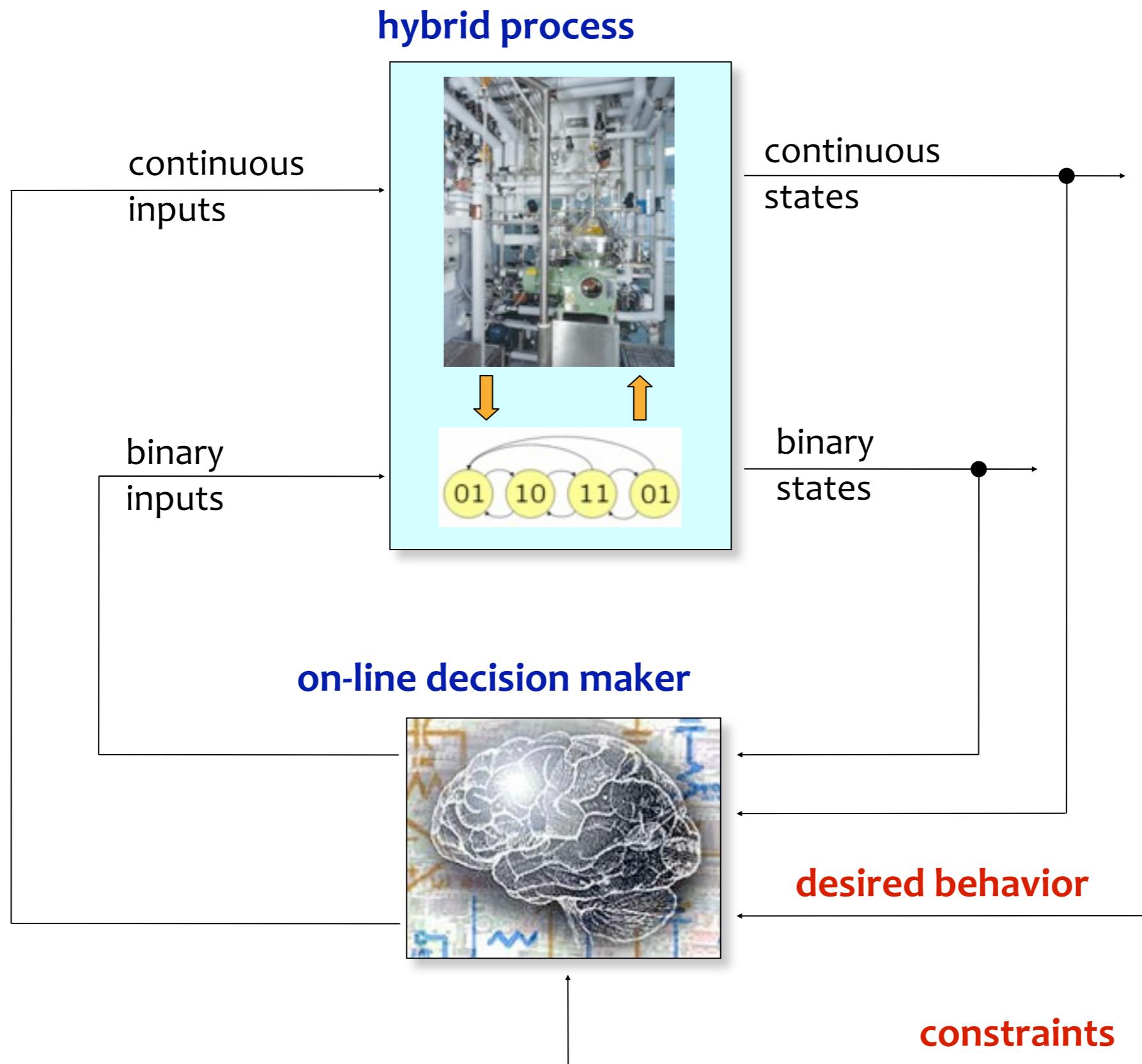


# Model Predictive Control of Hybrid Systems

# Hybrid Control Problem



# Model Predictive Control of Hybrid Systems

MLD model

$$x(t+1) = Ax(t) + B_1u(t) + B_2\delta(t) + B_3z(t)$$

$$y(t) = Cx(t) + D_1u(t) + D_2\delta(t) + D_3z(t)$$

$$E_2\delta(t) + E_3z(t) \leq E_4x(t) + E_1u(t) + E_5$$

Controller

Hybrid System

Reference

$$r(t)$$

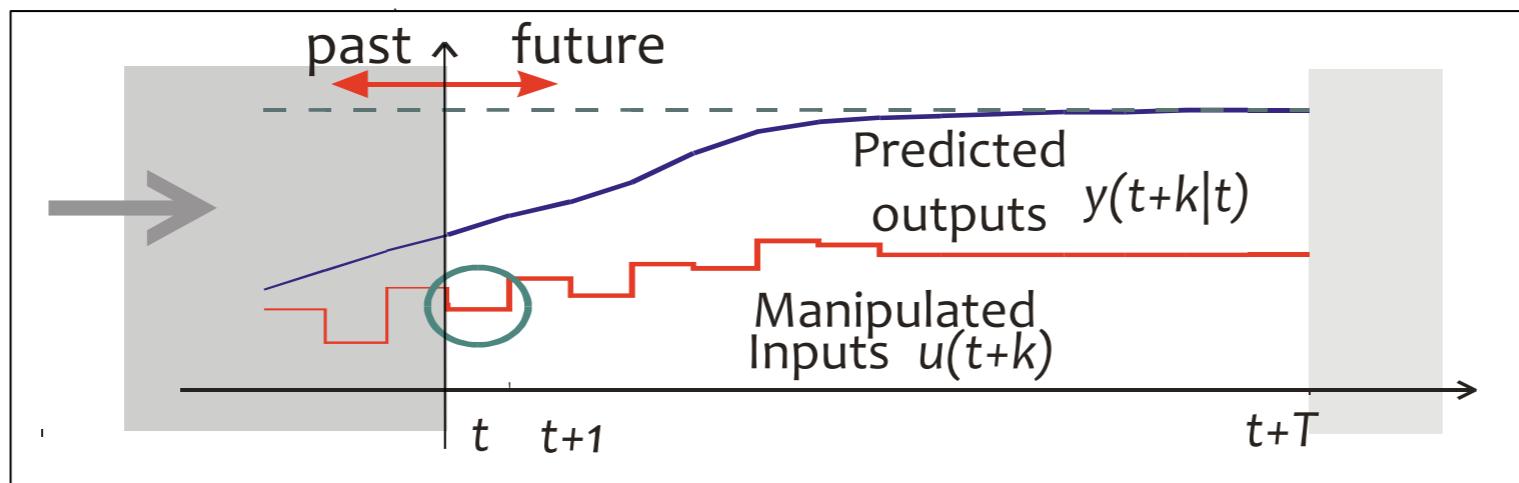
Input  
 $u(t)$

Output  
 $y(t)$

Measurements

- **MODEL:** use an MLD (or PWA) model of the plant to predict the future behavior of the hybrid system
- **PREDICTIVE:** optimization is still based on the predicted future evolution of the hybrid system
- **CONTROL:** the goal is to control the hybrid system

# Hybrid Model Predictive Control



Model  
Predictive (MPC)  
Control

- At time  $t$  solve with respect to  $U \triangleq \{u(t), u(t+1), \dots, u(t+T-1)\}$  the finite-horizon open-loop, optimal control problem:

$$\begin{aligned} \min_{u(t), \dots, u(t+T-1)} \quad & \sum_{k=0}^{T-1} \|y(t+k|t) - r(t)\| + \rho \|u(t+k) - u_r\| \\ & + \sigma (\|\delta(t+k) - \delta_r\| + \|z(t+k) - z_r\| + \|x(t+k|t) - x_r\|) \end{aligned}$$

subject to MLD model

$$x(t|t) = x(t)$$

$$x(t+T|t) = x_r$$

- Apply only  $u(t)=u_t^*$  (discard the remaining optimal inputs)
- At time  $t+1$ : get new measurements, repeat optimization

# Closed-loop Convergence

**Theorem 1** Let  $(x_r, u_r, \delta_r, z_r)$  be the equilibrium values corresponding to the set point  $r$ , and assume  $x(0)$  is such that the MPC problem is feasible at time  $t = 0$ . Then  $\forall Q, R \succ 0, \forall \sigma > 0$

$$\lim_{t \rightarrow \infty} y(t) = r$$

$$\lim_{t \rightarrow \infty} u(t) = u_r$$

$\lim_{t \rightarrow \infty} x(t) = x_r, \lim_{t \rightarrow \infty} \delta(t) = \delta_r, \lim_{t \rightarrow \infty} z(t) = z_r$ ,  
and all constraints are fulfilled.

(Bemporad, Morari 1999)

Proof: Easily follows from standard Lyapunov arguments

More stability results: see (Lazar, Heemels, Weiland, Bemporad, 2006)

# Convergence Proof

- Assume we set the terminal constraint  $x(t + T|t) = x_r$  in the optimal control problem
- Let  $\mathcal{U}_t^*$  denote the optimal control sequence  $\{u_t^*(0), \dots, u_t^*(T-1)\}$
- Let  $V(t) \triangleq J(\mathcal{U}_t^*, x(t))$  = value function  $\rightarrow$  Lyapunov function
- By construction,  $\mathcal{U}_1 = \{u_t^*(1), \dots, u_t^*(T-1), u_r\}$  is feasible @  $t+1$
- Hence,

$$V(t+1) \underset{\text{red circle}}{\leq} J(\mathcal{U}_1, x(t+1)) = V(t) - \|y(t) - r\|_Q - \|u(t) - u_r\|_R - \sigma(\|\delta(t) - \delta_r\| - \|z(t) - z_r\| - \|x(t) - x_r\|)$$

- Hence  $V(t)$  is decreasing and lower-bounded by 0  $\Rightarrow \exists V_\infty = \lim_{t \rightarrow \infty} V(t)$   
 $\Rightarrow V(t+1) - V(t) \rightarrow 0$
- Hence,  $\|y(t) - r\|_Q \rightarrow 0, \|u(t) - u_r\|_R \rightarrow 0, \dots, \|x(t) - x_r\| \rightarrow 0$

**Note: Global optimum not needed for convergence !**

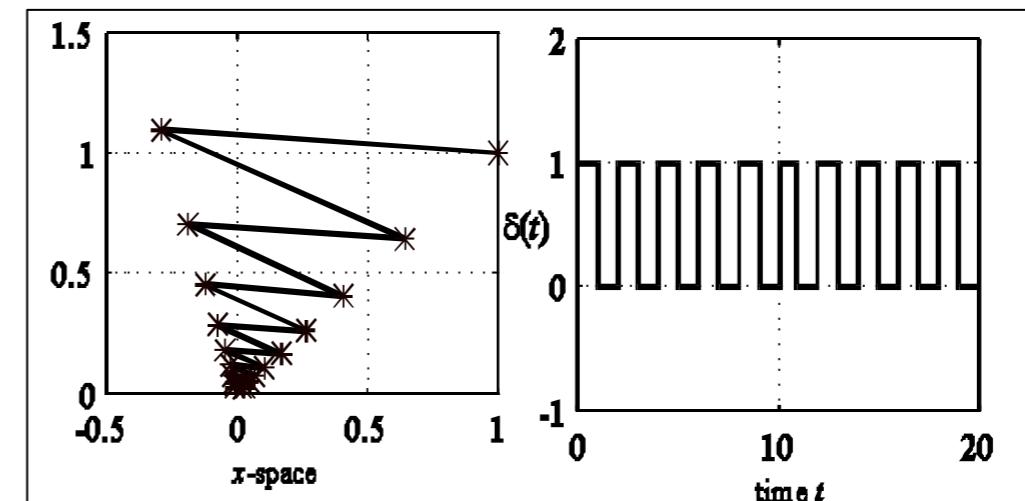
# Hybrid MPC - Example

PWA system:

$$x(t+1) = 0.8 \begin{bmatrix} \cos \alpha(t) & -\sin \alpha(t) \\ \sin \alpha(t) & \cos \alpha(t) \end{bmatrix} x(t) + \begin{bmatrix} 0 \\ 1 \end{bmatrix} u(t)$$
$$y(t) = x_2(t)$$
$$\alpha(t) = \begin{cases} \frac{\pi}{3} & \text{if } x_1(t) > 0 \\ -\frac{\pi}{3} & \text{if } x_1(t) \leq 0 \end{cases}$$

Constraint:  $-1 \leq u(t) \leq 1$

Open loop behavior



go to demo /demos/hybrid/bm99sim.m

# Hybrid MPC - Example

HYSDEL  
model

```
/* 2x2 PWA system - Example from the paper
A. Bemporad and M. Morari, ``Control of systems integrating logic, dynamics,
and constraints,'' Automatica, vol. 35, no. 3, pp. 407-427, 1999.
(C) 2003 by A. Bemporad, 2003 */

SYSTEM pwa {

INTERFACE {
    STATE { REAL x1 [-10,10];
             REAL x2 [-10,10]; }

    INPUT { REAL u [-1.1,1.1]; }

    OUTPUT{ REAL y; }

    PARAMETER {
        REAL alpha = 1.0472; /* 60 deg in radians */
        REAL C = cos(alpha);
        REAL S = sin(alpha); }
    }

IMPLEMENTATION {
    AUX { REAL z1,z2;
          BOOL sign; }

    AD { sign = x1<=0; }

    DA { z1 = { IF sign THEN 0.8*(C*x1+S*x2)
                ELSE 0.8*(C*x1-S*x2) };
         z2 = { IF sign THEN 0.8*(-S*x1+C*x2)
                ELSE 0.8*(S*x1+C*x2) }; }

    CONTINUOUS { x1 = z1;
                 x2 = z2+u; }

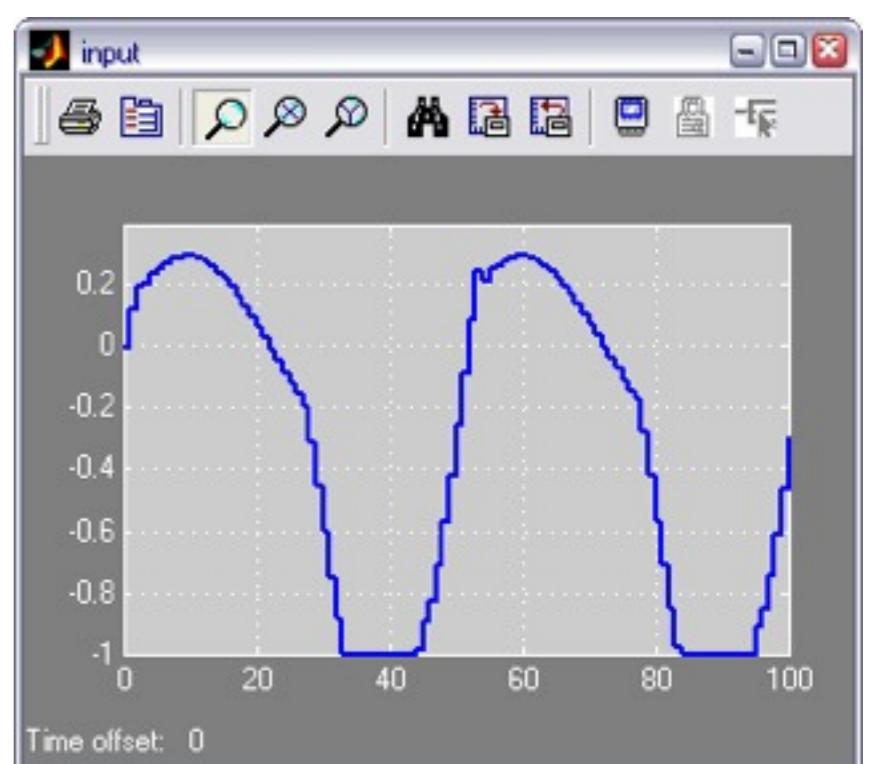
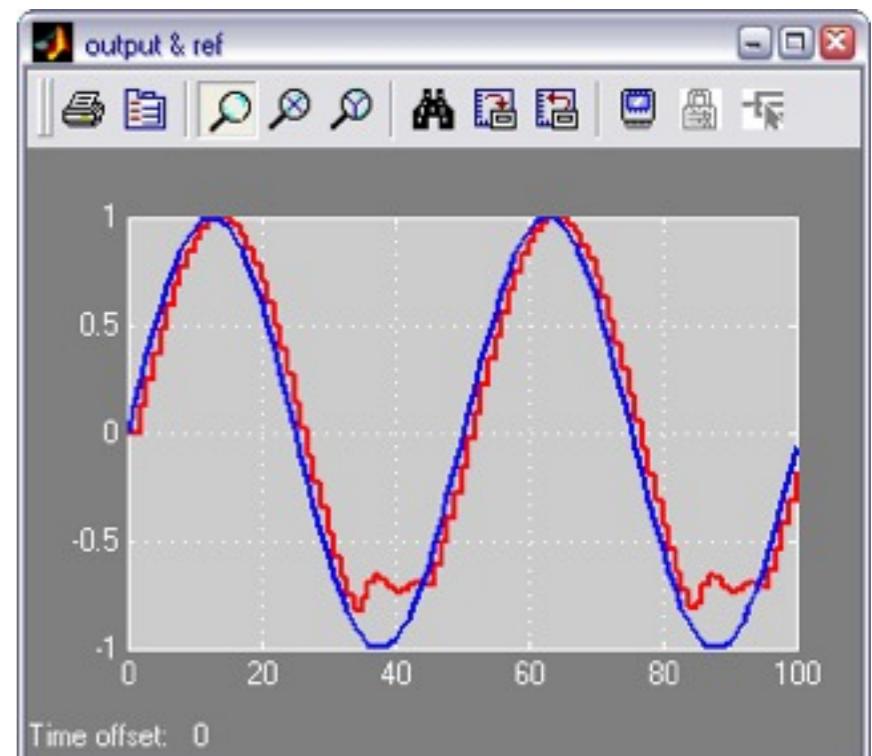
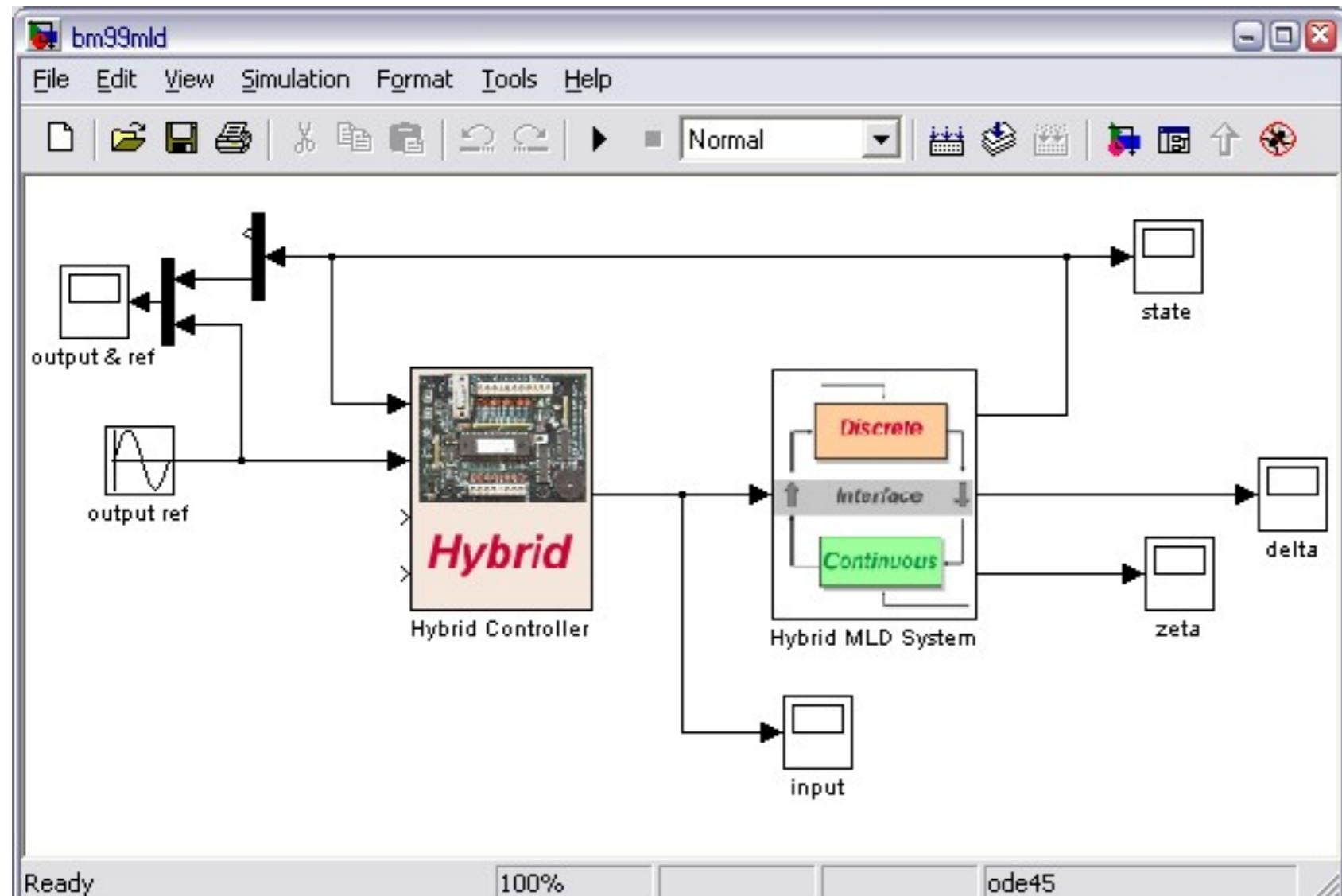
    OUTPUT { y = x2; }
}

}
```

