - \delta_3) & \geq 1 \\ \delta_2 + (1 - \delta_3) & \geq 1 \end{cases}$$

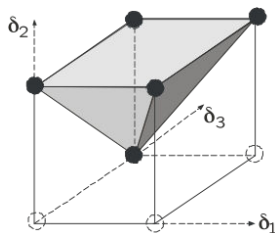

# LOGIC $\rightarrow$ INEQUALITIES: GEOMETRIC APPROACH

- Consider the Boolean statement  $F(X_1, \dots, X_n) = \text{true}$  and collect the rows of the **truth table**  $T(F)$  of  $F$

The **convex hull**  $P = \{\delta \in \mathbb{R}^n : A\delta \leq b\}$  of the **rows in**  $T(F)$  is the smallest polytope equivalent to the Boolean statement  $F$

(Mignone, Bemporad, Morari, 1999)

	$X_1$	$X_2$	...	$X_n$
$T(F)$ :	0	0	...	1
	0	1	...	1
	$\vdots$	$\vdots$	$\vdots$	$\vdots$
	1	1	...	0

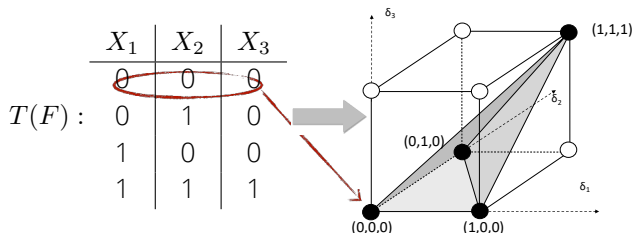


- Convex hull packages: `cdd`, `lrs`, `qhull`, `chD`, `Hull`, `Porto`  
CDDMEX package by K. Fukuda included in the Hybrid Toolbox



# LOGIC $\rightarrow$ INEQUALITIES: GEOMETRIC APPROACH

- Example:  $F(X_1, X_2, X_3) = [X_3 \leftrightarrow X_1 \wedge X_2]$  (logic **and**)



- Key idea:** white points cannot be inside the convex hull of black points

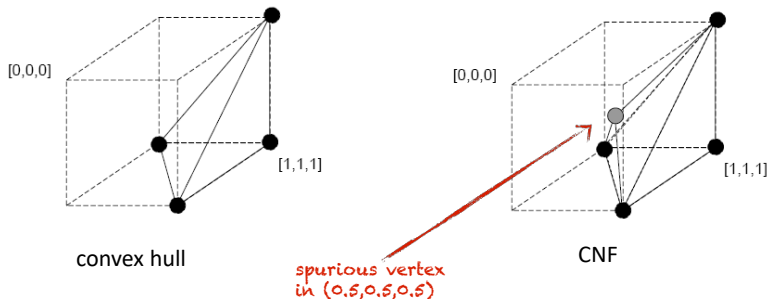
$$\text{conv} \left( \begin{bmatrix} 0 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 0 \\ 1 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 0 \\ 0 \end{bmatrix}, \begin{bmatrix} 1 \\ 1 \\ 1 \end{bmatrix} \right) = \left\{ \delta \in \mathbb{R}^3 : \begin{array}{l} -\delta_1 + \delta_3 \leq 0 \\ -\delta_2 + \delta_3 \leq 0 \\ \delta_1 + \delta_2 - \delta_3 \leq 1 \end{array} \right\}$$

```
>> V=struct('V',[0 0 0;0 1 0;1 0 0;1 1 1]);  
>> H=cddmex('hull',V);A=H.A,b=H.B
```

# GEOMETRIC VS SYMBOLIC APPROACH

- The polyhedron obtained via convex hull is the smallest one
- The one obtained via CNF may be larger. Example:

$$(X_1 \vee X_2) \wedge (X_1 \vee X_3) \wedge (X_2 \vee X_3) = \text{true}$$



- **Note:** no other example with 3 vars but  
 $(X_1 \vee X_2) \wedge (X_1 \vee X_3) \wedge (X_2 \vee X_3) \wedge (\neg X_1 \vee \neg X_2 \vee \neg X_3) = \text{true}$

# BIG-M TECHNIQUE (IFF)

- Consider the **if-and-only-if** condition

$$[\delta = 1] \leftrightarrow [a'x_c - b \leq 0]$$

$$\begin{aligned} x_c &\in \mathcal{X} \\ \delta &\in \{0, 1\} \end{aligned}$$

- Assume  $\mathcal{X} \subset \mathbb{R}^{n_c}$  bounded. Let  $M$  and  $m$  such that  $\forall x_c \in \mathcal{X}$

$$M > a'x_c - b$$

$$m < a'x_c - b$$

- The if-and-only-if condition is equivalent to

$$\begin{cases} a'x_c - b \leq M(1 - \delta) \\ a'x_c - b > m\delta \end{cases}$$

- We can replace the second constraint with  $a'x_c - b \geq \epsilon + (m - \epsilon)\delta$  to avoid strict inequalities, where  $\epsilon > 0$  is a small number (e.g., the machine precision)

# COMPUTING THE BIG-M

- If  $\mathcal{X}$  is a polyhedron, we can use linear programming

$$m < \min_{x_c \in \mathcal{X}} \{a'x_c\} - b$$

$$M > -\min_{x_c \in \mathcal{X}} \{-a'x_c\} - b$$

- If  $\mathcal{X}$  is a box,  $\mathcal{X} = [\ell, u]$ , we can use the following simpler method

(Lee, Kouvaritakis, 2000) (Bemporad, 2022)

$$a_+ = \max\{a, 0\}$$

$$a_- = a_+ - a = \max\{-a, 0\}$$

$$m < a'_+ \ell - a'_- u - b$$

$$M > a'_+ u - a'_- \ell - b$$

# BIG-M TECHNIQUE (IF-THEN-ELSE)

- Consider the **if-then-else** condition

$$z = \begin{cases} a'_1 x_c - b_1 & \text{if } \delta = 1 \\ a'_2 x_c - b_2 & \text{otherwise} \end{cases}$$

$$\begin{aligned} x_c &\in \mathcal{X} \\ \delta &\in \{0, 1\} \\ z &\in \mathbb{R} \end{aligned}$$

- Assume  $\mathcal{X} \subset \mathbb{R}^{n_c}$  bounded. Let  $M_1, M_2$  and  $m_1, m_2$  such that  $\forall x_c \in \mathcal{X}$

$$\begin{aligned} M_1 &> a'_1 x_c - b_1 > m_1 \\ M_2 &> a'_2 x_c - b_2 > m_2 \end{aligned}$$