/demos/hybrid/bm99.hys

# Hybrid MPC - Example

Closed-loop:



Performance index:

$$\min \sum_{k=1}^2 |y(t+k|t) - r(t)|$$

# Hybrid MPC – Temperature Control

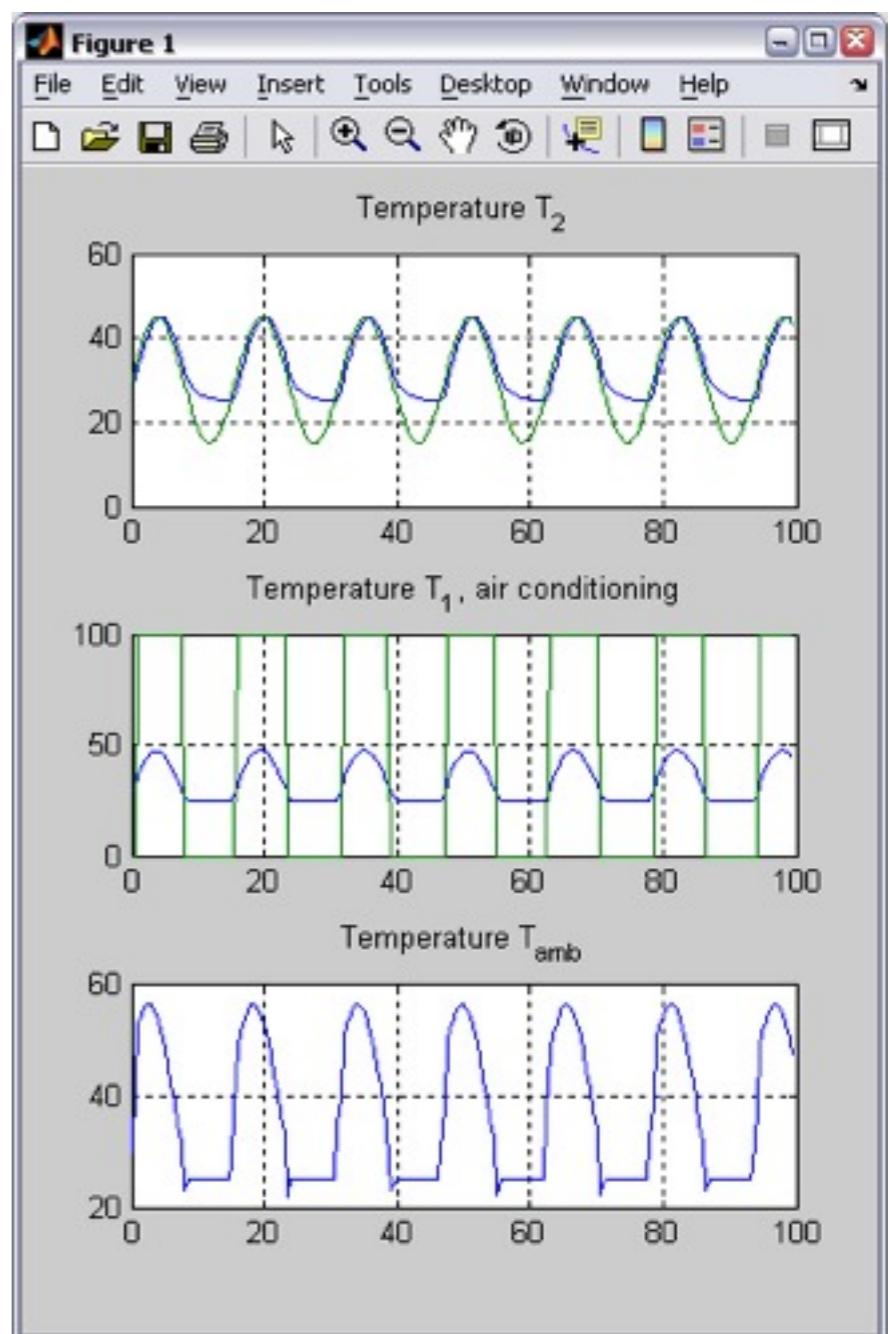
```
>>refs.x=2; % just weight state #2  
>>Q.x=1;  
>>Q.rho=Inf; % hard constraints  
>>Q.norm=2; % quadratic costs  
>>N=2; % optimization horizon  
>>limits.xmin=[25;-Inf];
```

```
>>C=hybcon(S,Q,N,limits,refs);
```

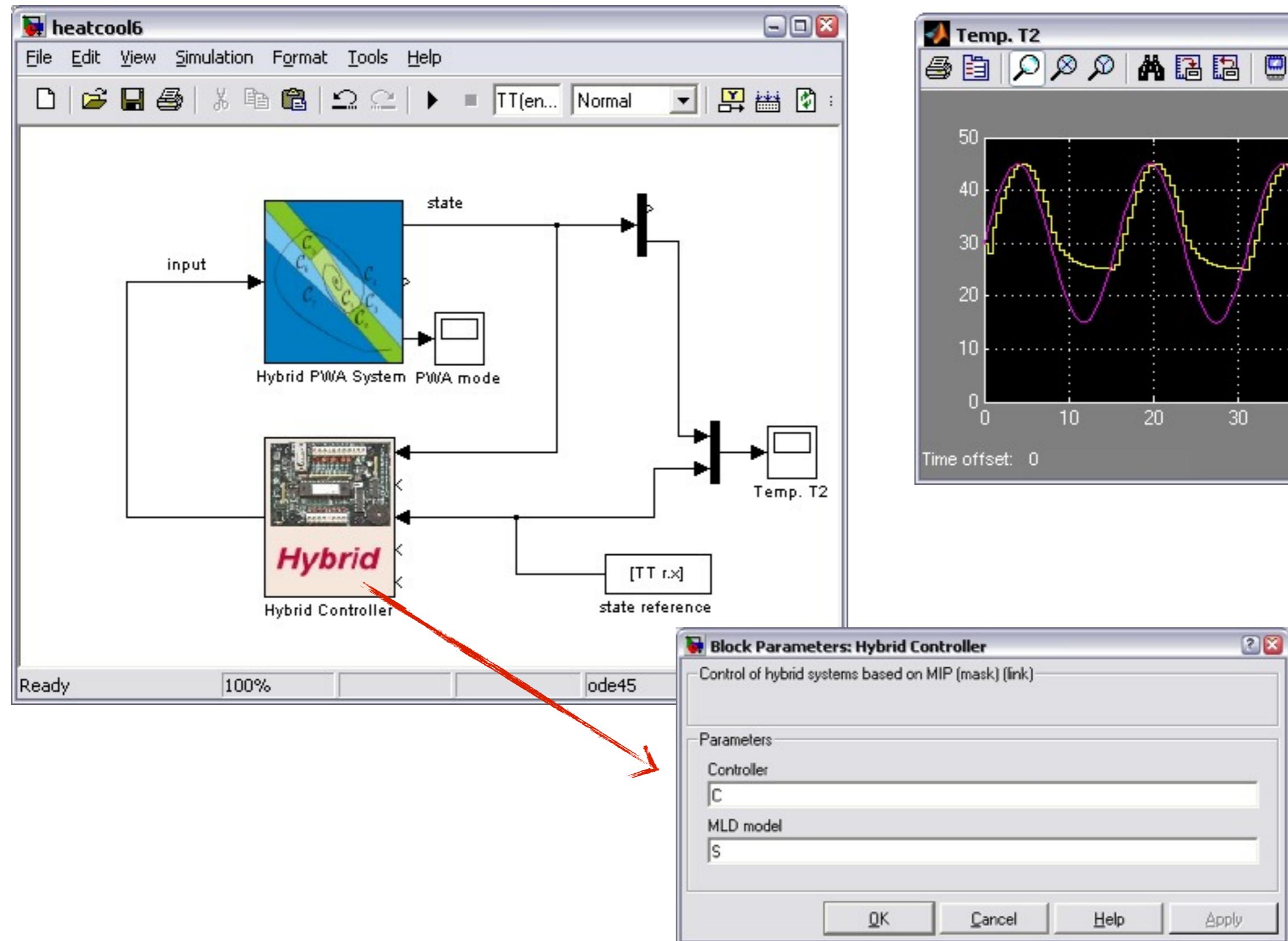
```
>> C  
  
Hybrid controller based on MLD model S <heatcoolmodel.hys>  
  
2 state measurement(s)  
0 output reference(s)  
0 input reference(s)  
1 state reference(s)  
0 reference(s) on auxiliary continuous z-variables  
  
20 optimization variable(s) (8 continuous, 12 binary)  
46 mixed-integer linear inequalities  
sampling time = 0.5, MILP solver = 'glpk'  
  
Type "struct(C)" for more details.  
>>
```

```
>>[XX,UU,DD,ZZ,TT]=sim(C,S,r,x0,Tstop);
```

$$\begin{aligned} \text{min } & \sum_{k=1}^2 (x_2(k) - r)^2 \\ \text{s.t. } & x_1(k) \geq 25 \quad k = 1, 2 \\ & \text{MLD model} \end{aligned}$$



# Hybrid MPC – Temperature Control



# Optimal Control of Hybrid Systems: Computational Aspects



# MIQP Formulation of MPC

(Bemporad, Morari, 1999)

$$\begin{aligned} \min_{\xi} J(\xi, x(0)) &= \sum_{t=0}^{T-1} y'(t) Q y(t) + u'(t) R u(t) \\ \text{subject to } & \begin{cases} x(t+1) = Ax(t) + B_1 u(t) + B_2 \delta(t) + B_3 z(t) + B_5 \\ y(t) = Cx(t) + D_1 u(t) + D_2 \delta(t) + D_3 z(t) + D_5 \\ E_2 \delta(t) + E_3 z(t) \leq E_4 x(t) + E_1 u(t) + E_5 \end{cases} \end{aligned}$$

- Optimization vector:

$$\xi = [u(0), \dots, u(T-1), \delta(0), \dots, \delta(T-1), z(0), \dots, z(T-1)]'$$

→ 
$$\begin{aligned} \min_{\xi} & \frac{1}{2} \xi' H \xi + x(0)' F \xi + \frac{1}{2} x'(0)' Y x(0) \\ \text{subj. to } & G \xi \leq W + S x(0) \end{aligned}$$

**Mixed Integer Quadratic Program (MIQP)**

$$u \in \mathbb{R}^{n_u}, \delta \in \{0, 1\}^{n_\delta}, z \in \mathbb{R}^{n_z}$$



$$\xi \in \mathbb{R}^{(n_u+n_z)T} \times \{0, 1\}^{n_\delta T}$$

$\xi$  has both real and  $\{0, 1\}$  components

# MIQP Formulation of MPC

(Bemporad, Borrelli, Morari, 2000)

$$\min_{\xi} J(\xi, x(0)) = \sum_{t=0}^{T-1} \|Qy(t)\|_{\infty} + \|Ru(t)\|_{\infty}$$

subject to MLD model

- Basic trick: introduce slack variables

$$\min_x |x| \rightarrow \begin{array}{ll} \min_{x, \epsilon} & \epsilon \\ \text{s.t.} & \epsilon \geq x \\ & \epsilon \geq -x \end{array}$$

$$\left\{ \begin{array}{ll} \epsilon_k^x \geq \|Qy(t+k|t)\|_{\infty} & i = 1, \dots, p \quad k = 0, \dots, T-1 \\ \epsilon_k^u \geq \|Ru(t+k)\|_{\infty} & i = 1, \dots, m \quad k = 0, \dots, T-1 \end{array} \right. \rightarrow \left\{ \begin{array}{ll} \epsilon_k^x \geq Q^i y(t+k|t) & i = 1, \dots, p \quad k = 0, \dots, T-1 \\ \epsilon_k^x \geq -Q^i y(t+k|t) & i = 1, \dots, p \quad k = 0, \dots, T-1 \\ \epsilon_k^u \geq R^i u(t+k) & i = 1, \dots, m \quad k = 0, \dots, T-1 \\ \epsilon_k^u \geq -R^i u(t+k) & i = 1, \dots, m \quad k = 0, \dots, T-1 \end{array} \right.$$

- Optimization vector:

$$\xi = [\epsilon_1^x, \dots, \epsilon_{T-1}^x, \epsilon_0^u, \dots, \epsilon_{T-1}^u, u(0), \dots, u(T-1), \delta(0), \dots, \delta(T-1), z(0), \dots, z(T-1)]'$$

$Q^i = i$ th row of matrix  $Q$

$$\min_{\xi} J(\xi, x(0)) = \sum_{k=0}^{T-1} \epsilon_k^x + \epsilon_k^u$$

s.t.  $G\xi \leq W + Sx(0)$

**Mixed Integer Linear Program (MILP)**

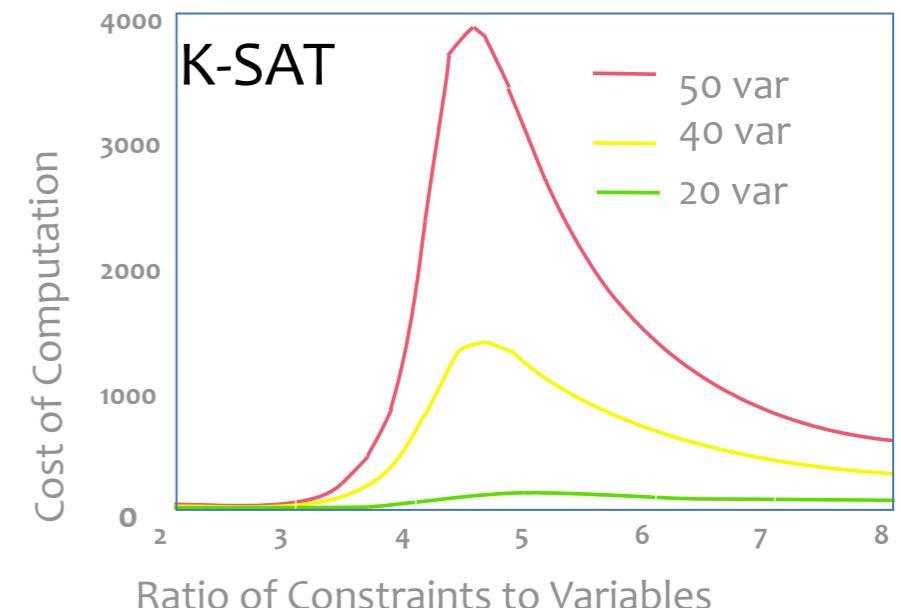
$\xi$  has both real and  $\{0, 1\}$  components

# Mixed-Integer Program Solvers

- Mixed-Integer Programming is NP-complete

*Phase transitions* have been found in computationally hard problems.

BUT



(Monasson et al., Nature, 1999)

- General purpose Branch & Bound/Branch & Cut solvers available for **MILP** and **MIQP** (CPLEX, Xpress-MP, BARON, GLPK, ...)

More solvers and benchmarks: <http://plato.la.asu.edu/bench.html>

- No need to reach global optimum (see proof of the theorem), although performance deteriorates

# Solving Mixed-Integer Programs

$$\begin{aligned} \min \quad & f'x + d'\delta \\ \text{s.t.} \quad & Ax + B\delta \leq c \\ & x \in \mathbb{R}^n, \quad \delta \in \{0, 1\}^m \end{aligned}$$

$$\begin{aligned} \min \quad & \frac{1}{2} [\begin{matrix} x \\ \delta \end{matrix}]' H [\begin{matrix} x \\ \delta \end{matrix}] + f'x + d'\delta \\ \text{s.t.} \quad & Ax + B\delta \leq c \\ & x \in \mathbb{R}^n, \quad \delta \in \{0, 1\}^m \end{aligned}$$

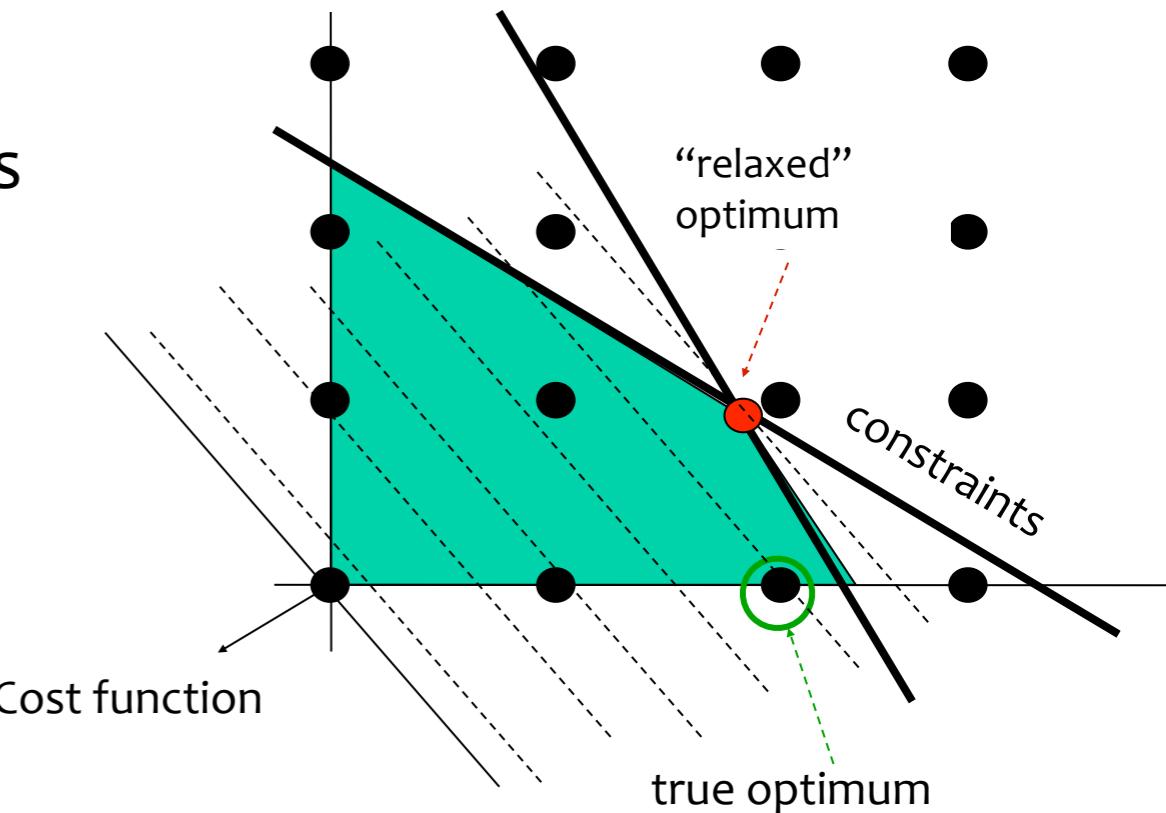
- Some variables are continuous, some are discrete (0/1)
- In general, it is a NP-Hard problem
- Naive solution: enumerate all possible integer solutions and choose the best one.

But  $m$  binary variables lead to  $2^m$  solutions, each of which requires a LP

→ Impossible but for extremely small  $m$  !

# Branch & Bound Algorithm

1. Solve a “**relaxed**” problem with all binary variables treated as continuous,  $0 \leq \delta_i \leq 1$ . This gives a (lower) bound on the “best possible” solution. Unfortunately, some  $\delta_i$  may have fractional parts.
2. **Branch** on one binary variable: set  $\delta_j=0$  and  $\delta_j=1$  in two separate solutions, for some  $j$
3. Use a **bound** on the optimal cost to eliminate in one shot a large number of combinations that are certainly not leading to the optimum (=advantage over full enumeration)
4. Branch again on another variable, and so on, until no further branching is possible.



# A Simple Example in Supply Chain Management

manufacturer A



manufacturer B



manufacturer C



inventory 1



$x_{11}(k), x_{12}(k)$

$u_{11}(k)$

$u_{12}(k)$

retailer 1



$y_1(k)$

$y_2(k)$

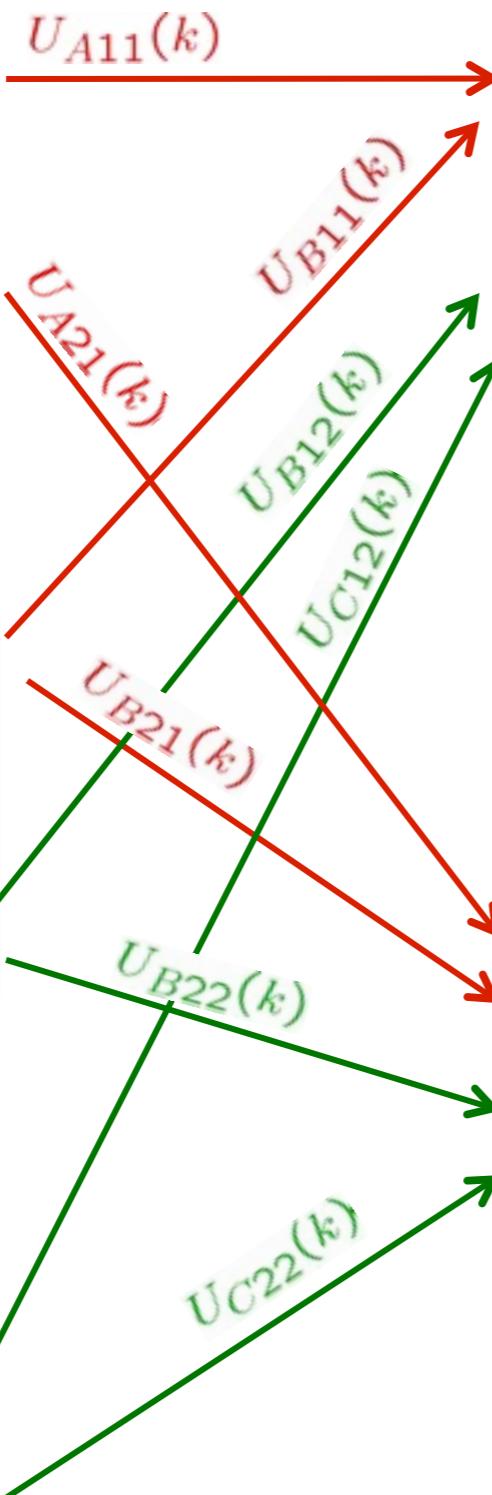
inventory 2



$x_{21}(k), x_{22}(k)$

$u_{22}(k)$

$u_{21}(k)$



# System Variables

- continuous states:

$x_{ij}(k)$  = amount of  $j$  hold in inventory  $i$  at time  $k$  ( $i=1,2, j=1,2$ )

- continuous outputs:

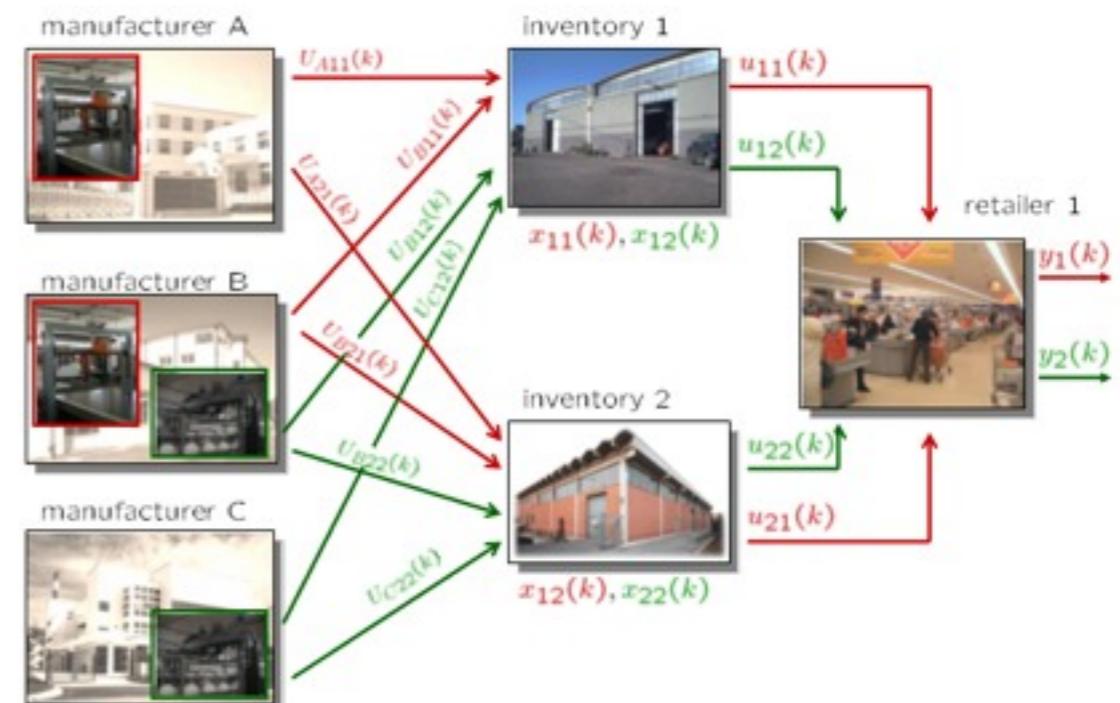
$y_j(k)$  = amount of  $j$  sold at time  $k$  ( $j=1,2$ )

- continuous inputs:

$u_{ij}(k)$  = amount of  $j$  taken from inventory  $i$  at time  $k$  ( $i=1,2, j=1,2$ )

- binary inputs:

$U_{Xij}(k)$  = 1 if manufacturer  $X$  produces and send  $j$  to inventory  $i$  at time  $k$



# Constraints

- Max capacity of inventory  $i$ :

$$0 \leq \sum_j x_{ij}(k) \leq x_{Mi}$$

Numerical values:

$$x_{M1}=10, x_{M2}=10$$

- Max transportation from inventories:

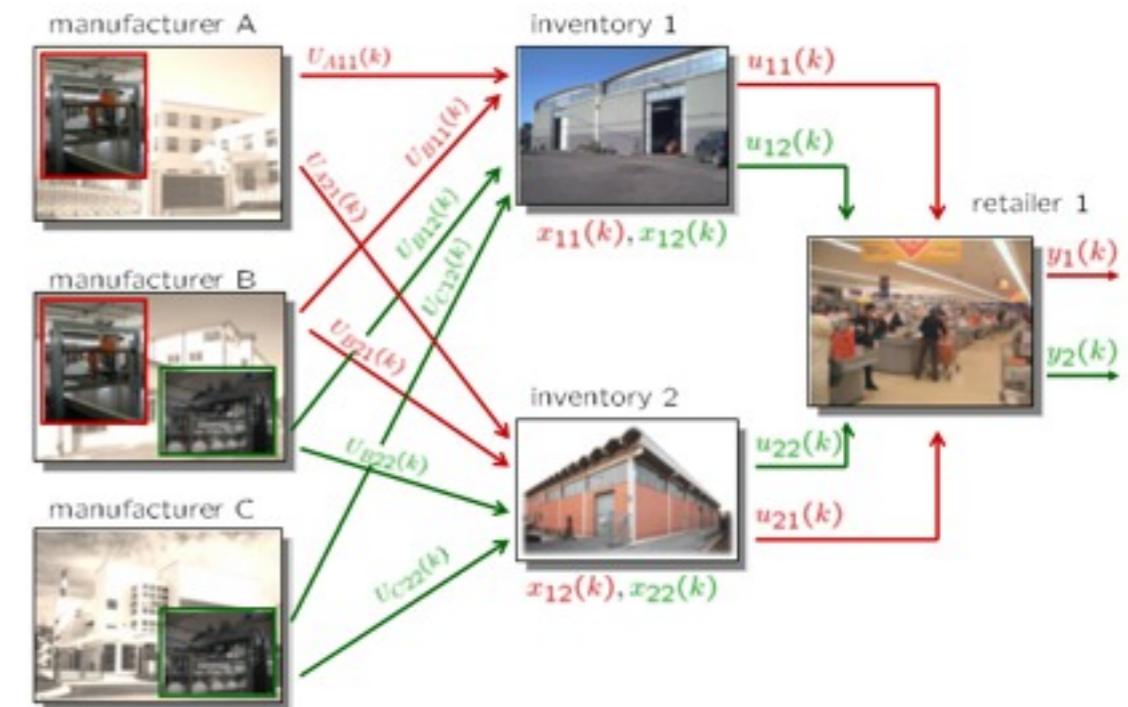
$$0 \leq u_{ij}(k) \leq u_M$$

- A product can only be sent to one inventory:

$UA11(k)$  and  $UA21(k)$  cannot be =1 at the same time

$UB11(k)$  and  $UB21(k)$  cannot be =1 at the same time  $UB12(k)$  and  $UB22(k)$  cannot be =1 at the same time

$UC12(k)$  and  $UC22(k)$  cannot be =1 at the same time



- A manufacturer can only produce one type of product at one time:

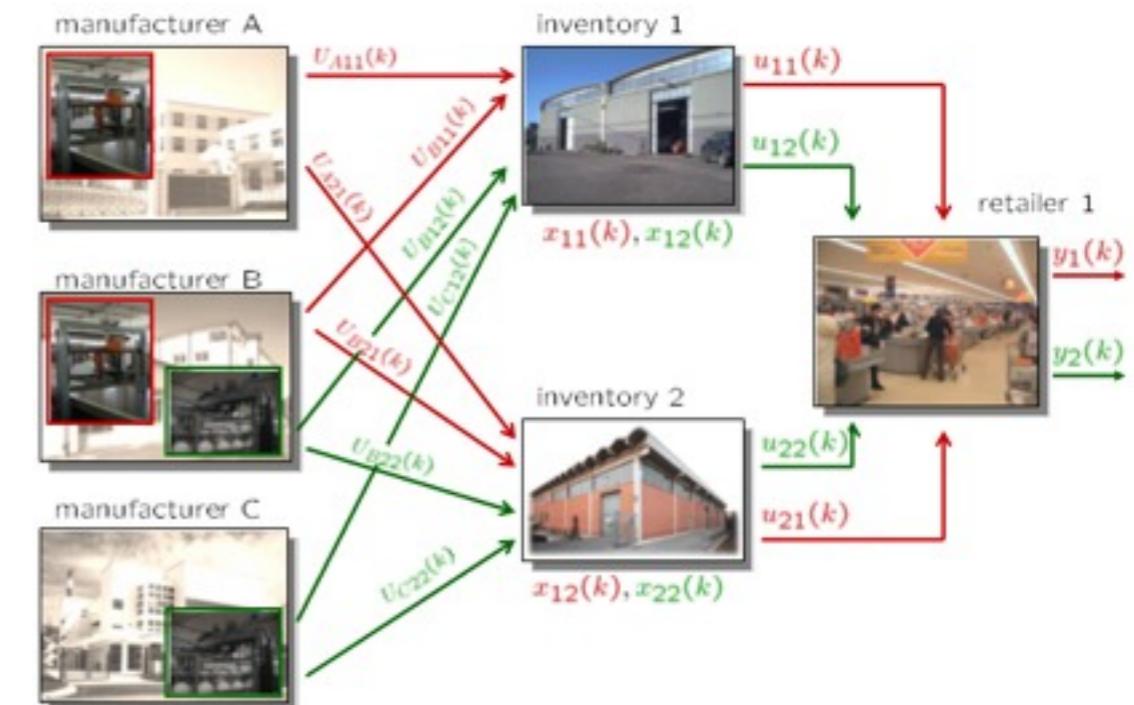
$[UB11(k)=1 \text{ or } UB21(k)=1]$  and  $[UB12(k)=1 \text{ or } UB22(k)=1]$  cannot be true

# Dynamics

$P_{A1}, P_{B1}, P_{B2}, P_{C2}$  = amount of type 1(2) produced by  $A$  ( $B,C$ ) in one time interval

Numerical values:

$$P_{A1}=4, P_{B1}=6, P_{B2}=7, P_{C2}=3$$



- Level of inventories:

$$\begin{cases} x_{11}(k+1) = x_{11}(k) + P_{A1}U_{A11}(k) + P_{B1}U_{B11}(k) - u_{11}(k) \\ x_{12}(k+1) = x_{12}(k) + P_{B2}U_{B12}(k) + P_{C2}U_{C12}(k) - u_{12}(k) \\ x_{21}(k+1) = x_{21}(k) + P_{A1}U_{A21}(k) + P_{B1}U_{B21}(k) - u_{21}(k) \\ x_{22}(k+1) = x_{22}(k) + P_{B2}U_{B22}(k) + P_{C2}U_{C22}(k) - u_{22}(k) \end{cases}$$

# Hybrid Dynamical Model

```

SYSTEM supply_chain{
INTERFACE {
    STATE { REAL x11 [0,10];
             REAL x12 [0,10];
             REAL x21 [0,10];
             REAL x22 [0,10]; }

    INPUT { REAL u11 [0,10];
            REAL u12 [0,10];
            REAL u21 [0,10];
            REAL u22 [0,10];
            BOOL UA11,UA21,UB11,UB12,UB21,UB22,UC12,UC22; }

    OUTPUT {REAL y1,y2; }

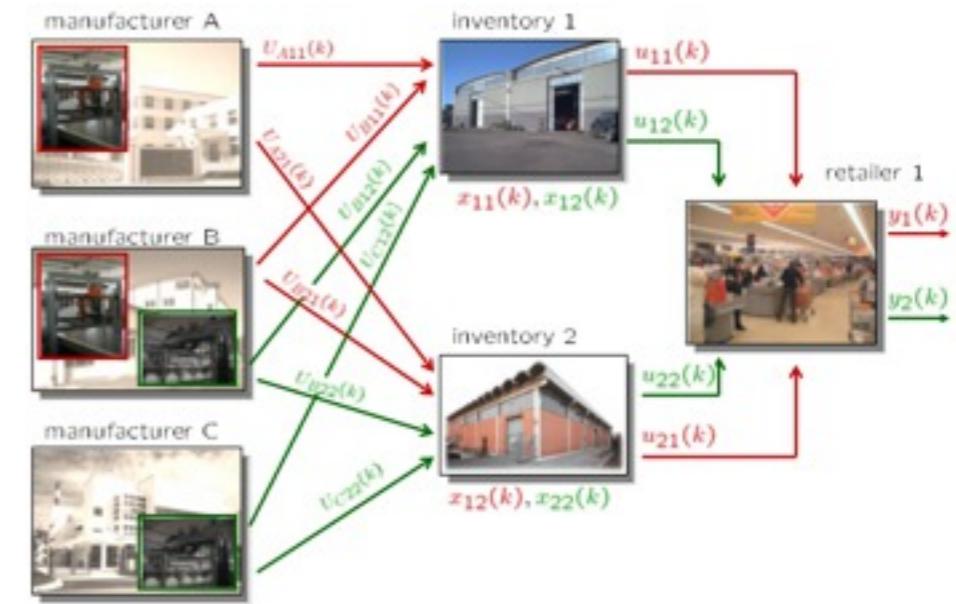
    PARAMETER { REAL PA1,PB1,PB2,PC2,xM1,xM2; }
}

IMPLEMENTATION {

AUX { REAL zA11, zB11, zB12, zC12, zA21, zB21, zB22, zC22; }

DA { zA11 = {IF UA11 THEN PA1 ELSE 0};
     zB11 = {IF UB11 THEN PB1 ELSE 0};
     zB12 = {IF UB12 THEN PB2 ELSE 0};
     zC12 = {IF UC12 THEN PC2 ELSE 0};
     zA21 = {IF UA21 THEN PA1 ELSE 0};
     zB21 = {IF UB21 THEN PB1 ELSE 0};
     zB22 = {IF UB22 THEN PB2 ELSE 0};
     zC22 = {IF UC22 THEN PC2 ELSE 0}; }
}
}

```



```

CONTINUOUS {x11 = x11 + zA11 + zB11 - u11;
            x12 = x12 + zB12 + zC12 - u12;
            x21 = x21 + zA21 + zB21 - u21;
            x22 = x22 + zB22 + zC22 - u22; }

OUTPUT { y1 = u11 + u21;
          y2 = u12 + u22; }

MUST { ~ (UA11 & UA21);
        ~ (UC12 & UC22);
        ~ ((UB11 | UB21) & (UB12 |
UB22));
        ~ (UB11 & UB21);
        ~ (UB12 & UB22);
        x11+x12 <= xM1;
        x11+x12 >=0;
        x21+x22 <= xM2;
        x21+x22 >=0; }

} }