- The if-then-else condition is equivalent to

$$\begin{cases} z \leq a'_1 x_c - b_1 - (m_1 - M_2)(1 - \delta) \\ z \geq a'_1 x_c + b_1 + (m_2 - M_1)(1 - \delta) \\ z \leq a'_2 x_c - b_2 - (m_2 - M_1)\delta \\ z \geq a'_2 x_c - b_2 + (m_1 - M_2)\delta \end{cases}$$

# SWITCHED AFFINE SYSTEM

- The state-update equation of a SAS can be rewritten as

$$x_c(k+1) = \sum_{i=1}^s z_i(k) \quad z_i(k) \in \mathbb{R}^{n_c}$$

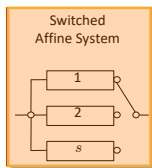
with

$$z_1(k) = \begin{cases} A_1 x_c(k) + B_1 u_c(k) + f_1 & \text{if } \delta_1(k) = 1 \\ 0 & \text{otherwise} \end{cases}$$
$$\vdots$$
$$z_s(k) = \begin{cases} A_s x_c(k) + B_s u_c(k) + f_s & \text{if } \delta_s(k) = 1 \\ 0 & \text{otherwise} \end{cases}$$

and with  $\delta_i(k) \in \{0, 1\}$  subject to the **exclusive or** condition

$$\sum_{i=1}^s \delta_i(k) = 1 \text{ or equivalently } \begin{cases} \sum_{i=1}^s \delta_i(k) \geq 1 \\ \sum_{i=1}^s \delta_i(k) \leq 1 \end{cases}$$

- Output eq.  $y_c(k) = C_i x_c(k) + D_i u_c(k) + g_i$  admits a similar transformation



# TRANSFORMATION OF A DHA INTO LINEAR (IN)EQUALITIES

$$X_1 \vee X_2 = \text{TRUE} \quad \longrightarrow \quad \delta_1 + \delta_2 \geq 1, \quad \delta_1, \delta_2 \in \{0, 1\}$$

Any logic statement

$$f(X) = \text{TRUE}$$

$$\bigwedge_{j=1}^m \left( \bigvee_{i \in P_j} X_i \vee \bigvee_{i \in N_j} \neg X_i \right) \quad (\text{CNF})$$

$N_j, P_j \subseteq \{1, \dots, n\}$

$$\begin{cases} 1 \leq \sum_{i \in P_1} \delta_i + \sum_{i \in N_1} (1 - \delta_i) \\ \vdots \\ 1 \leq \sum_{i \in P_m} \delta_i + \sum_{i \in N_m} (1 - \delta_i) \end{cases}$$

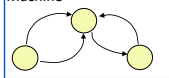
$$[\delta_e^i(k) = 1] \leftrightarrow [H^i x_c(k) \leq W^i]$$

$$\begin{cases} H^i x_c(k) - W^i \leq M^i (1 - \delta_e^i(k)) \\ H^i x_c(k) - W^i > m^i \delta_e^i(k) \end{cases}$$

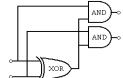
$$\begin{aligned} \text{IF } [\delta = 1] \text{ THEN } z &= a_1^T x + b_1^T u + f_1 \\ \text{ELSE } z &= a_2^T x + b_2^T u + f_2 \end{aligned}$$

$$\begin{cases} (m_1 - M_2)(1 - \delta) + z \leq a_1 x + b_1 u + f_1 \\ (m_2 - M_1)(1 - \delta) - z \leq -a_1 x - b_1 u - f_1 \\ (m_2 - M_1)\delta + z \leq a_2 x + b_2 u + f_2 \\ (m_1 - M_2)\delta - z \leq -a_2 x - b_2 u - f_2 \end{cases}$$

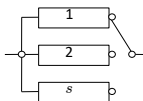
Finite State Machine



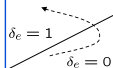
Mode Selector



Switched Affine System



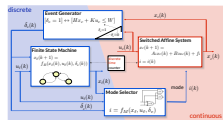
Event Generator



# MIXED LOGICAL DYNAMICAL (MLD) SYSTEMS

(Bemporad, Morari, 1999)

- By converting logic relations into mixed-integer linear inequalities a DHA can be rewritten as the **Mixed Logical Dynamical (MLD)** system



$$\begin{cases} x(k+1) &= Ax(k) + B_1u(k) + B_2\delta(k) + B_3z(k) + B_5 \\ y(k) &= Cx(k) + D_1u(k) + D_2\delta(k) + D_3z(k) + D_5 \\ E_2\delta(k) + E_3z(k) &\leq E_4x(k) + E_1u(k) + E_5 \end{cases}$$



$$\begin{aligned} x &\in \mathbb{R}^{n_c} \times \{0, 1\}^{n_b}, \quad u \in \mathbb{R}^{m_c} \times \{0, 1\}^{m_b} \\ y &\in \mathbb{R}^{p_c} \times \{0, 1\}^{p_b}, \quad \delta \in \{0, 1\}^{r_b}, \quad z \in \mathbb{R}^{r_c} \end{aligned}$$

- The translation from DHA to MLD can be automatized, see e.g. the language **HYSDEL** (HYbrid Systems DEscription Language) (Torrise, Bemporad, 2004)
- MLD models allow solving MPC, verification, state estimation, and fault detection problems via **mixed-integer programming**



# A SIMPLE EXAMPLE OF MLD SYSTEM

- PWA system<sup>1</sup>: 
$$x(k+1) = \begin{cases} 0.8x(k) + u(k) & \text{if } x(k) \geq 0 \\ -0.8x(k) + u(k) & \text{if } x(k) < 0 \end{cases}$$

- Introduce event variable  $[\delta(k) = 1] \leftrightarrow [x(k) \geq 0]$  and use big-M technique:

$$\begin{aligned} \longrightarrow \quad & x(k) \geq m(1 - \delta(k)) & M = -m = 10 \\ & x(k) \leq -\epsilon + (M + \epsilon)\delta(k) & \epsilon > 0 \text{ "small"} \end{aligned}$$

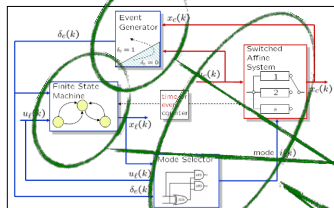
- Since  $x(k+1) = 1.6\delta(k)x(k) - 0.8x(k) + u(k)$ , introduce the aux variable

$$\begin{aligned} z(k) &\leq M\delta(k) \\ z(k) &\geq m\delta(k) \\ z(k) &\leq x(k) - m(1 - \delta(k)) \\ z(k) &\geq x(k) - M(1 - \delta(k)) \end{aligned} \quad \delta(k) \in \{0, 1\}$$

- Linear state update:  $x(k+1) = -0.8x(k) + 1.6z(k) + u(k)$

<sup>1</sup> This is the nonlinear system  $x(k+1) = 0.8|x(k)| + u(k)$

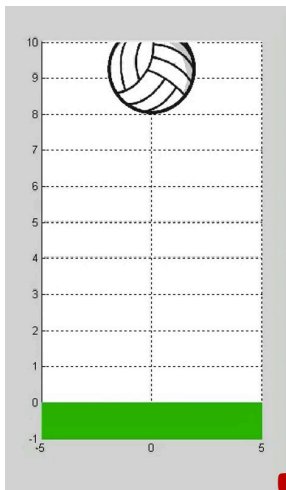
# DHA AND HYSDEL MODELS



Additional relations  
constraining system's  
variables

```
SYSTEM name {  
  INTERFACE {  
    STATE {  
      REAL xc [xmin,xmax];  
      BOOL xl; }  
    INPUT {  
      REAL uc [umin,umax];  
      BOOL ul; }  
    PARAMETER {  
      REAL param1 = 1;}  
  } /* end of interface */  
  
  IMPLEMENTATION {  
    AUX { BOOL d;  
          REAL z; }  
  
    AUTOMATA { xl = xl & ~ul; }  
  
    AD { d = xc - 1 <= 0; }  
  
    DA { z = { IF d THEN 2*xc ELSE -xc }; }  
  
    CONTINUOUS {  
      xc = z; }  
  
    MUST {  
      xc + uc <= 2;  
      ~(xl & ul); }  
  } /* end implementation */  
} /* end system */
```

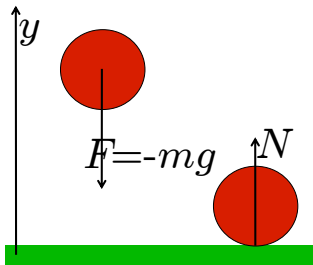
# BOUNCING BALL EXAMPLE



$$\ddot{y} = -g$$

$$y \leq 0 \Rightarrow \dot{y}(t^+) = -(1 - \alpha)\dot{y}(t^-)$$

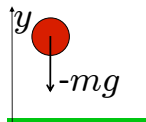
$$\alpha \in [0, 1]$$



How to model the bouncing ball as a discrete-time hybrid system ?

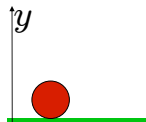
# BOUNCING BALL – TIME DISCRETIZATION

- Case  $y(k) > 0$  (ball falling):  
$$v(k) \approx \frac{y(k) - y(k-1)}{T_s}$$
$$-g \approx \frac{v(k) - v(k-1)}{T_s}$$



$$\begin{cases} v(k+1) = v(k) - T_s g \\ y(k+1) = y(k) + T_s v(k+1) \\ \quad = y(k) + T_s v(k) - T_s^2 g \end{cases}$$

- Case  $y(k) \leq 0$  (ground level):  
$$v(k+1) = -(1 - \alpha)v(k)$$
$$y(k+1) = y(k) = y(k) - T_s v(k)$$



$$\begin{cases} v(k+1) = -(1 - \alpha)v(k) \\ y(k+1) = y(k) - T_s v(k) \end{cases}$$

- We need a binary variable  $[\delta(k) = 1] \leftrightarrow [y(k) \leq 0]$

# BOUNCING BALL - HYSDEL MODEL

```
SYSTEM bouncing_ball {
INTERFACE {
/* Description of variables and constants */
STATE { REAL height [-10,10];
        REAL velocity [-100,100];  }

PARAMETER {
        REAL g;
        REAL alpha; /* 0=elastic, 1=completely anelastic */
        REAL Ts; }
}
IMPLEMENTATION {
AUX {  BOOL negative;
        REAL hnext;
        REAL vnext;  }

AD {   negative = height <= 0; }

DA {   hnext = { IF negative THEN height-Ts*velocity
                 ELSE height+Ts*velocity-Ts*Ts*g};
        vnext = { IF negative THEN -(1-alpha)*velocity
                 ELSE velocity-Ts*g};  }

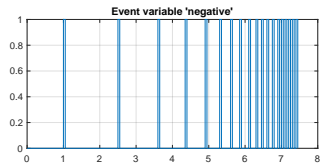
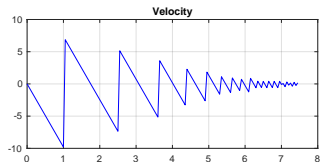
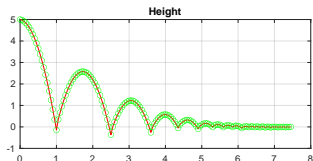
CONTINUOUS {
        height = hnext;
        velocity = vnext;}
}}
```

go to `demo demos/hybrid/bball.m`

# BOUNCING BALL - SIMULATION

```
>> Ts=0.05;  
>> g=9.8;  
>> alpha=0.3;  
  
>> S=mld('bouncing_ball',Ts);  
  
>> N=150;  
>> U=zeros(N,0);  
>> x0=[5 0]';  
  
>> [X,T,D]=sim(S,x0,U);
```

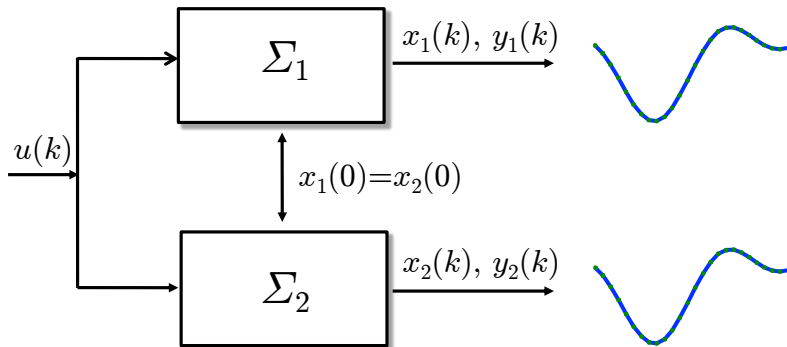
- **Note:** no **Zeno effect** in discrete time !