```

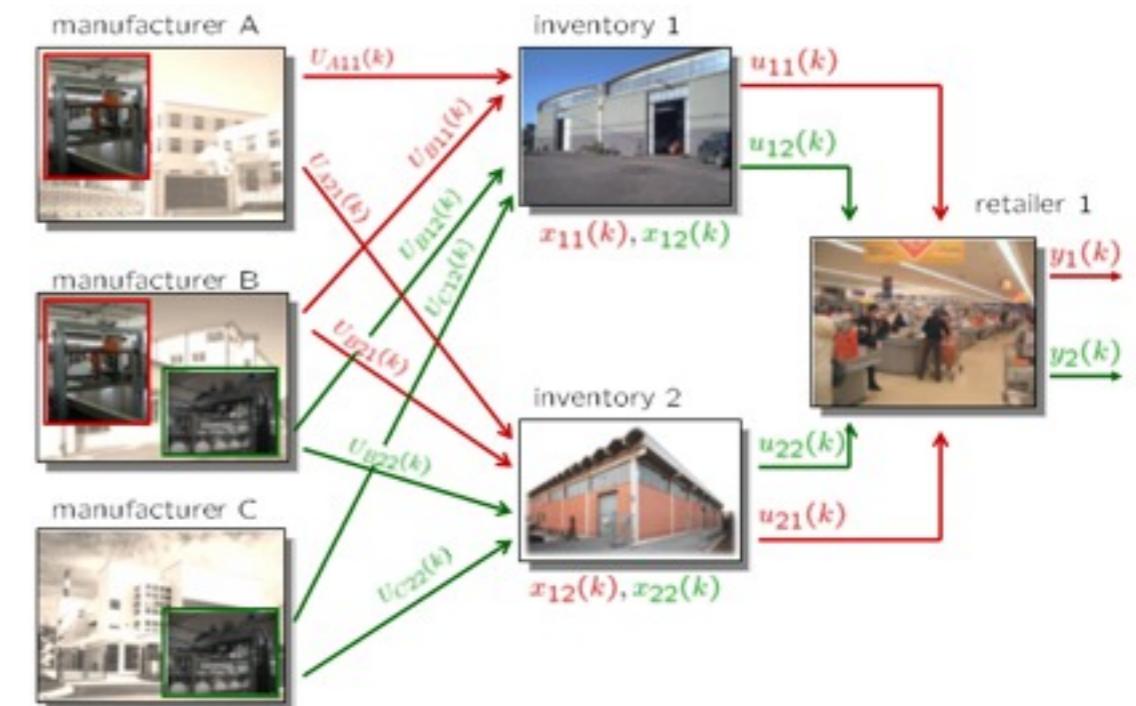
/demos/hybrid/supply\_chain.m

# Objectives

- Meet customer demand as much as possible:  $y_1 \approx r_1, y_2 \approx r_2$

- Minimize transportation costs

- Fulfill all constraints



# Performance Specs

$$\min \sum_{k=0}^{N-1} \left( 10(|y_1(k) - r_1(k)| + |y_2(k) - r_2(k)|) + \right.$$

*penalty on demand tracking error*

$$+ 4(|u_{11}(k)| + |u_{12}(k)|) +$$

*cost for shipping  
from inv.#1 to market*

$$+ 2(|u_{21}(k)| + |u_{22}(k)|) +$$

*cost for shipping  
from inv.#2 to market*

$$+ 1(|U_{A11}(k)| + |U_{A21}(k)|) +$$

*cost from A to  
inventories*

$$+ 4(|U_{B11}(k)| + |U_{B12}(k)| + |U_{B21}(k)| + |U_{B22}(k)|) +$$

*cost from B to  
inventories*

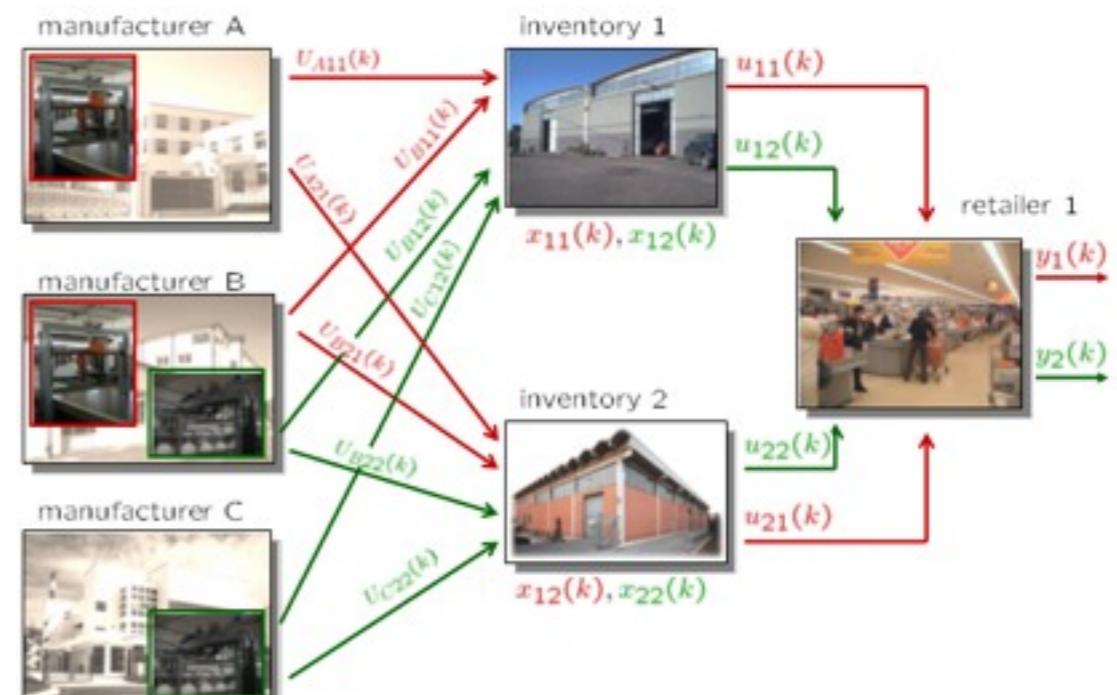
$$\left. + 10(|U_{C12}(k)| + |U_{C22}(k)|) \right)$$

*cost from C to  
inventories*

# Simulation setup

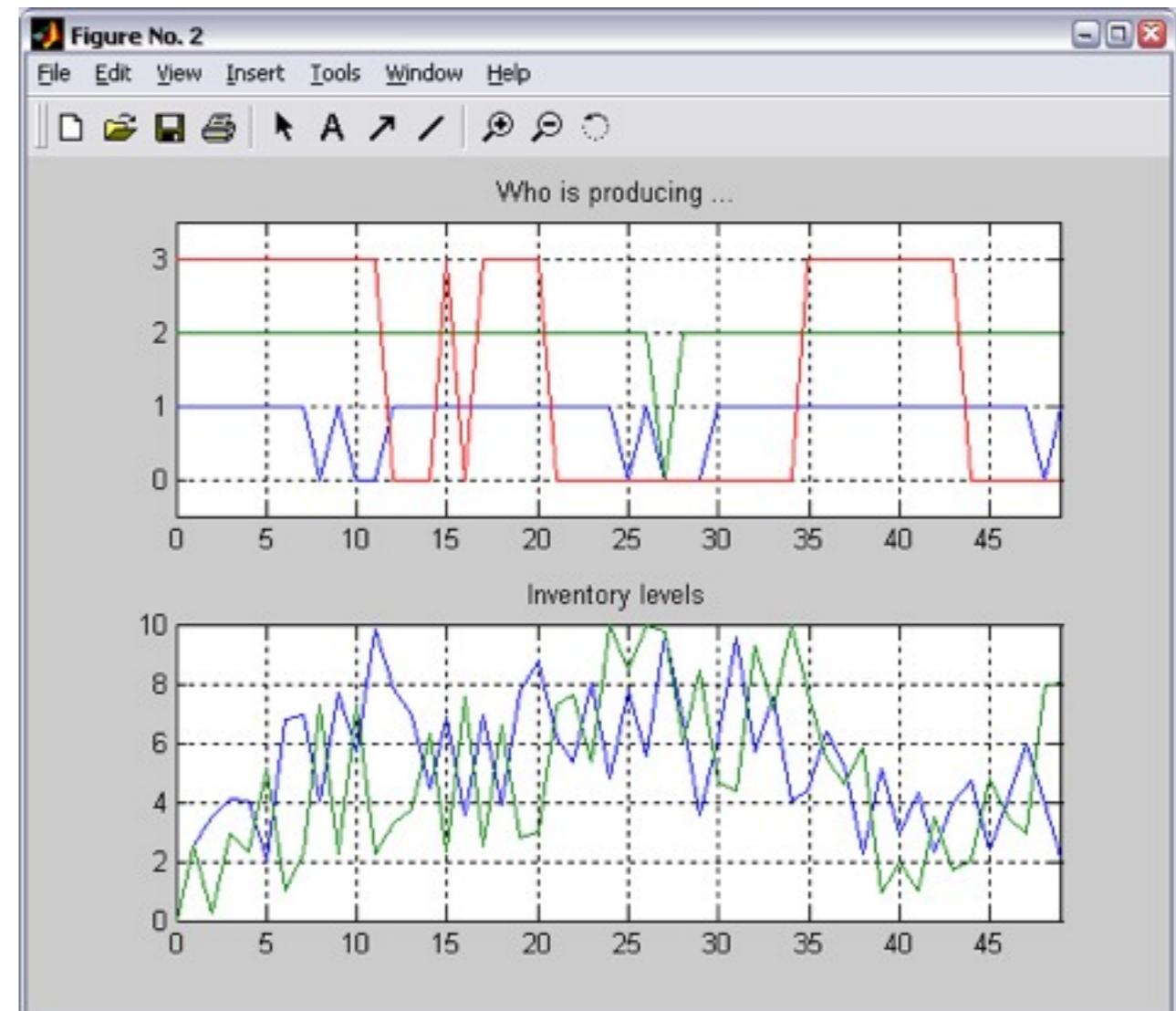
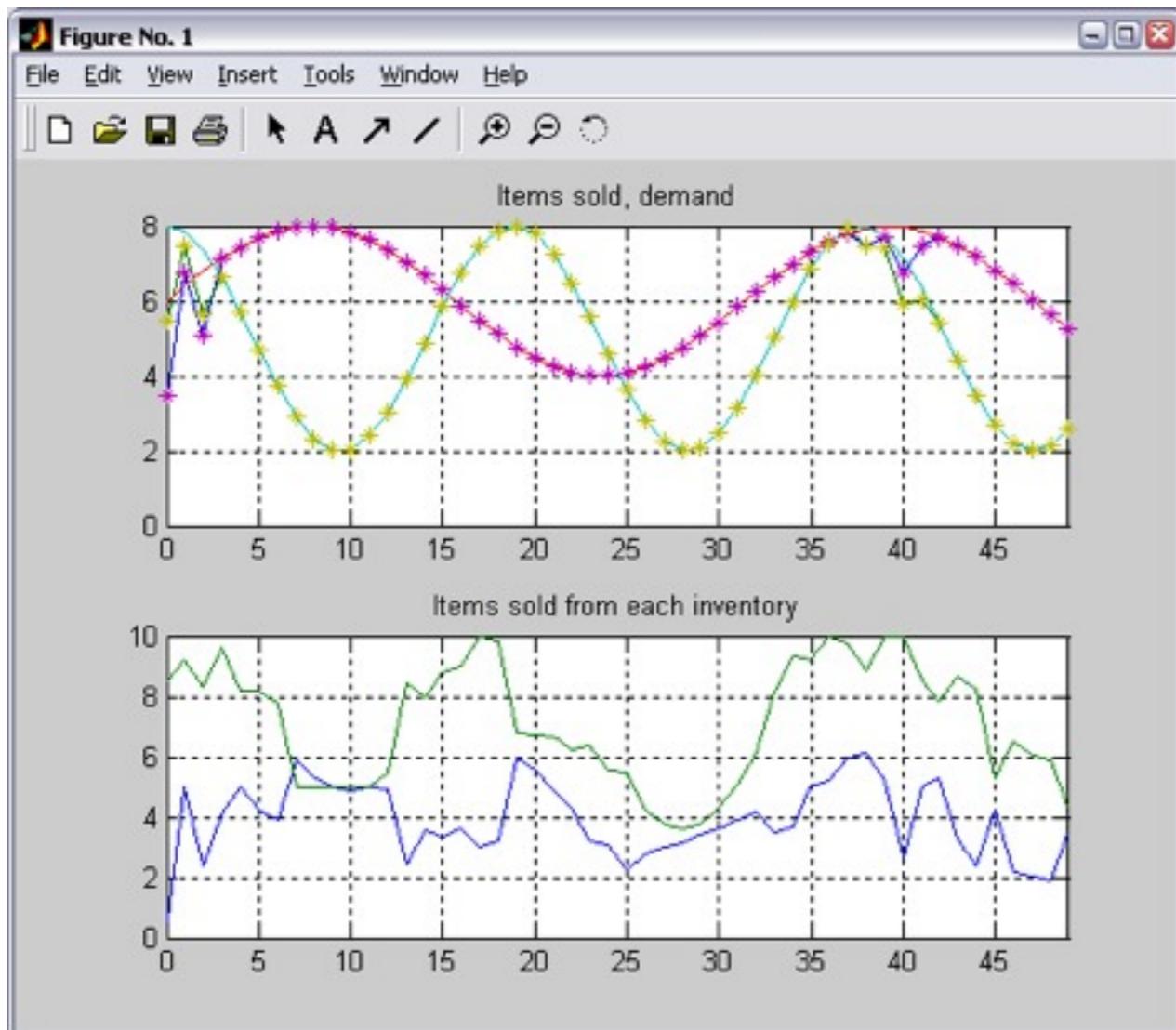
```
>>refs.y=[1 2]; % weights output2 #1,#2  
>>Q.y=diag([10 10]); % output weights  
...  
>>Q.norm=Inf; % infinity norms  
>>N=2; % optimization horizon  
>>limits.umin=umin; % constraints  
>>limits.umax=umax;  
>>limits.xmin=xmin;  
>>limits xmax=xmax;  
  
>>C=hybcon(S,Q,N,limits,refs);
```

```
>> C  
  
Hybrid controller based on MLD model S <supply_chain.hys> [Inf-norm]  
  
4 state measurement(s)  
2 output reference(s)  
12 input reference(s)  
0 state reference(s)  
0 reference(s) on auxiliary continuous z-variables  
  
44 optimization variable(s) (28 continuous, 16 binary)  
176 mixed-integer linear inequalities  
sampling time = 1, MILP solver = 'glpk'  
  
Type "struct(C)" for more details.  
  
>>
```



# Simulation results

```
>>x0=[0;0;0;0]; % Initial condition  
>>r.y=[6+2*sin((0:Tstop-1)'/5)  
      5+3*cos((0:Tstop-1)'/3)]; % Reference trajectories  
  
>>[XX,UU,DD,ZZ,TT]=sim(C,S,r,x0,Tstop);
```



CPU time:  $\approx$  30ms per time step (using GLPK on this machine)

# Explicit Hybrid MPC

# Explicit Hybrid MPC (MLD)

$$\min_{\xi} J(\xi, \boxed{x(t)}, \boxed{r(t)}) = \sum_{k=0}^{T-1} \|Q(y_k - \boxed{r(t)})\|_{\infty} + \|Ru_k\|_{\infty}$$

subject to

$$\begin{cases} x_{k+1} = Ax_k + B_1u_k + B_2\delta_k + B_3z_k + B_5 \\ y_k = Cx_k + D_1u_k + D_2\delta_k + D_3z_k + D_5 \\ E_2\delta_k + E_3z_k \leq E_4x_k + E_1u_k + E_5 \\ x_0 = \boxed{x(t)} \end{cases}$$

- On-line optimization: solve the problem **for each given  $x(t)$**

Mixed-Integer Linear Program (MILP)

- Off-line optimization: solve the MILP **for all  $x(t)$  in advance**

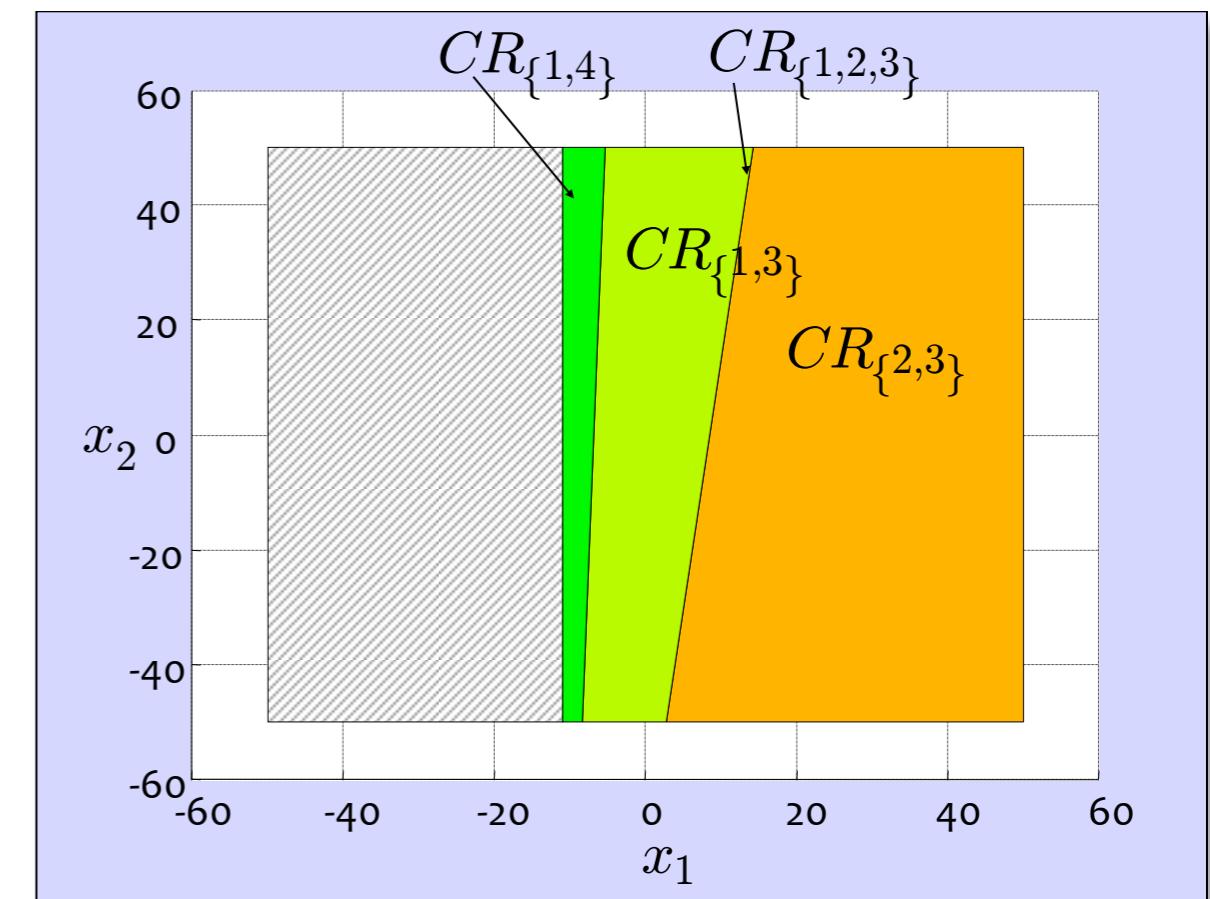
$$\begin{aligned} \min_{\xi} & \sum_{k=0}^{T-1} \epsilon_i^x + \epsilon_i^u \\ \text{s.t. } & G\xi \leq W + S \begin{bmatrix} \boxed{x(t)} \\ r(t) \end{bmatrix} \end{aligned}$$

multi-parametric Mixed Integer Linear Program (mp-MILP)

# Example of Multiparametric Solution

## Multiparametric LP

$$\begin{aligned} \min_{\xi} \quad & -3\xi_1 - 8\xi_2 \\ \text{s.t.} \quad & \left\{ \begin{array}{lcl} \xi_1 + \xi_2 & \leq & 13 + x_1 \\ 5\xi_1 - 4\xi_2 & \leq & 20 \\ -8\xi_1 + 22\xi_2 & \leq & 121 + x_2 \\ -4\xi_1 - \xi_2 & \leq & -8 \\ -\xi_1 & \leq & 0 \\ -\xi_2 & \leq & 0 \end{array} \right. \end{aligned}$$



$$\xi(x) = \begin{cases} [0.00 \ 0.05]x + [11.85] & \text{if } \begin{bmatrix} 0.02 & 0.00 \\ 0.00 & 0.02 \\ 0.00 & -0.02 \\ -0.12 & 0.01 \end{bmatrix}x \leq \begin{bmatrix} 1.00 \\ 1.00 \\ 1.00 \\ -1.00 \end{bmatrix} \quad CR_{\{2,3\}} \\ [0.73 \ -0.03]x + [5.50] & \text{if } \begin{bmatrix} 0.00 & 0.02 \\ 0.00 & -0.02 \\ 0.12 & -0.01 \\ -0.15 & 0.00 \end{bmatrix}x \leq \begin{bmatrix} 1.00 \\ 1.00 \\ 1.00 \\ 1.00 \end{bmatrix} \quad CR_{\{1,3\}} \\ [-0.33 \ 0.00]x + [-1.67] & \text{if } \begin{bmatrix} 0.00 & 0.02 \\ 0.00 & -0.02 \\ 0.15 & -0.00 \\ -0.09 & 0.00 \end{bmatrix}x \leq \begin{bmatrix} 1.00 \\ 1.00 \\ -1.00 \\ 1.00 \end{bmatrix} \quad CR_{\{1,4\}} \end{cases}$$

# Multiparametric MILP

$$\begin{aligned} \min_{\xi=\{\xi_c, \xi_d\}} \quad & f' \xi_c + d' \xi_d \\ \text{s.t.} \quad & G \xi_c + E \xi_d \leq W + Fx \end{aligned}$$

$$\xi_c \in \mathbb{R}^n$$

$$\xi_d \in \{0, 1\}^m$$

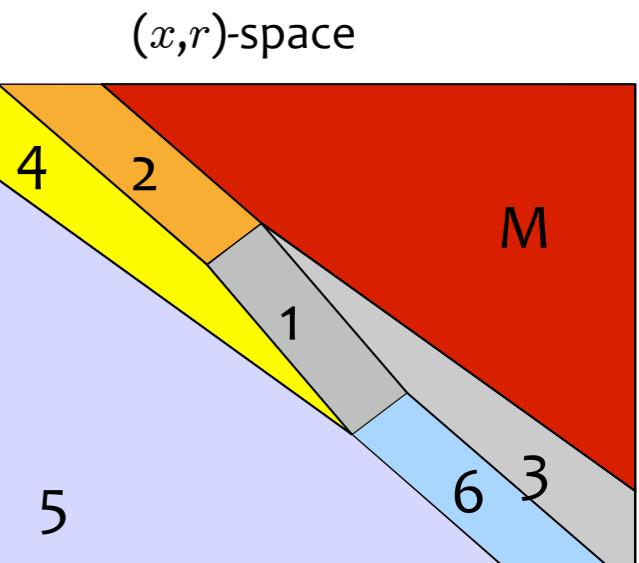
- mp-MILP can be solved (by alternating MILPs and mp-LPs)

(Dua, Pistikopoulos, 1999)

- **Theorem:** The multiparametric solution is piecewise  $\xi^*(x)$

- The MPC controller is piecewise affine in  $x, r$

$$u(x, r) = \begin{cases} F_1 x + E_1 r + g_1 & \text{if } H_1 [\frac{x}{r}] \leq K_1 \\ \vdots & \vdots \\ F_M x + E_M r + g_M & \text{if } H_M [\frac{x}{r}] \leq K_M \end{cases}$$



# Explicit Hybrid MPC (PWA)

$$\min_U J(U, \underline{x}, \underline{r}) = \sum_{k=0}^{T-1} \|R(y(k) - \underline{r})\|_p + \|Qu(k)\|_p$$

subject to  $\begin{cases} \text{PWA model} \\ x(0) = \underline{x} \end{cases}$

$$p = 1, 2, \infty$$

$$\|v\|_2 = v'v$$

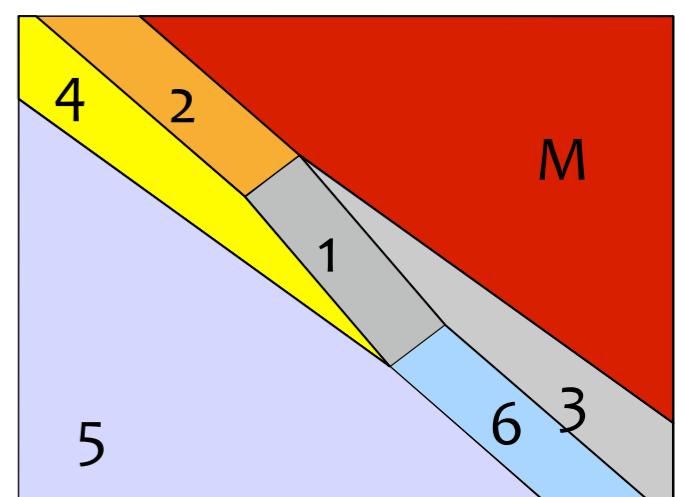
$$\|v\|_\infty = \max |v_i|$$

$$\|v\|_1 = \sum v_i$$

- The MPC controller is piecewise affine in  $x, r$

$$u(x, r) = \begin{cases} F_1x + E_1r + g_1 & \text{if } H_1 [\begin{matrix} x \\ r \end{matrix}] \leq K_1 \\ \vdots & \vdots \\ F_Mx + E_Mr + g_M & \text{if } H_M [\begin{matrix} x \\ r \end{matrix}] \leq K_M \end{cases}$$

$(x, r)$ -space



Note: in the 2-norm case the partition may not be fully polyhedral

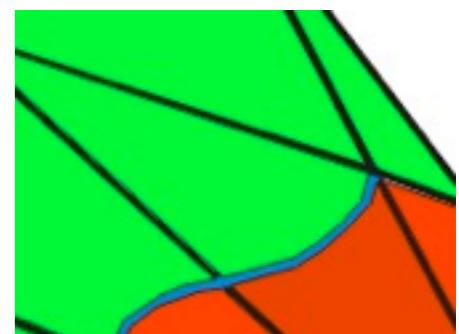
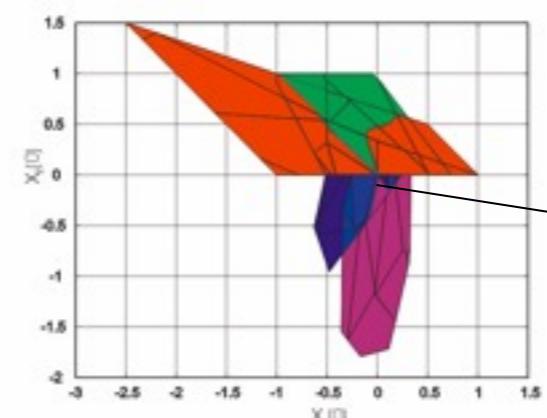
# Computation of Explicit Hybrid MPC (PWA)

Method A: (Borrelli, Baotic, Bemporad, Morari, *Automatica*, 2005)

Use a combination of **DP (dynamic programming)** and **mpLP**  
(1-norm,  $\infty$ -norm), or **mpQP** (quadratic forms)

Method B: (Bemporad, *Hybrid Toolbox*, 2003) (Alessio, Bemporad, ADHS 2006)(Mayne, ECC 2001)

- 1 - Use backwards (=DP) **reachability analysis** for enumerating all feasible mode sequences  $I = \{ i(0), i(1), \dots, i(T) \}$ ;
- 2 - For each fixed sequence  $I$ , solve the explicit finite-time optimal control problem for the corresponding linear time-varying system (**mpQP** or **mpLP**);
- 3 - Case 1/ $\infty$ -norm: Compare value functions and split regions.  
Quadratic case: keep overlapping regions (possibly eliminate overlaps that are never optimal) and compare on-line (if needed).



Note: in the 2-norm case, the fully explicit partition may not be polyhedral

# Hybrid Control Examples (Revisited)

# Hybrid control example

PWA system:

$$x(t+1) = 0.8 \begin{bmatrix} \cos \alpha(t) & -\sin \alpha(t) \\ \sin \alpha(t) & \cos \alpha(t) \end{bmatrix} x(t) + \begin{bmatrix} 0 \\ 1 \end{bmatrix} u(t)$$

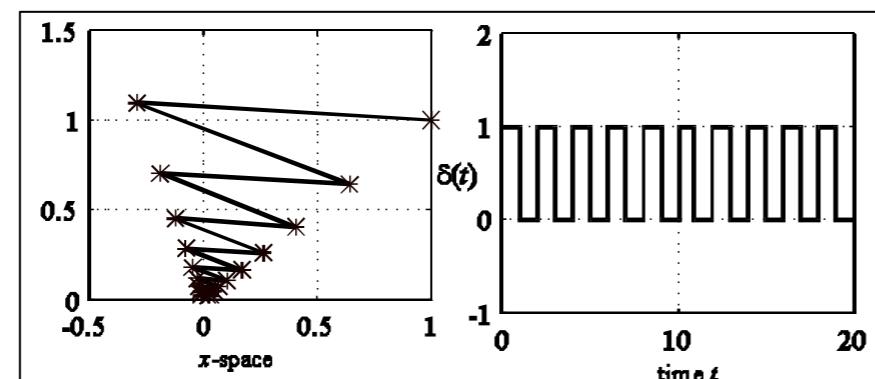
$$y(t) = x_2(t)$$

$$\alpha(t) = \begin{cases} \frac{\pi}{3} & \text{if } x_1(t) \geq 0 \\ -\frac{\pi}{3} & \text{if } x_1(t) < 0 \end{cases}$$

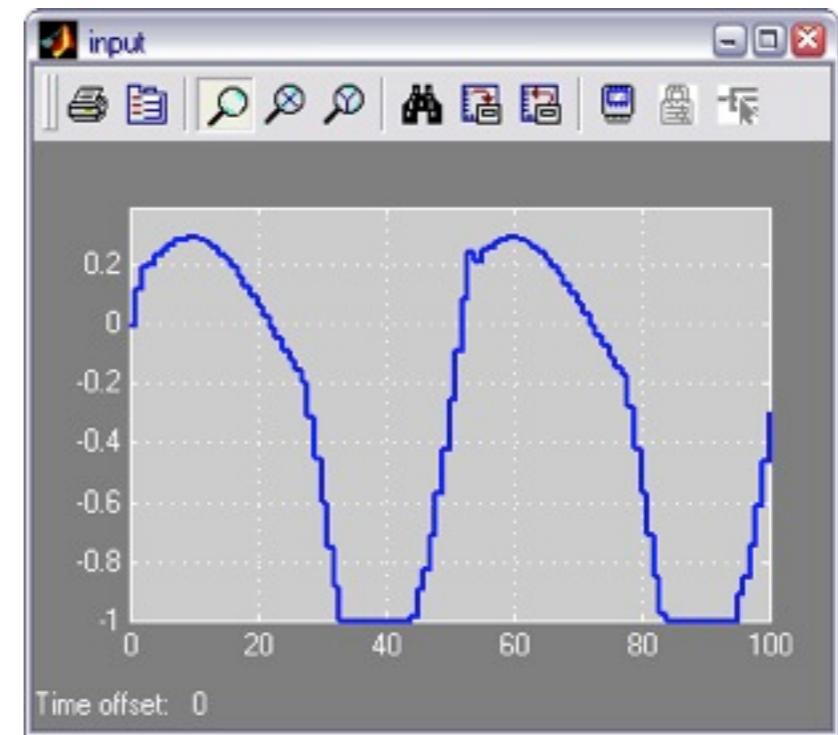
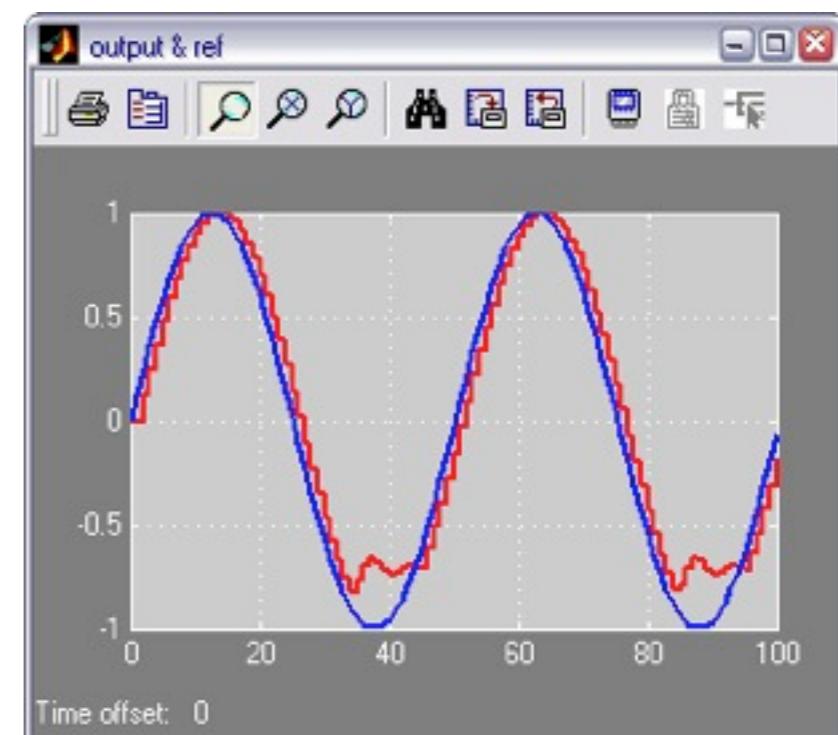
Constraints:  $-1 \leq u(t) \leq 1$

Objective:  $\min \sum_{k=1}^2 |y(t+k|t) - r(t)|$

Open loop behavior:



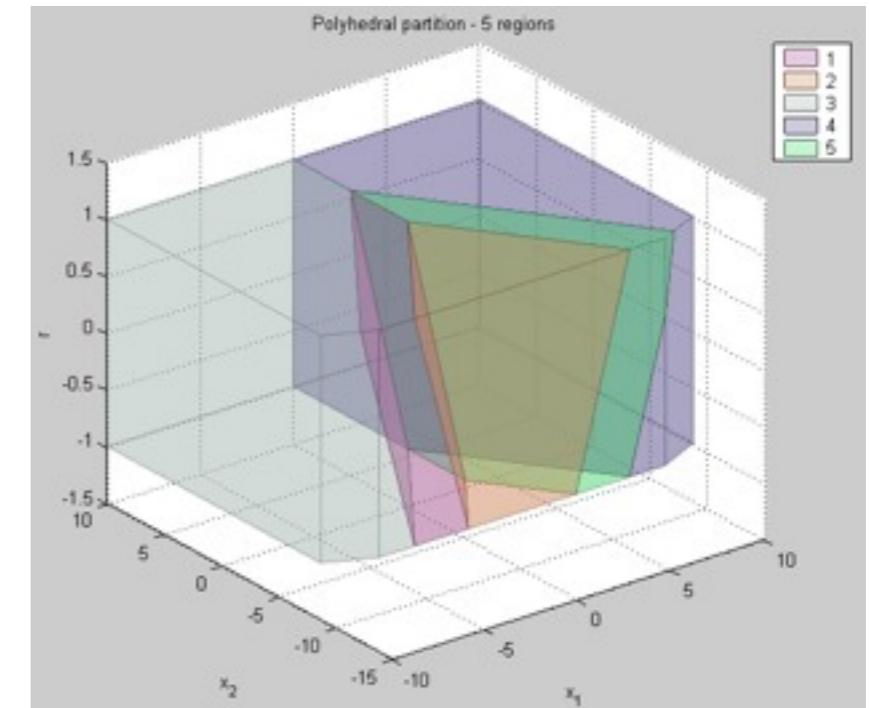
Closed loop:



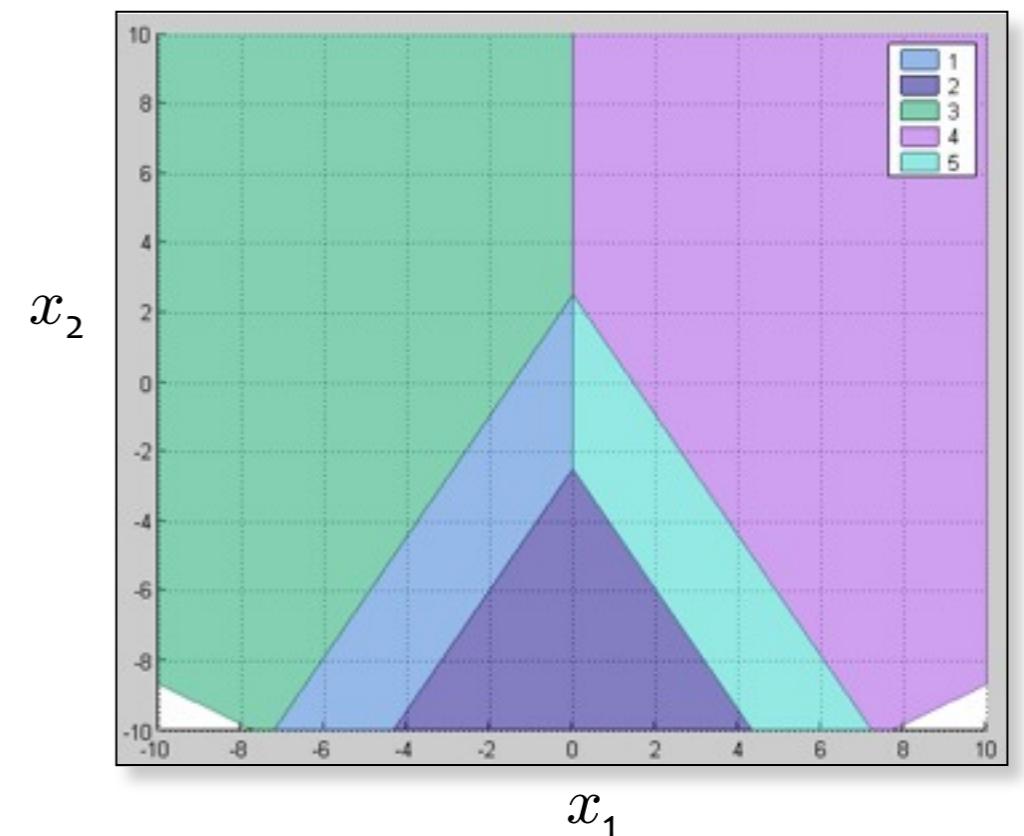
HybTbx: [/demos/hybrid/bm99sim.m](#)

# Explicit PWA Controller

$$u(x, r) = \begin{cases} [0.6928 -0.4 1] \begin{bmatrix} x \\ r \end{bmatrix} & \text{if } \begin{bmatrix} 0.6928 & -0.4 & 1 \\ -0.4 & -0.6928 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -1 \\ -0.6928 & 0.4 & -1 \\ 0 & 0 & 1 \\ 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} 1 \\ 10 \\ 1 \\ 1 \\ 1 \\ 1e-006 \\ 10 \end{bmatrix} \\ & \text{(Region \#1)} \\ 1 & \text{if } \begin{bmatrix} -0.6928 & 0.4 & -1 \\ 0.6928 & 0.4 & -1 \\ 0 & 0 & 1 \\ 0 & 0 & -1 \\ 0 & -1 & 0 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} -1 \\ -1 \\ 1 \\ 1 \\ 1 \\ 10 \end{bmatrix} \\ & \text{(Region \#2)} \\ -1 & \text{if } \begin{bmatrix} -0.4 & -0.6928 & 0 \\ 0 & -1 & 0 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0.6928 & -0.4 & 1 \\ 0 & 0 & -1 \\ -1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} 10 \\ 10 \\ 1e-006 \\ 10 \\ -1 \\ 1 \\ 10 \end{bmatrix} \\ & \text{(Region \#3)} \\ -1 & \text{if } \begin{bmatrix} -1 & 0 & 0 \\ 0 & 1 & 0 \\ 0.4 & -0.6928 & 0 \\ 0 & -1 & 0 \\ -0.6928 & -0.4 & 1 \\ 0 & 0 & 1 \\ 0 & 0 & -1 \\ 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} 0 \\ 10 \\ 10 \\ 10 \\ -1 \\ 1 \\ 10 \end{bmatrix} \\ & \text{(Region \#4)} \\ [-0.6928 -0.4 1] \begin{bmatrix} x \\ r \end{bmatrix} & \text{if } \begin{bmatrix} -0.6928 & -0.4 & 1 \\ 0.4 & -0.6928 & 0 \\ 0 & -1 & 0 \\ 0 & 0 & -1 \\ 0.6928 & 0.4 & -1 \\ 0 & 0 & 1 \\ -1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} 1 \\ 10 \\ 1 \\ 1 \\ 1 \\ 1 \\ 0 \end{bmatrix} \\ & \text{(Region \#5)} \end{cases}$$