# **EQUIVALENCE OF HYBRID MODELS**

# EQUIVALENCE OF HYBRID MODELS

- Two hybrid models  $\Sigma_1, \Sigma_2$  are **equivalent** if for all initial states  $x_1(0) = x_2(0)$  and input excitations  $u_1(k) \equiv u_2(k)$ , the corresponding trajectories  $x_1(k) \equiv x_2(k)$  and  $y_1(k) \equiv y_2(k), \forall k = 0, 1, \dots$





# EQUIVALENCE OF HYBRID MODELS

- **MLD** and **PWA** systems are equivalent (Bemporad, Ferrari-Trecate, Morari, 2000)

Proof: For a given combination  $(x_\ell, u_\ell, \delta)$  of an MLD model, the state and output equation are linear and valid in a polyhedron.

Conversely, a PWA system can be modeled as MLD system (see next slide)

- Efficient **conversion algorithms** from MLD to PWA form exist  
(Bemporad, 2004) (Geyer, Torrisi, Morari, 2003)
- Further equivalences exist with other classes of hybrid dynamical systems, such as **Linear Complementarity (LC)** systems (Heemels, De Schutter, Bemporad, 2001)

# MODELING PWA SYSTEMS IN MLD FORM

- PWA system with bounded states  $x \in \mathcal{X}$  and inputs  $u \in \mathcal{U}$  with  $s$  regions

$$\begin{aligned}x(k+1) &= A_{i(k)}x(k) + B_{i(k)}u(k) + f_{i(k)} \\y(k) &= C_{i(k)}x(k) + D_{i(k)}u(k) + g_{i(k)} \\i(k) &= \text{such that } \begin{bmatrix} x(k) \\ u(k) \end{bmatrix} \in \mathcal{C}_{i(k)}\end{aligned} \quad i(k) \in \{1, 2, \dots, s\}$$

- The sets  $\mathcal{C}_i = \{ \begin{bmatrix} x \\ u \end{bmatrix} : H_i x + J_i u \leq K_i \}$  make a **polyhedral partition** of  $\mathcal{X} \times \mathcal{U}$  ( $\cup_{i=1}^s \mathcal{C}_i = \mathcal{X} \times \mathcal{U}$  and  $\overset{\circ}{\mathcal{C}}_i \cap \overset{\circ}{\mathcal{C}}_j = \emptyset, \forall i \neq j, i, j = 1, \dots, s$ )
- Introduce  $s$  binary variables  $\delta_i, i = 1, \dots, s$  and set the logic constraints

$$[\delta_i = 1] \rightarrow [H_i x + J_i u \leq K_i]$$

$$\bigoplus_{i=1}^s [\delta_i = 1] = \text{true}$$



$$\begin{aligned}H_i x + J_i u &\leq K_i + M_i(1 - \delta_i) \\ \sum_{i=1}^s \delta_i &= 1\end{aligned}$$

where the vector  $M_i$  of upper-bounds can be computed, e.g., via LP

# MODELING PWA SYSTEMS IN MLD FORM

- Introduce auxiliary real vectors  $z_i, w_i$  defined by if-then-else rules

$$z_i = \begin{cases} A_i x + B_i u + f_i & \text{if } \delta_i = 1 \\ 0 & \text{otherwise} \end{cases} \quad w_i = \begin{cases} C_i x + D_i u + g_i & \text{if } \delta_i = 1 \\ 0 & \text{otherwise} \end{cases}$$

and convert into mixed-integer inequalities (**big-M technique**):

$$\begin{cases} z_i \leq A_i x + B_i u + f_i - m_i(1 - \delta_i) \\ z_i \geq A_i x + B_i u + f_i - M_i(1 - \delta_i) \\ z_i \leq M_i \delta_i \\ z_i \geq m_i \delta_i \end{cases}$$

$$m \leq A_i x + B_i u + f_i \leq M_i$$

- Finally, write the state update and output equations

$$x(k+1) = \sum_{i=1}^s z_i(k) \quad y(k) = \sum_{i=1}^s w_i(k)$$

# MODELING PWA SYSTEMS VIA DISJUNCTIVE PROGRAMMING

(Balas, 1985)

- A PWA system with bounded states and inputs is equivalent to the **disjunction**

$$\bigvee_{i=1}^s \left[ \begin{array}{l} H_i x(k) + J_i u(k) \leq K_i \\ x(k+1) = A_i x(k) + B_i u(k) + f_i \end{array} \right] \quad \begin{array}{l} x_{lb} \leq x(k) \leq x_{ub} \\ u_{lb} \leq u(k) \leq u_{ub} \end{array}$$

- Introduce  $s$  binary variables  $\delta_1(k), \dots, \delta_s(k)$  subject to  $\sum_{i=1}^s \delta_i(k) = 1$
- Introduce the **convex hull relaxation** of the disjunction

$$\begin{aligned} x(k) &= \sum_{i=1}^s v_i(k), & x_{lb}\delta_i(k) &\leq v_i(k) \leq x_{ub}\delta_i(k) \\ u(k) &= \sum_{i=1}^s w_i(k), & u_{lb}\delta_i(k) &\leq w_i(k) \leq u_{ub}\delta_i(k) \end{aligned}$$

and set

$$x(k+1) = \sum_{i=1}^s A_i v_i(k) + B_i w_i(k) + f_i \delta_i(k), \quad H_i v_i(k) + J_i w_i(k) \leq K_i \delta_i(k)$$

- Only introduce  $s$  binary variables  $\delta_i, i = 1, \dots, s$  and set:

$$\underbrace{m_i^x(1 - \delta_i(k)) \leq x(k+1) - A_i x(k) - B_i u(k) - f_i \leq M_i^x(1 - \delta_i(k))}_{[\delta_i(k) = 1] \rightarrow [x(k+1) = A_i x(k) + B_i u(k) + f_i]}$$

$$\underbrace{m_i^y(1 - \delta_i(k)) \leq y(k) - C_i x(k) - D_i u(k) - g_i \leq M_i^y(1 - \delta_i(k))}_{[\delta_i(k) = 1] \rightarrow [y(k) = C_i x(k) + D_i u(k) + g_i]}$$

$$H_i x(k) + J_i u(k) \leq K_i + M_i(1 - \delta_i(k)) \quad [\delta_i(k) = 1] \rightarrow [H_i x(k) + J_i u(k) \leq K_i]$$

$$\sum_{i=1}^s \delta_i(k) = 1$$

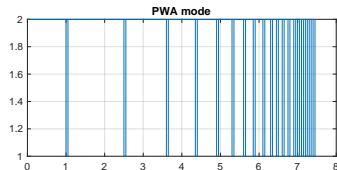
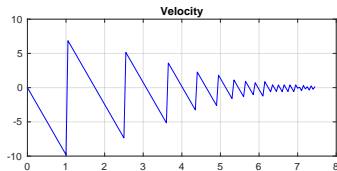
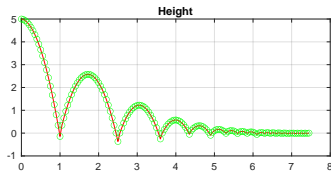
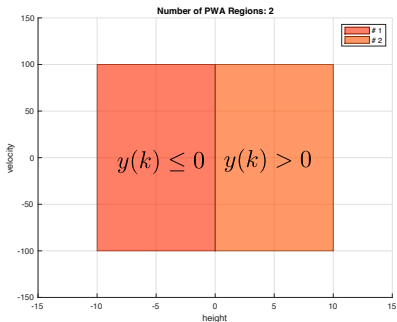
$$\bigoplus_{i=1}^s [\delta_i(k) = 1] = \text{true}$$

where  $m_i^x, M_i^x, m_i^y, M_i^y$  are suitably defined upper and lower bounds

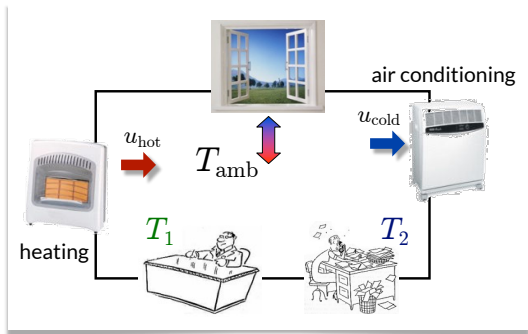
# BOUNCING BALL - PWA EQUIVALENT

```
>> P=pwa(S);  
>> plot(P)  
  
>> [X,T,I]=sim(P,x0,U);
```

(Bemporad, 2004)



# EXAMPLE: ROOM TEMPERATURE CONTROL



## discrete dynamics

- #1 = cold  $\rightarrow$  heater = on
- #2 = cold  $\rightarrow$  heater = on **unless** #1 hot
- A/C activation has similar rules

## continuous dynamics

$$\frac{dT_i}{dt} = -\alpha_i(T_i - T_{amb}) + k_i(u_{hot} - u_{cold})$$

$$i = 1, 2$$

go to demo `demos/hybrid/heatcool.m`

# EXAMPLE: ROOM TEMPERATURE CONTROL

```
SYSTEM heatcool {  
  
INTERFACE {  
    STATE { REAL T1 [-10,50];  
            REAL T2 [-10,50]; }  
    INPUT { REAL Tamb [-50,50]; }  
    PARAMETER {  
        REAL Ts, alpha1, alpha2, k1, k2;  
        REAL Thot1, Tcold1, Thot2, Tcold2, Uc, Uh; }  
    }  
  
IMPLEMENTATION {  
    AUX { REAL uhot, ucold;  
          BOOL hot1, hot2, cold1, cold2; }  
  
    AD { hot1 = T1>=Thot1;  
         hot2 = T2>=Thot2;  
         cold1 = T1<=Tcold1;  
         cold2 = T2<=Tcold2; }  
  
    DA { uhot = { IF cold1 | (cold2 & ~hot1) THEN Uh ELSE 0 };  
          ucold = { IF hot1 | (hot2 & ~cold1) THEN Uc ELSE 0 }; }  
  
    CONTINUOUS { T1 = T1+Ts*(-alpha1*(T1-Tamb)+k1*(uhot-ucold));  
                 T2 = T2+Ts*(-alpha2*(T2-Tamb)+k2*(uhot-ucold)); }  
    }  
}
```

```
>> S=mld('heatcoolmodel',Ts);
```

get the MLD model in MATLAB

```
>> [XX,TT]=sim(S,x0,U);
```

simulate the MLD model



# EXAMPLE: ROOM TEMPERATURE CONTROL

- MLD model of the room temperature system

$$\begin{cases} x(k+1) &= Ax(k) + B_1u(k) + B_2\delta(k) + B_3z(k) + B_5 \\ y(k) &= Cx(k) + D_1u(k) + D_2\delta(k) + D_3z(k) + D_5 \\ E_2\delta(k) + E_3z(k) &\leq E_4x(k) + E_1u(k) + E_5 \end{cases}$$

- 2 continuous states (temperature  $T_1, T_2$ )
- 1 continuous input (room temperature  $T_{\text{amb}}$ )
- 2 auxiliary continuous vars (power flows  $u_{\text{hot}}, u_{\text{cold}}$ )
- 6 auxiliary binary vars (4 threshold events + 2 for the OR condition)
- 20 mixed-integer inequalities

- In principle, we have  $2^6 = 64$  possible combinations of binary variables

# EXAMPLE: ROOM TEMPERATURE CONTROL

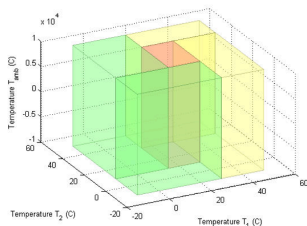
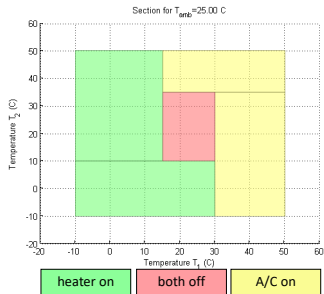
- PWA model of the room temperature system

$$x(k+1) = A_{i(k)}x(k) + B_{i(k)}u(k) + f_{i(k)}$$

$$y(k) = C_{i(k)}x(k) + D_{i(k)}u(k) + g_{i(k)}$$

>> P=pwa(S);

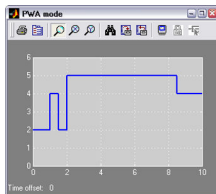
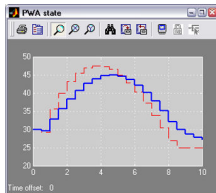
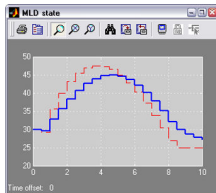
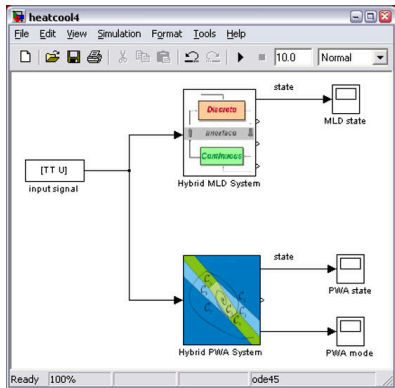
$$i(k) \text{ s.t. } H_{i(k)}x(k) + J_{i(k)}u(k) \leq K_{i(k)}$$



5 polyhedral regions  
(partition does not depend on input)