Section with  $r=0$

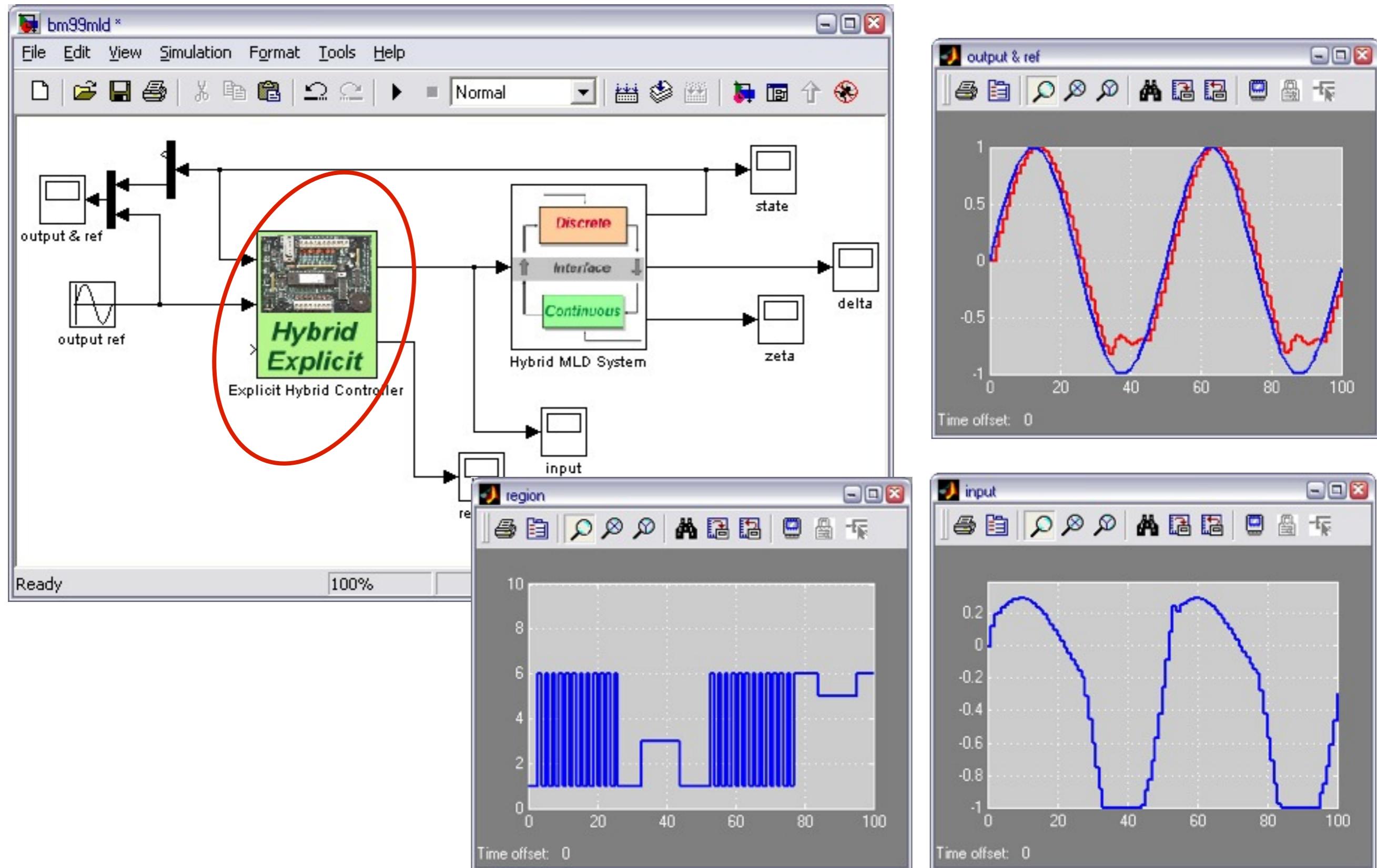


HybTbx: **/demos/hybrid/bm99sim.m**  
 (CPU time: 1.51 s, Pentium M 1.4GHz)

PWA law  $\equiv$  MPC law !

# Hybrid control example

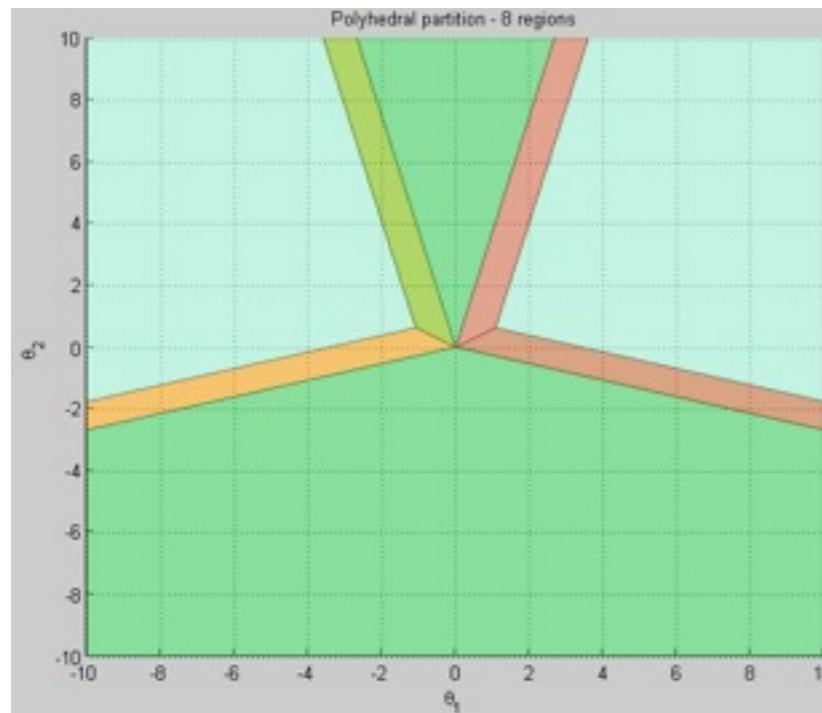
Closed loop:



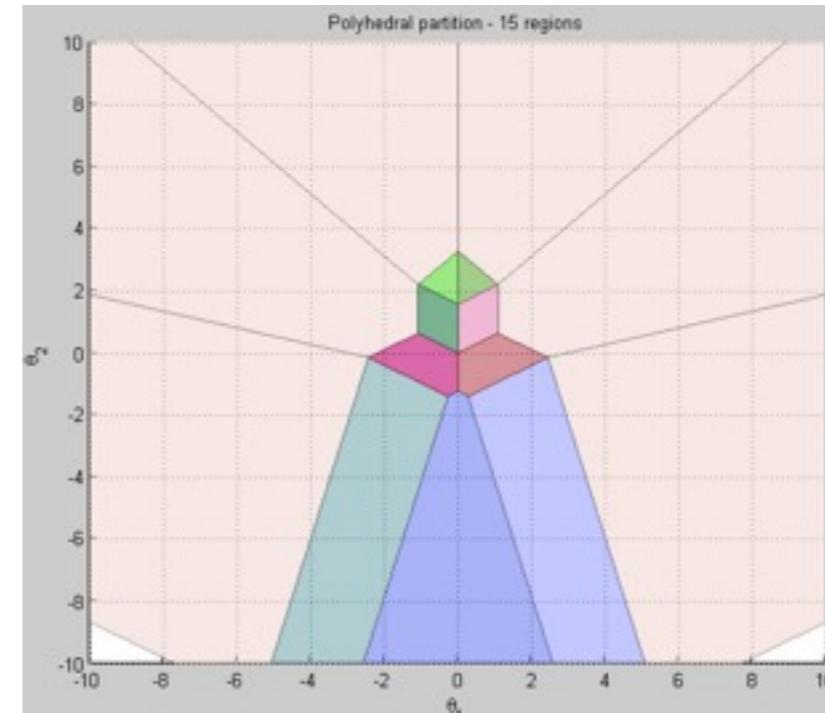
# Explicit PWA Regulator

Objective:  $\min \sum_{k=1}^N \|x(t+k|t)\|_\infty$

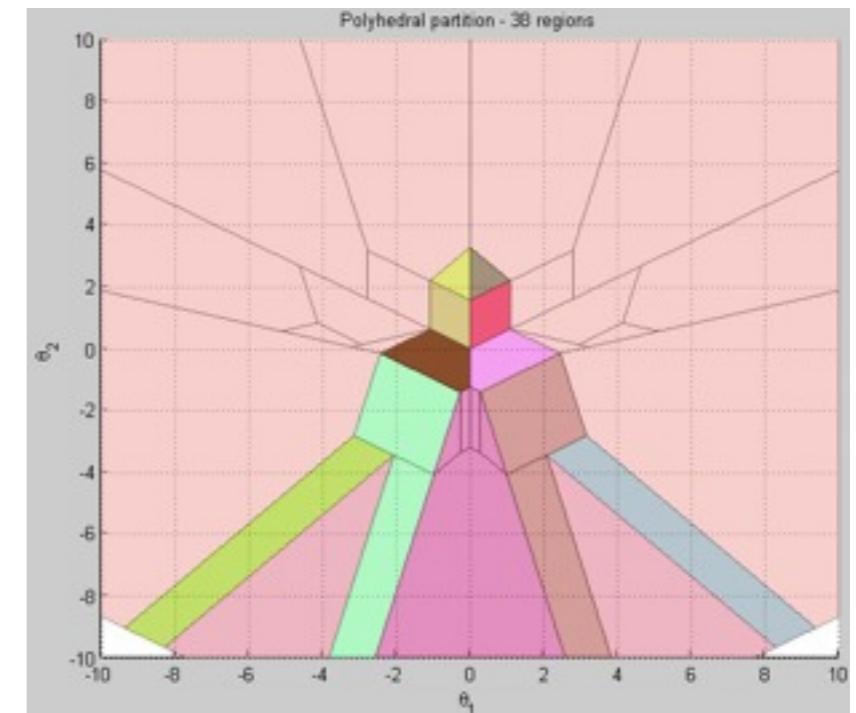
Prediction horizon  $N=1$



Prediction horizon  $N=2$



Prediction horizon  $N=3$



HybTbx: [`/demos/hybrid/bm99benchmark.m`](#)

# Explicit MPC – Temperature Control

```
>> E=expcon(C, range, options);
```

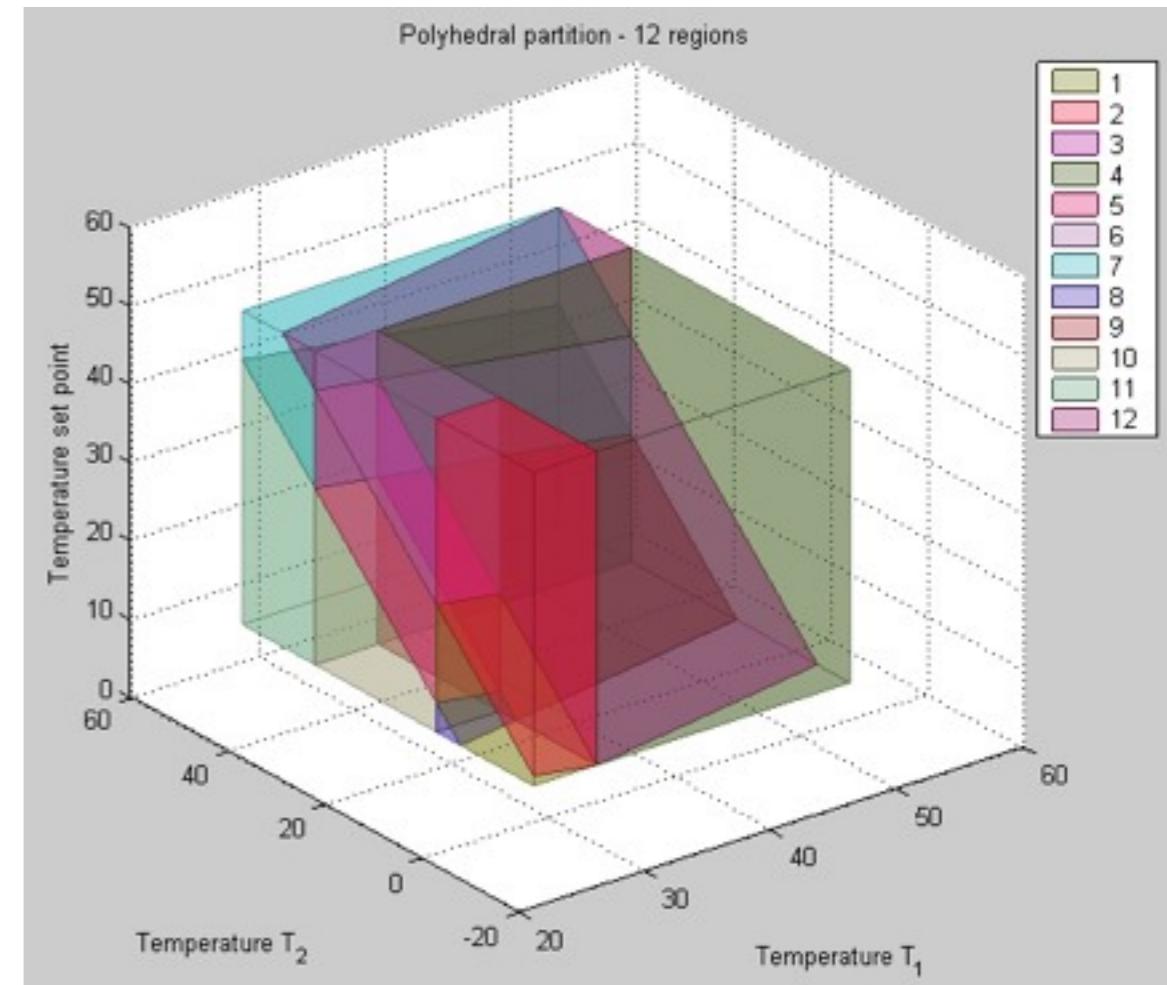
```
>> E
```

Explicit controller (based on hybrid controller C)  
3 parameter(s)  
1 input(s)  
12 partition(s)  
sampling time = 0.5

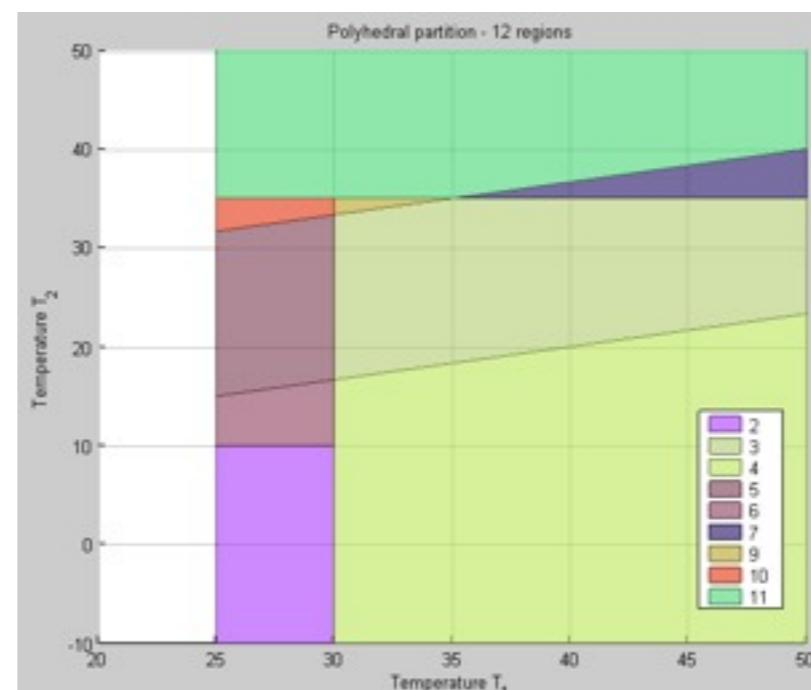
The controller is for hybrid systems (tracking)  
This is a state-feedback controller.

Type "struct(E)" for more details.

```
>>
```

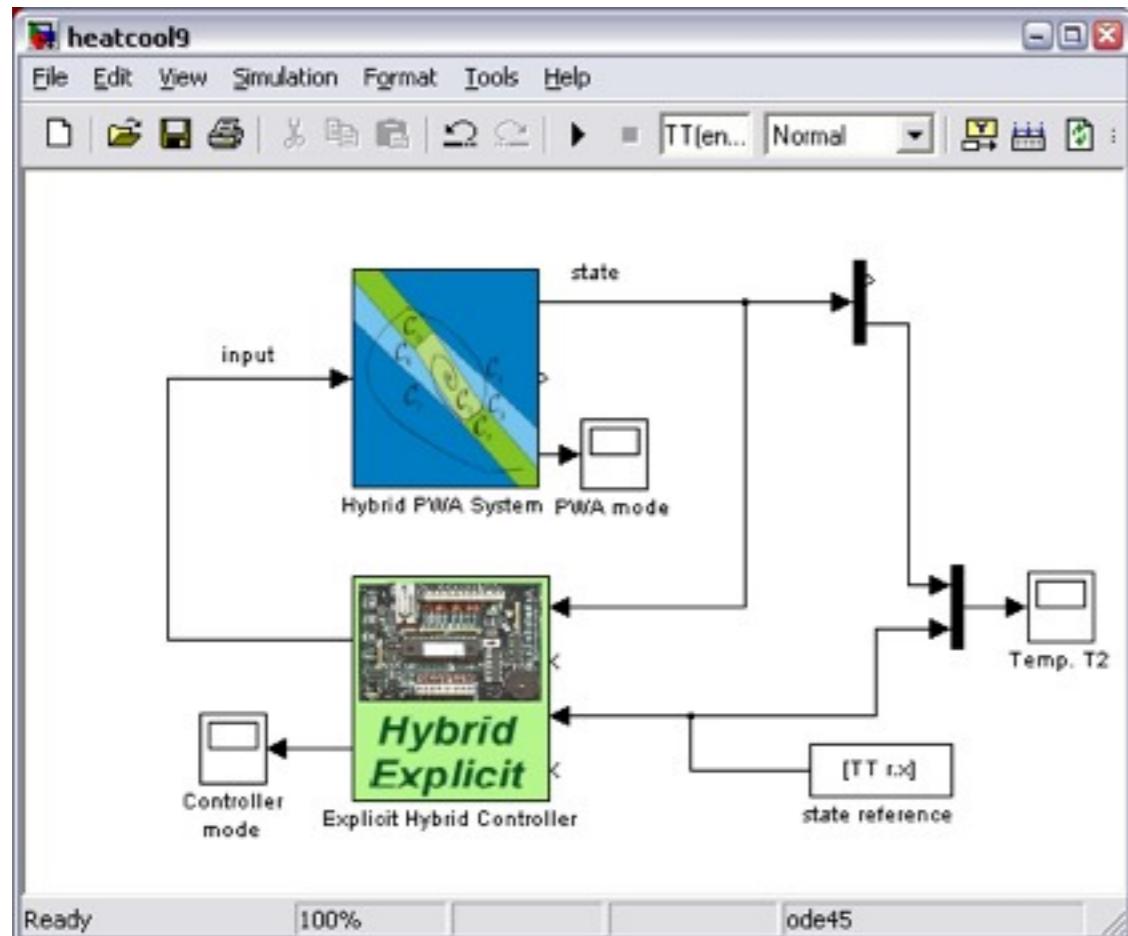


$$\begin{aligned} \min \quad & \sum_{k=1}^2 (x_2(k) - r)^2 \\ \text{s.t.} \quad & x_1(k) \geq 25 \quad k = 1, 2 \\ & \text{PWA model} \end{aligned}$$