2 continuous states ( $T_1, T_2$ )  
1 continuous input ( $T_{amb}$ )

# EXAMPLE: ROOM TEMPERATURE CONTROL



- MLD and PWA models are equivalent, hence simulated states are the same

# USING PWA EQUIVALENCE FOR MODEL ANALYSIS

- Assume plant + controller can be modeled as DHA:
  - **plant** = approximated as PWA system (e.g.: nonlinear switched model)
  - **controller** = switched linear controller (e.g: combination of threshold conditions, logic, linear feedback laws, ...)
- Convert DHA to MLD form, then to PWA form
- The resulting closed-loop PWA model reveals how the closed-loop system behaves in different regions of the state-space
- Can analyze **closed-loop stability** analysis using **piecewise quadratic Lyapunov functions** (Johansson, Rantzer, 1998) (Mignone, Ferrari-Trecate, Morari, 2000)

# OTHER EXISTING HYBRID MODELS

(Heemels, De Schutter, Bemporad, 2001)

- **Linear complementarity (LC)** systems (Heemels, 1999)

$$\begin{aligned}x(k+1) &= Ax(k) + B_1u(k) + B_2w(k) \\y(k) &= Cx(k) + D_1u(k) + D_2w(k) \\v(k) &= E_1x(k) + E_2u(k) + E_3w(k) + E_4 \\0 &\leq v(k) \perp w(k) \geq 0\end{aligned}$$

**Examples:**

mechanical systems,  
electrical circuits



- **Min-max-plus-scaling (MMPS)** systems (De Schutter, Van den Boom, 2000)

$$\begin{aligned}x(k+1) &= M_x(x(k), u(k), w(k)) \\y(k) &= M_y(x(k), u(k), w(k)) \\0 &\geq M_c(x(k), u(k), w(k))\end{aligned}$$

**Example:**

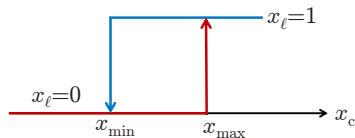
discrete-event system  
 $k$  = event counter

where  $M_{\circ}$  are MMPS functions defined by the grammar

$$M := x_i | \alpha | \max(M_1, M_2) | \min(M_1, M_2) | M_1 + M_2 | \beta M_1$$

**Example:**  $x(k+1) = 2 \max(x(k), 0) + \min(\frac{1}{2}u(k), 1)$

# MODELING HYSTERESIS



- Hysteresis between  $x_{\min} \leq x_c(k) \leq x_{\max}$
- Introduce two binary variables

$$[\delta_{\min}(k) = 1] \leftrightarrow [x_c(k) \leq x_{\min}]$$

$$[\delta_{\max}(k) = 1] \leftrightarrow [x_c(k) \geq x_{\max}]$$

- Introduce logic state  $x_\ell \in \{0, 1\}$  with dynamics

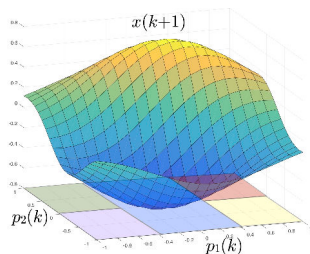
$$x_\ell(k+1) = (x_\ell(k) \wedge \neg \delta_{\min}(k)) \vee (\neg x_\ell(k) \wedge \delta_{\max}(k))$$

# IDENTIFICATION OF HYBRID SYSTEMS

# HYBRID SYSTEM IDENTIFICATION

- A hybrid model of the process may not be available from physical principles
- Therefore, a model must be either
  - **estimated from data** (model is unknown)
  - or **hybridized** (model is known but nonlinear)
- If one linear model is enough: easy problem (SYS-ID TBX) (Ljung, 1999)
- If switching sequence known: easy, just identify one linear model per mode
- If modes & dynamics must be identified simultaneously, we need **hybrid system identification** (or **piecewise affine regression**)

In industrial MPC most effort is spent in **identifying (multiple) linear prediction models** from data



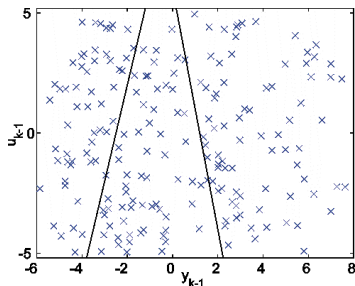


# LEARNING PWA MODELS FROM DATA

**Problem:** Estimate from data **both** the **parameters** of the affine submodels and the **partition** of the PWA map

Example: Let the data be generated by the PWARX system

$$y_k = \begin{cases} \begin{bmatrix} -0.4 & 1 & 1.5 \end{bmatrix} \phi_k + \epsilon_k \\ \text{if } \begin{bmatrix} 4 & -1 & 10 \end{bmatrix} \phi_k < 0 \\ \begin{bmatrix} 0.5 & -1 & -0.5 \end{bmatrix} \phi_k + \epsilon_k \\ \text{if } \begin{bmatrix} -4 & 1 & 10 \\ 5 & 1 & -6 \end{bmatrix} \phi_k \leq 0 \\ \begin{bmatrix} -0.3 & 0.5 & -1.7 \end{bmatrix} \phi_k + \epsilon_k \\ \text{if } \begin{bmatrix} -5 & -1 & 6 \end{bmatrix} \phi_k < 0 \end{cases}$$



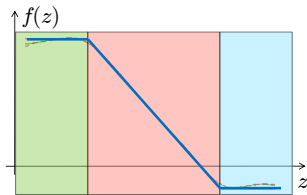
with  $\phi_k = [y_{k-1} \ u_{k-1} \ 1]'$ ,  $|u_k| \leq 5$ , and  $|\epsilon_k| \leq 0.1$

# PWA REGRESSION PROBLEM

- **Problem:** Given input/output pairs  $\{x(k), y(k)\}, k = 1, \dots, N$  and number  $s$  of models, compute a **piecewise affine** (PWA) approximation  $y \approx f(x)$

$$v(k) = \begin{cases} F_1 z(k) + g_1 & \text{if } H_1 z(k) \leq K_1 \\ \vdots \\ F_s z(k) + g_s & \text{if } H_s z(k) \leq K_s \end{cases}$$

$$v(k) = \begin{bmatrix} x(k+1) \\ y(k) \end{bmatrix}, \quad z(k) = \begin{bmatrix} x(k) \\ u(k) \end{bmatrix}$$



- Need to learn **both** the parameters  $\{F_i, g_i\}$  of the affine submodels **and** the partition  $\{H_i, K_i\}$  of the PWA map from data (**offline learning**)
- Possibly update model+partition as new data are available (**online learning**)