Section in the  $(T_1, T_2)$ -space  
for  $T_{\text{ref}} = 30$

# Explicit MPC – Temperature Control



Generated  
C-code

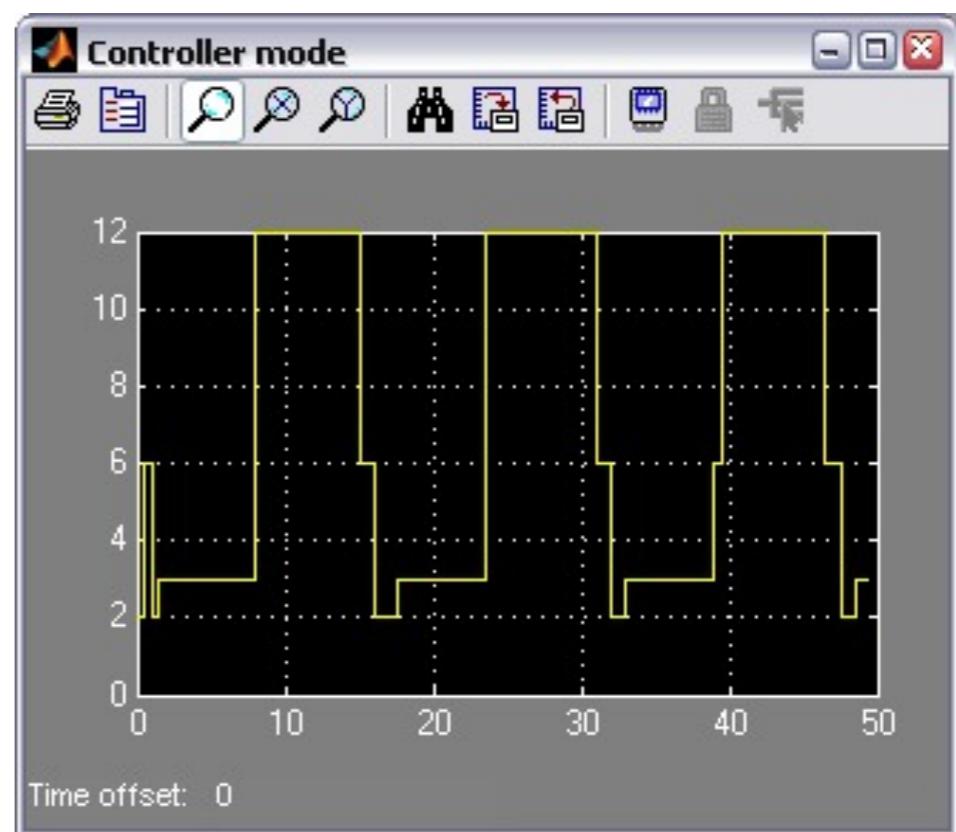
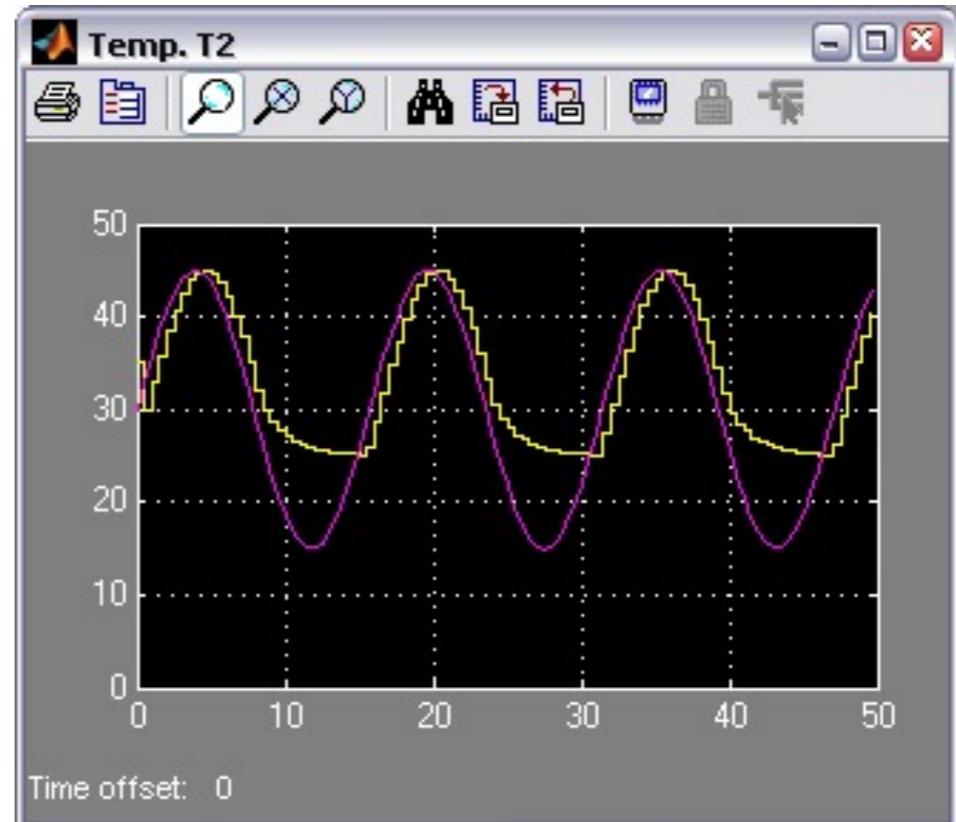


utils/expcon.h

```
#define EXPCON_REG 12
#define EXPCON_NTH 3
#define EXPCON_NYM 2
#define EXPCON_NH 72
#define EXPCON_NF 12
static double EXPCON_F[] = {
    -1,0,0,0,-1,0,
    -1,-1,-1,-1,-1,0,-3,-3,
    -3,0,-3,0,0,0,0,0,
    0,0,4,4,4,0,4,0,0,
    0,0,0,0};

static double EXPCON_G[] = {
    101.6,1.6,1.6,-1.6,98.4,0.001,0,100,51.6,
    101.6,51.6,48.4,50};

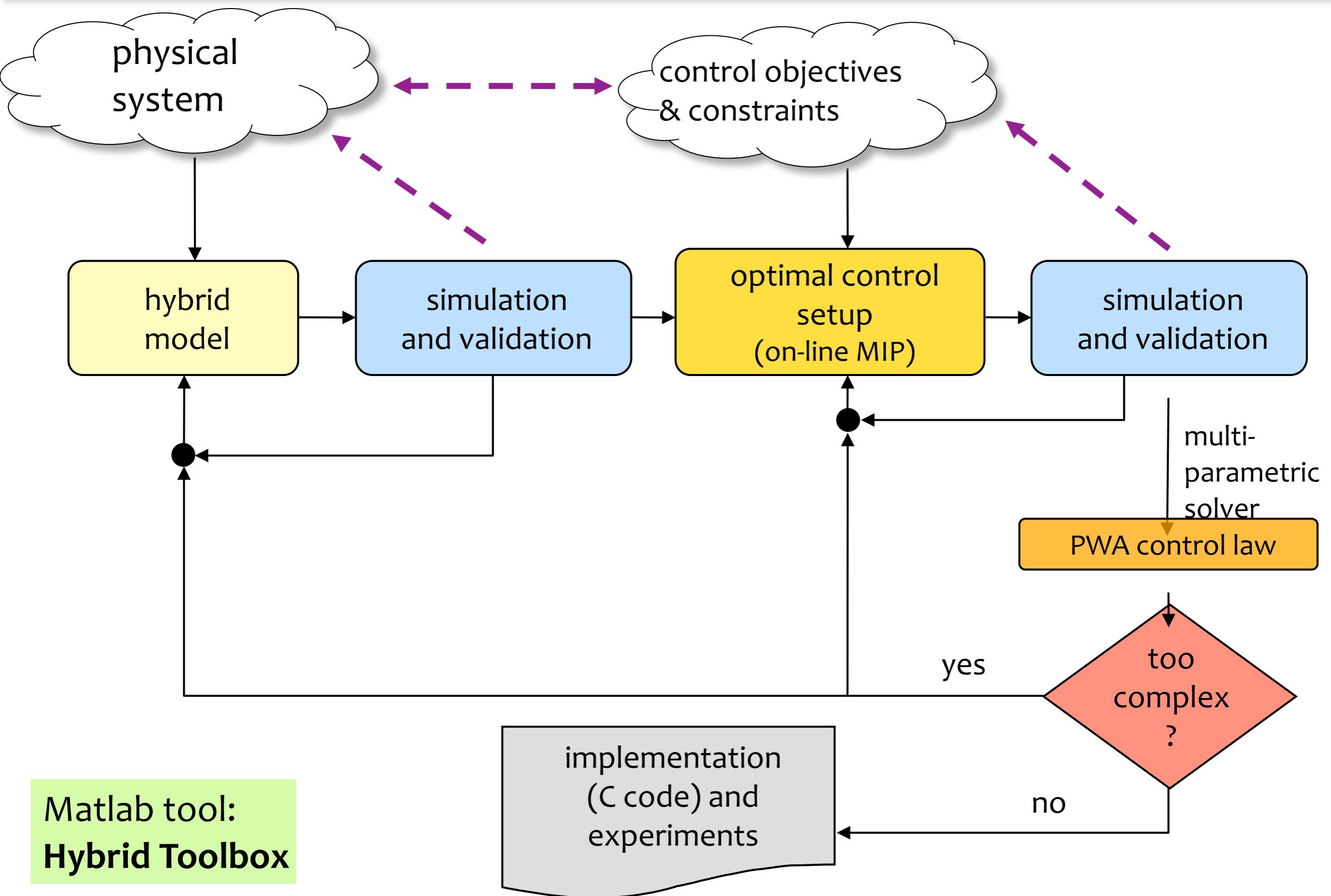
static double EXPCON_H[] = {
    0,0,0,-0.00999999,0,-0.0333333,
    0.02,0.00999999,-0.02,0,0,-0.0333333,0.02,0.00999999,
    0,0,-0.02,0.02,0,-1,0,0.00999999,0,
```



# Implementation Aspects of Hybrid MPC

- **Alternatives:** (1) **solve MIP** on-line  
(2) **evaluate a PWA function**
- **Small problems** (short horizon  $N=1,2$ , one or two inputs): explicit PWA control law preferable
  - **time** to evaluate the control law is shorter than MIP
  - **control code** is simpler (no complex solver must be included in the control software !)
  - more **insight** in controller's behavior
- **Medium/large problems** (longer horizon, many inputs and binary variables): MIP preferable

# Hybrid control design flow

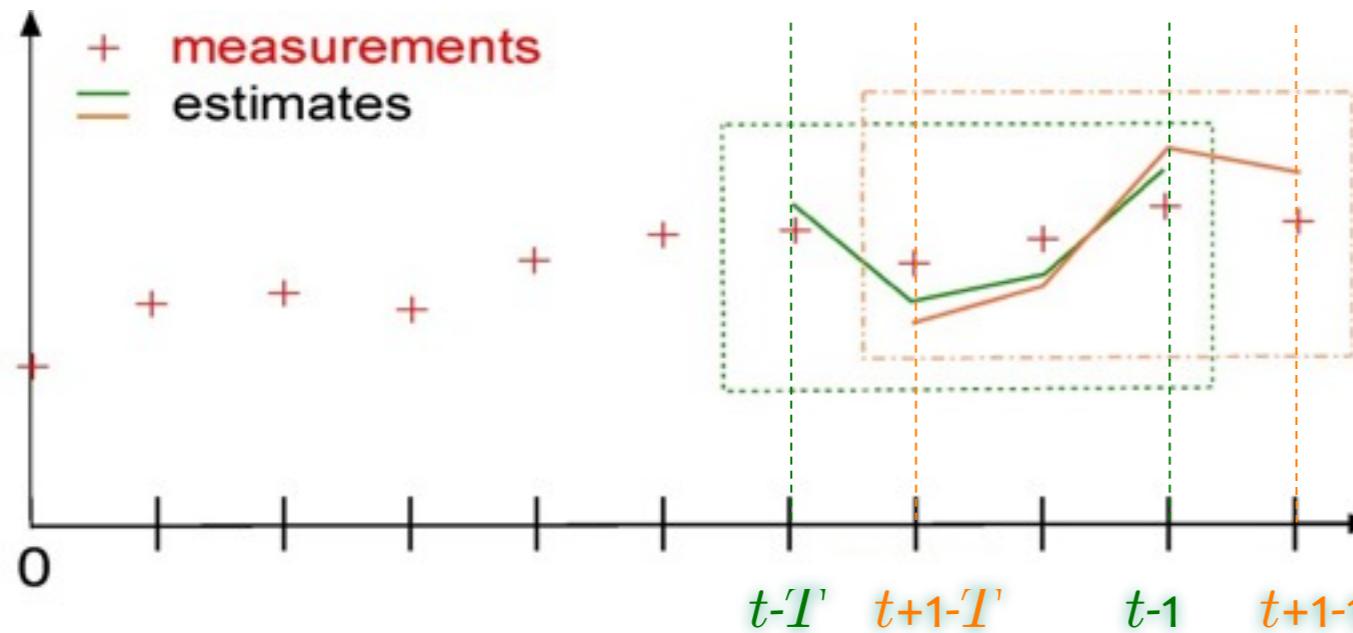


# Matlab tool: **Hybrid Toolbox**

# Moving Horizon Estimation Fault Detection & Isolation

# State Estimation / Fault Detection

- Problem: given past output measurements and inputs, estimate the current states and faults
- Solution: Use **Moving Horizon Estimation** for MLD systems (dual of MPC)



Augment the MLD model with:

- Input disturbances
- Output disturbances

$$\xi \in \mathbb{R}^n$$
$$\zeta \in \mathbb{R}^p$$

At each time  $t$   
solve the problem:

$$\min \sum_{k=1}^T \|\hat{y}(t-k|t) - y(t-k)\|^2 + \dots$$

and get estimate

$$\hat{x}(t)$$

→ MHE optimization = MIQP

(Bemporad, Mignone, Morari, ACC 1999)

→ Convergence can be guaranteed

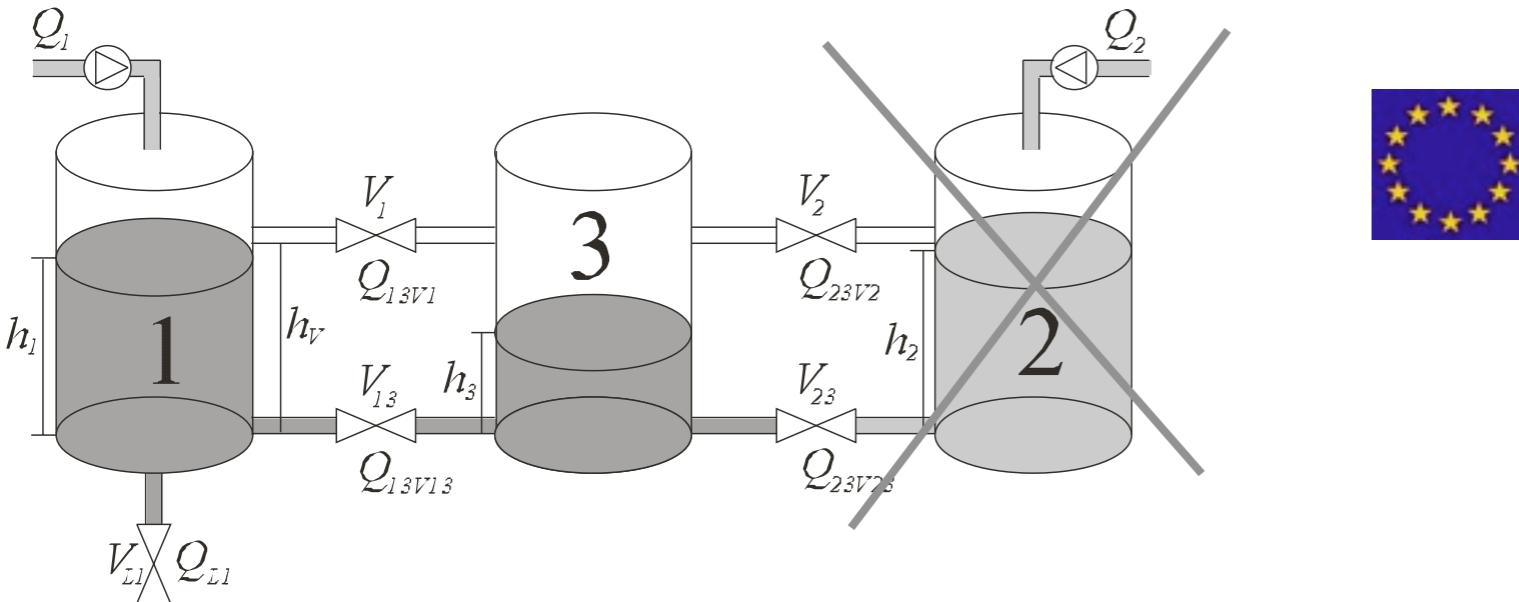
(Ferrari-T., Mignone, Morari, 2002)

**Fault detection:**

augment MLD with unknown **binary** disturbances

$$\phi \in \{0, 1\}^{n_f}$$

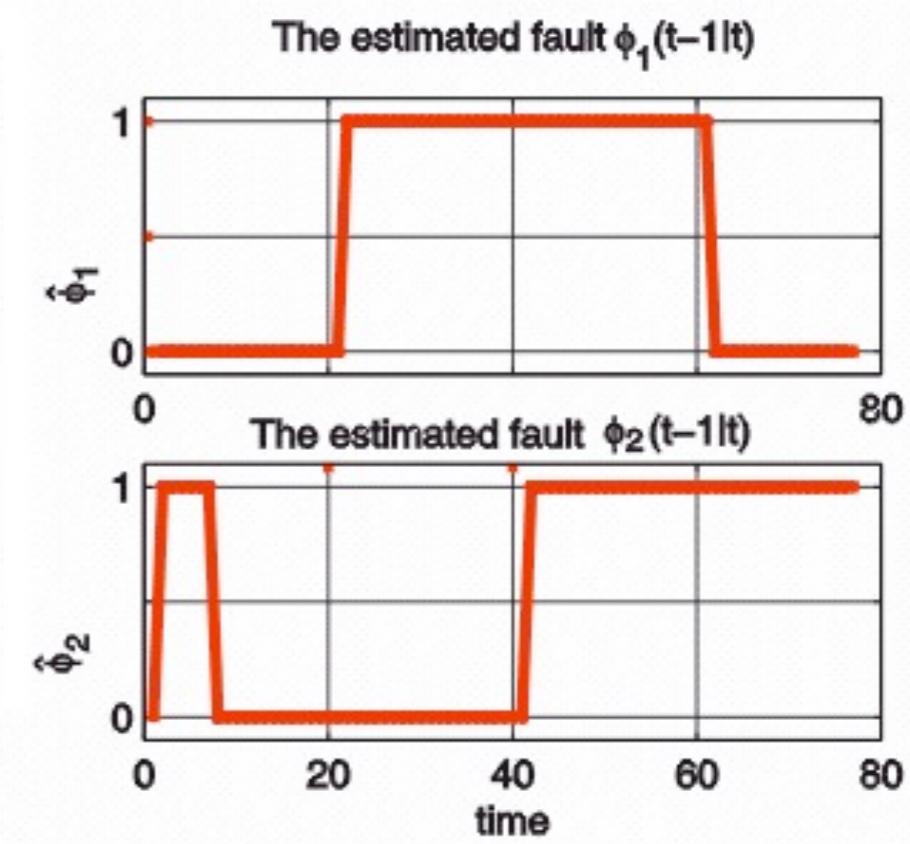
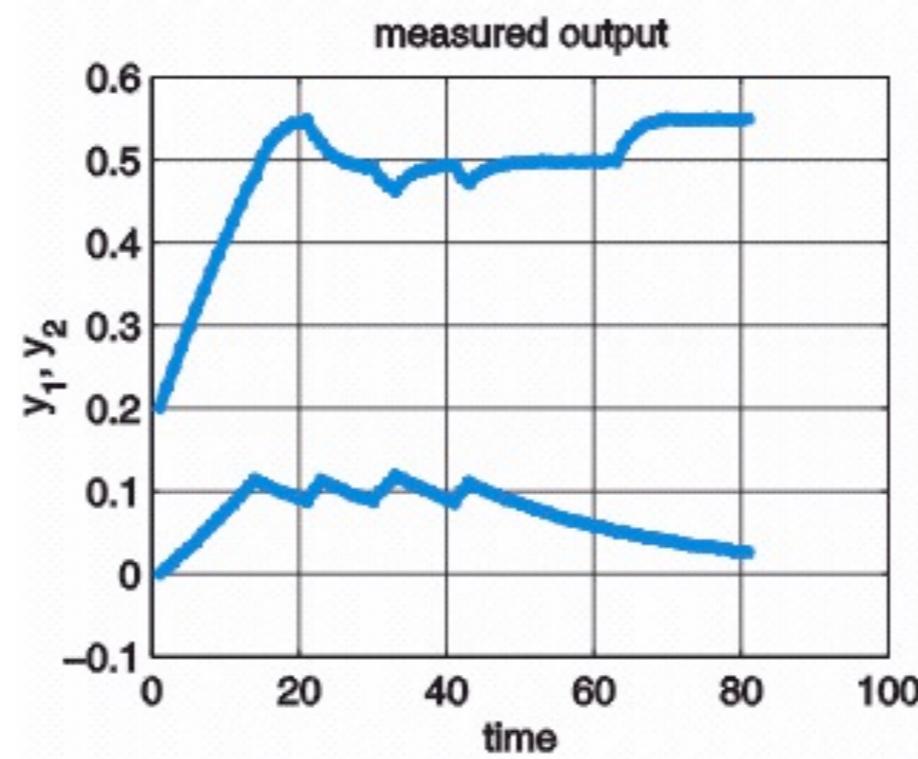
# Example: Three Tank System