# APPROACHES TO PWA IDENTIFICATION

- Mixed-integer linear or quadratic programming (Roll, Bemporad, Ljung, 2004)
- Partition of infeasible set of inequalities (Bemporad, Garulli, Paoletti, Vicino, 2005)
- K-means clustering in a feature space (Ferrari-Trecate, Muselli, Liberati, Morari, 2003)
- Bayesian approach (Juloski, Wieland, Heemels, 2004)
- Kernel-based approaches (Pillonetto, 2016)
- Hyperplane clustering in data space (Münz, Krebs, 2002)
- **Recursive multiple least squares & PWL separation** (Breschi, Piga, Bemporad, 2016)
- Jump models (Bemporad, Breschi, Piga, Boyd, 2018)
- **Piecewise affine regression and classification (PARC)** (Bemporad, 2022)
- Any **machine learning** technique producing a PWA model, such as (leaky)ReLU-NNs, decision trees, softmax regression, K-nearest neighbors, ...

# PWA REGRESSION ALGORITHM

(Breschi, Piga, Bemporad, 2016)

1. Estimate models  $\{F_i, g_i\}$  **recursively**. Let  $e_i(k) = v(k) - F_i z(k) - g_i$  and only update model  $i(k)$  such that

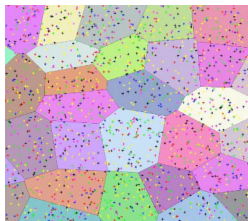
$$i(k) \leftarrow \arg \min_{i=1, \dots, s} \underbrace{e_i(k)' \Lambda_e^{-1} e_i(k)}_{\substack{\text{one-step prediction error} \\ \text{of model } \#i}} + \underbrace{(z(k) - c_i)' R_i^{-1} (z(k) - c_i)}_{\substack{\text{proximity to centroid} \\ \text{of cluster } \#i}}$$

using **recursive LS** and **inverse QR decomposition** (Alexander, Ghirnikar, 1993)

This also splits the data points  $z(k)$  in **clusters**  $C_i = \{z(k) : i(k) = i\}$

2. Compute a polyhedral partition  $\{H_i, K_i\}$  of the regressor space via **multi-category linear separation**

$$\phi(z) = \max_{i=1, \dots, s} \{w'_i z - \gamma_i\} \quad i(z) = \arg \max_i \{w'_i z - \gamma_i\}$$



3. Repeat the entire procedure  $M$  times

# PWA REGRESSION EXAMPLES

(Breschi, Piga, Bemporad, 2016)

- Identification of **piecewise-affine ARX** model

$$\begin{aligned} \begin{bmatrix} y_1(k) \\ y_2(k) \end{bmatrix} &= \begin{bmatrix} -0.83 & 0.20 \\ 0.30 & -0.52 \end{bmatrix} \begin{bmatrix} y_1(k-1) \\ y_2(k-1) \end{bmatrix} + \begin{bmatrix} -0.34 & 0.45 \\ -0.30 & 0.24 \end{bmatrix} \begin{bmatrix} u_1(k-1) \\ u_2(k-1) \end{bmatrix} \\ &+ \begin{bmatrix} 0.20 \\ 0.15 \end{bmatrix} + \max \left\{ \begin{bmatrix} 0.20 & -0.90 \\ 0.10 & -0.42 \end{bmatrix} \begin{bmatrix} y_1(k-1) \\ y_2(k-1) \end{bmatrix} \right. \\ &\left. + \begin{bmatrix} 0.42 & 0.20 \\ 0.50 & 0.64 \end{bmatrix} \begin{bmatrix} u_1(k-1) \\ u_2(k-1) \end{bmatrix} + \begin{bmatrix} 0.40 \\ 0.30 \end{bmatrix}, \begin{bmatrix} 0 \\ 0 \end{bmatrix} \right\} + e(k), \end{aligned}$$

- **Quality of fit:** best fit rate (BFR) =  $\max \left\{ 1 - \frac{\|y_i - \hat{y}_i\|_2}{\|y_i - \bar{y}_i\|_2}, 0 \right\}, i = 1, 2$

$y$  = measured,  $\hat{y}$  = open-loop simulated,  $\bar{y}$  = sample mean of  $y$

		$N = 4000$	$N = 20000$	$N = 100000$
$y_1$	(offline) RLP	96.0 %	96.5 %	99.0 %
	(Offline) RPSN	96.2 %	96.4 %	98.9 %
	(Online) ASGD	86.7 %	95.0 %	96.7 %
$y_2$	(offline) RLP	96.2 %	96.9 %	99.0 %
	(offline) RPSN	96.3 %	96.8 %	99.0 %
	(online) ASGD	87.4 %	95.2 %	96.4 %

BFR on validation data, open-loop validation

RLP = Robust linear programming

(Bennett, Mangasarian, 1994)

RPSN = Piecewise-smooth Newton method

(Bemporad, Bernardini, Patrinos, 2015)

ASGD = Averaged stochastic gradient descent

(Bottou, 2012)

- **CPU time for computing the partition:**

(i7 2.40-GHz Intel core)

	$N = 4000$	$N = 20000$	$N = 100000$
(Offline) RLP	0.308 s	3.227 s	112.435 s
(Offline) RPSN	0.016 s	0.086 s	0.365 s
(Online) ASGD	0.013 s	0.023 s	0.067 s

# PWA REGRESSION EXAMPLES

(Breschi, Piga, Bemporad, 2016)

- Identification of **linear parameter varying ARX** model

$$\begin{bmatrix} y_1(k) \\ y_2(k) \end{bmatrix} = \begin{bmatrix} \bar{a}_{1,1}(p(k)) & \bar{a}_{1,2}(p(k)) \\ \bar{a}_{2,1}(p(k)) & \bar{a}_{2,2}(p(k)) \end{bmatrix} \begin{bmatrix} y_1(k-1) \\ y_2(k-1) \end{bmatrix} + \begin{bmatrix} \bar{b}_{1,1}(p(k)) & \bar{b}_{1,2}(p(k)) \\ \bar{b}_{2,1}(p(k)) & \bar{b}_{2,2}(p(k)) \end{bmatrix} \begin{bmatrix} u_1(k-1) \\ u_2(k-1) \end{bmatrix} + e_o(k)$$

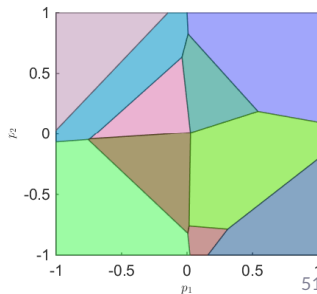
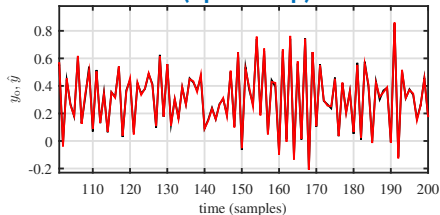
$\bar{a}(p)$  = PWA function of  $p$   
 $\bar{b}(p)$  has **quadratic** and **sin** terms

- **Quality of fit (BFR):**

	$y_1$	$y_2$
PWA regression	87 %	84 %
parametric LPV*	80 %	70 %

\* (Bamieh, Giarré, 2002)

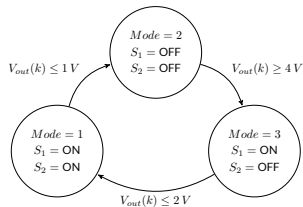
- **Validation data (open-loop):**



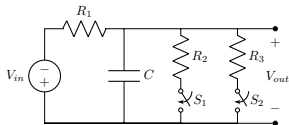
# IDENTIFICATION OF HYBRID SYSTEMS WITH LOGIC STATES

(Breschi, Piga, Bemporad, 2016)

- Identification of a **hybrid model** with logic states



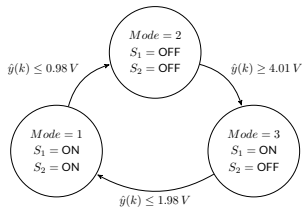
true system



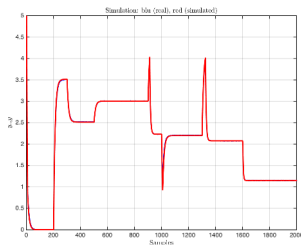
**Quality of fit:** BFR=96.64 % (validation)

**CPU time** for identification: 78 ms

(2000 samples, MacBook Pro 2.8 GHz)



identified system



# PARC – PIECEWISE AFFINE REGRESSION AND CLASSIFICATION

(Bemporad, 2022)

- **Piecewise Affine Regression and Classification (PARC)** algorithm
- Training dataset:
  - **feature vector**  $z \in \mathbb{R}^n$  (categorical features **one-hot encoded** in  $\{0, 1\}$ )
  - **target vector**  $v_c \in \mathbb{R}^{m_c}$  (numeric),  $v_{di} \in \{w_{di}^1, \dots, w_{di}^{m_i}\}$  (categorical)
- PARC iteratively **clusters** training data in  $K$  sets and **fits** linear predictors:
  1. fit  $v_c = a_j z + b_j$  by **ridge regression** ( $=\ell_2$ -regularized least squares)
  2. fit  $v_{di} = w_{di}^{h_*}$ ,  $h_* = \arg \max \{a_{dih}^h z + b_{di}^h\}$  by **softmax regression**
  3. fit a convex **PWL separation function** by **softmax regression**

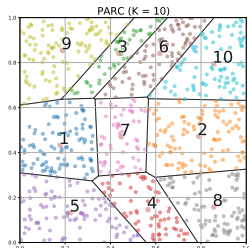
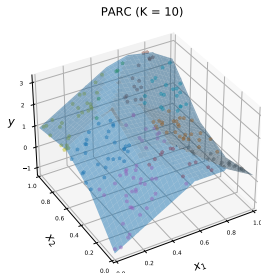
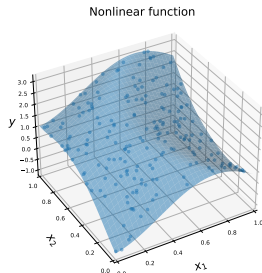
$$\Phi(z) = \omega^{j(z)} z + \gamma^{j(z)}, \quad j(z) = \min \left\{ \arg \max_{j=1, \dots, K} \{\omega^j z + \gamma^j\} \right\}$$


- Data reassigned to clusters based on weighted fit/PWL separation criterion
- PARC is a **block-coordinate descent** algorithm  $\Rightarrow$  (local) convergence ensured



# PARC – PIECEWISE AFFINE REGRESSION AND CLASSIFICATION

- Simple PWA regression example:
  - 1000 samples of  $y = \sin(4x_1 - 5(x_2 - 0.5)^2) + 2x_2$  (use 80% for training)
  - PWA approximation over  $K = 10$  polyhedral regions



• Code:  `pip install pyparc`

[github.com/bemporad/PyPARC](https://github.com/bemporad/PyPARC)

# PARC - CART & BUMPERS EXAMPLE

- **Example:** moving cart and bumpers + heat transfer during bumps.

Spring and viscous forces are **nonlinear**.

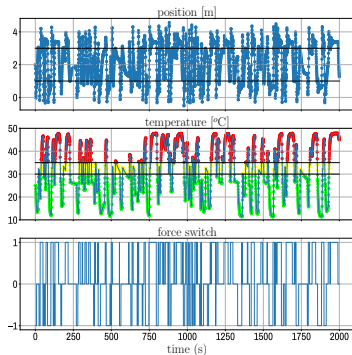
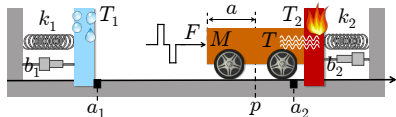
- Categorical input  $F \in \{-\bar{F}, 0, \bar{F}\}$  and categorical output  $c \in \{\text{green}, \text{yellow}, \text{red}\}$

- 4000 training samples

- Feature vector  $z_k = [y_k, \dot{y}_k, T_k, F_k]$

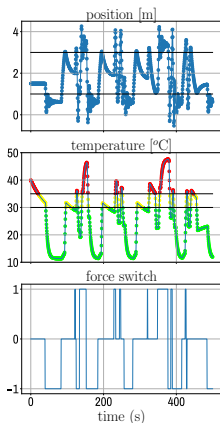
- Target vector  $v_k = [y_{k+1}, \dot{y}_{k+1}, T_{k+1}, c_k]$

- Hybrid model learned by **PARC** ( $K = 5$  regions)

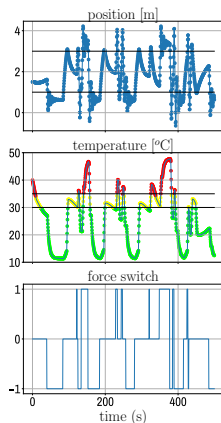


# PARC - CART & BUMPERS EXAMPLE

- **Open-loop** simulation on 500 s **test** data:



continuous-time system



discrete-time PWA model

- Model fit is good enough for MPC design purposes (see later ...)