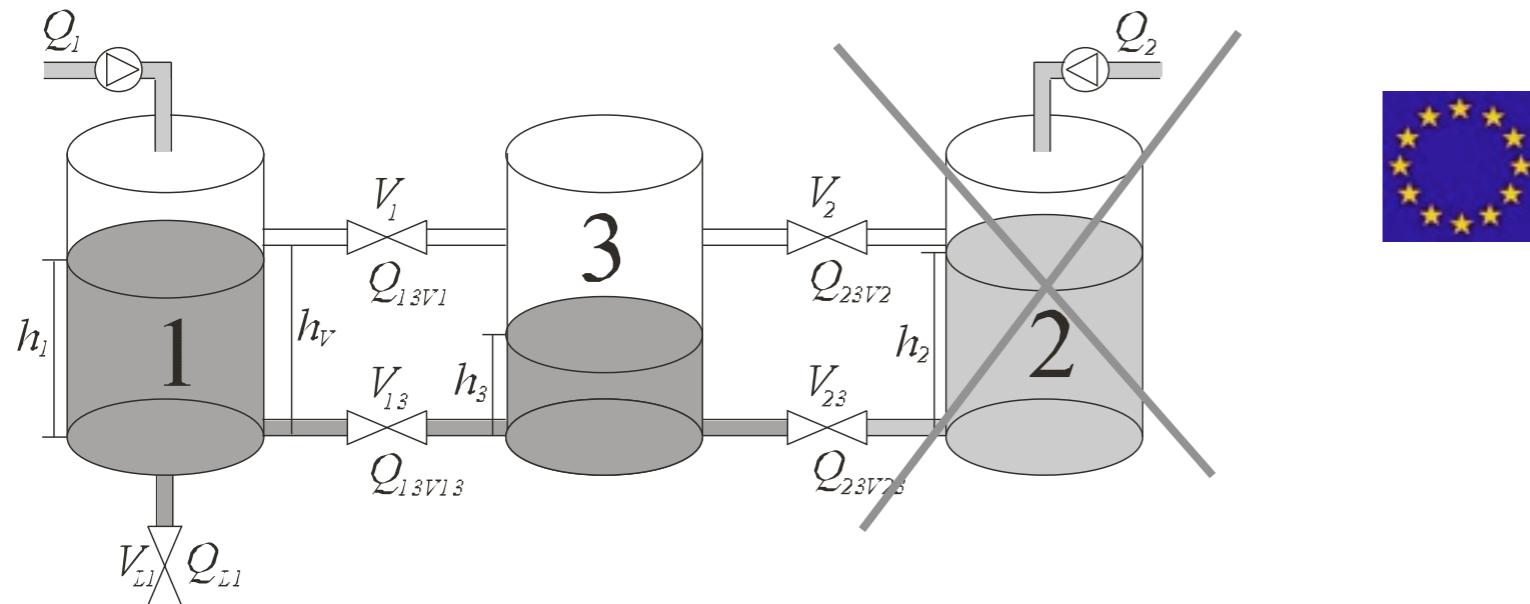
COSY Benchmark problem, ESF



- $\phi_1$ : leak in tank 1  
for  $20s \leq t \leq 60s$
- $\phi_2$ : valve  $V_1$  blocked  
for  $t \geq 40s$



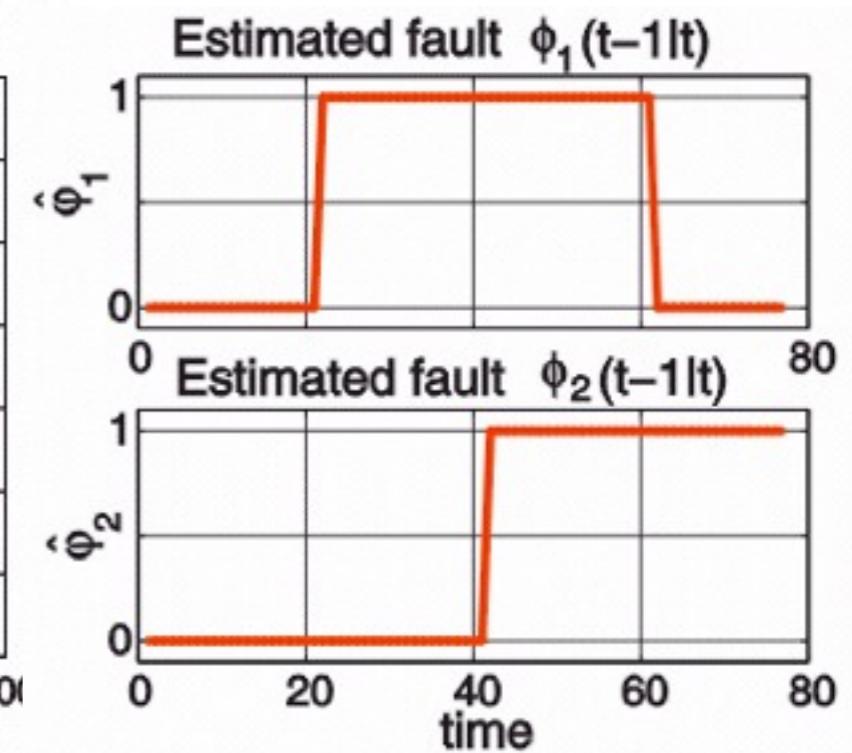
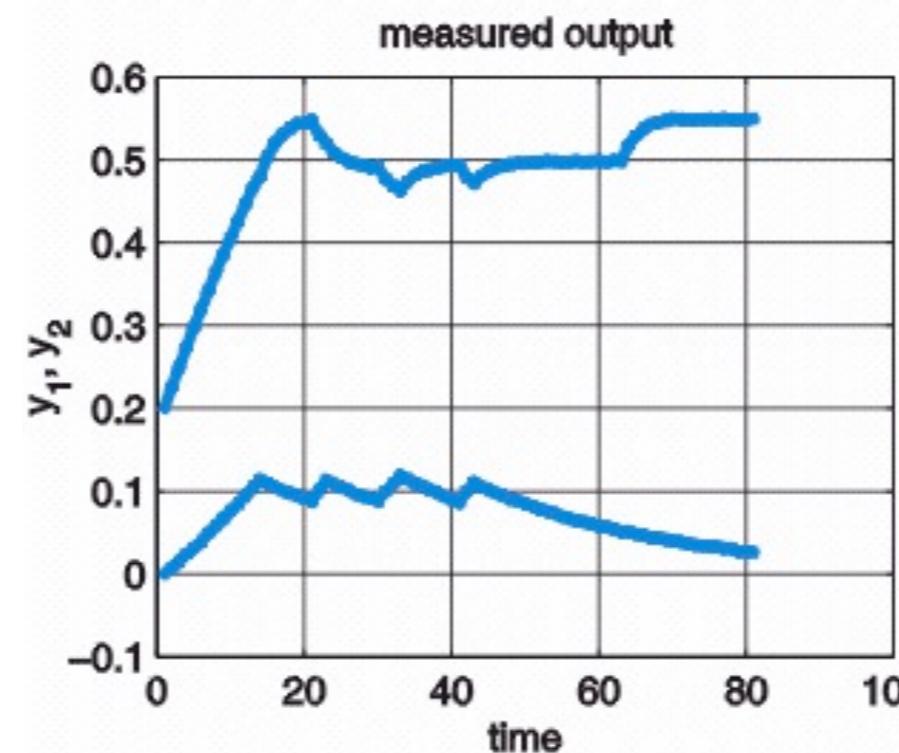
# Example: Three Tank System



COSY Benchmark problem, ESF

- $\phi_1$ : leak in tank 1  
for  $20s \leq t \leq 60s$
  - $\phi_2$ : valve  $V_1$  blocked  
for  $t > 40s$

- Add logic constraint



## A Few Hybrid MPC Tricks ...

# Measured Disturbances

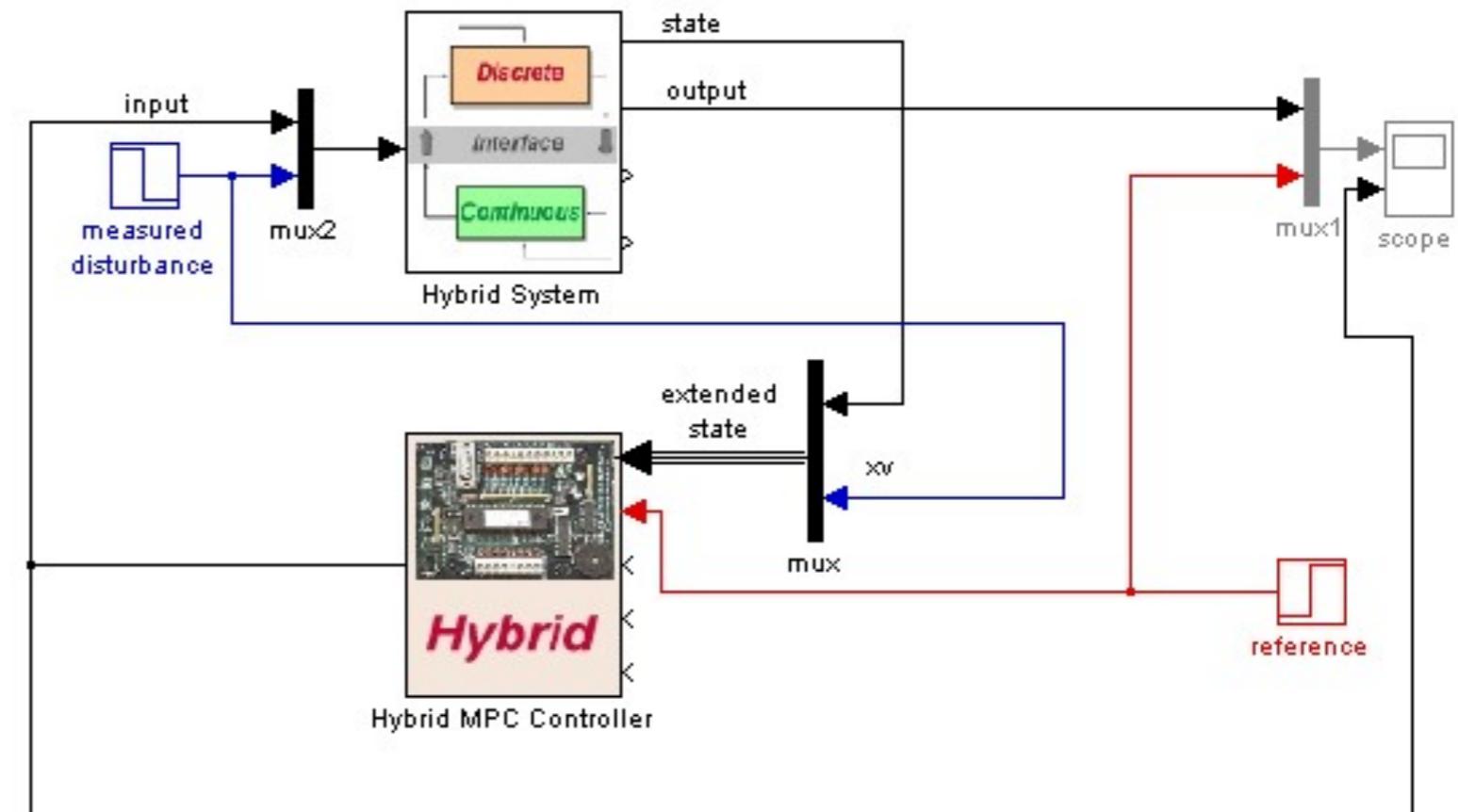
- Disturbance  $v(k)$  can be measured at time  $k$
- Augment the hybrid prediction model with a constant state

$$x_v(k+1) = x_v(k)$$

- In Hysdel:

```
INTERFACE {
  STATE {
    REAL x      [-1e3, 1e3];
    REAL xv     [-1e3, 1e3];
  }
  ...
}

IMPLEMENTATION {
  CONTINUOUS {
    x = A*x + B*u + Bv*xv
    xv= xv;
    ...
  }
}
```



/demos/hybrid/hyb\_meas\_dist.m

Note: same trick applies to linear MPC

# Hybrid MPC - Tracking

- Optimal control problem (quadratic performance index):

$$\min_{\Delta U} \quad \sum_{k=0}^{N-1} \|W^y(y(k+1) - r(t))\|^2 + \|W^{\Delta u}\Delta u(k)\|^2$$
$$[\Delta u(k) \triangleq u(k) - u(k-1)]$$

$$\text{subj. to } u_{\min} \leq u(k) \leq u_{\max}, \quad k = 0, \dots, N-1$$

$$\Delta u_{\min} \leq \Delta u(k) \leq \Delta u_{\max}, \quad k = 0, \dots, N-1$$

$$y_{\min} \leq y(k) \leq y_{\max}, \quad k = 1, \dots, N$$

- Optimization problem:  
(MIQP)

$$\min_{\Delta U} \quad J(\Delta U, x(t)) = \frac{1}{2} \Delta U' H \Delta U + [x'(t) \ r'(t) \ u'(t-1)] F \Delta U$$
$$\text{s.t. } G \Delta U \leq W + K \begin{bmatrix} x(t) \\ r(t) \\ u(t-1) \end{bmatrix}$$

Note: same trick as in linear MPC

# Integral Action in Hybrid MPC

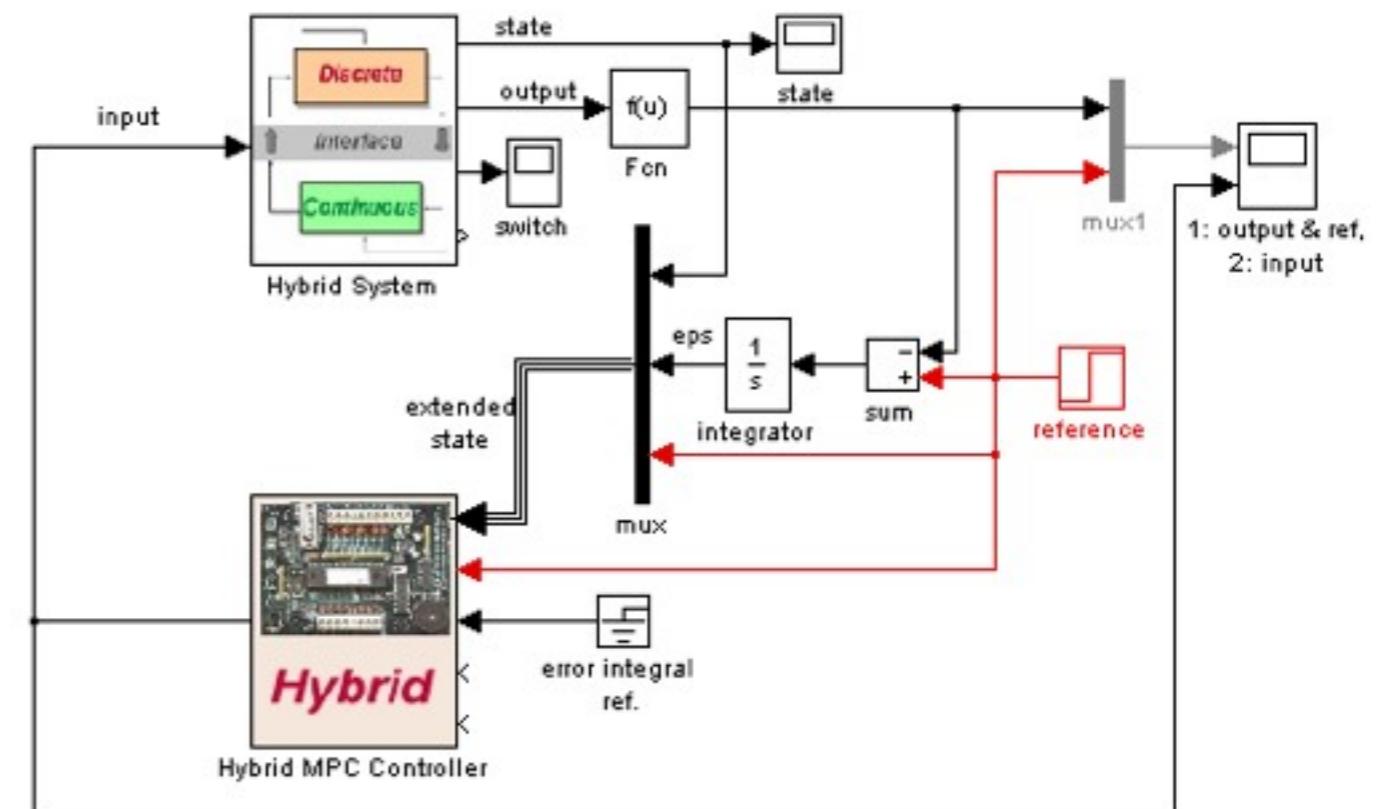
- Augment the hybrid prediction model with integrators of output errors as additional states:

$$\epsilon(k+1) = \epsilon(k) + T_s \cdot (r(k) - y(k))$$

$T_s$  = sampling time

- Treat  $r(k)$  as a measured disturbance (=additional constant state)
- Add weight on  $\epsilon(k)$  in cost function to make  $\epsilon(k) \rightarrow 0$
- In HYSDEL:

```
INTERFACE {
    STATE {
        REAL x [-100,100];
        ...
        REAL epsilon [-1e3, 1e3];
        REAL r [0, 100]; }
    OUTPUT {
        REAL y; }
    ...
}
IMPLEMENTATION {
    CONTINUOUS {
        epsilon=epsilon+Ts*(r-(c*x));
        r=r;
        ...
    }
    OUTPUT{
        y=c*x; } }
```



/demos/hybrid/hyb\_integral\_action.m

Note: same trick applies to linear MPC

# Variable Constraints

Problem: change upper (and/or lower) bounds on line

$$u(t) \leq u_{\max}(t)$$

1. Add a constant state and a new output in the prediction model:

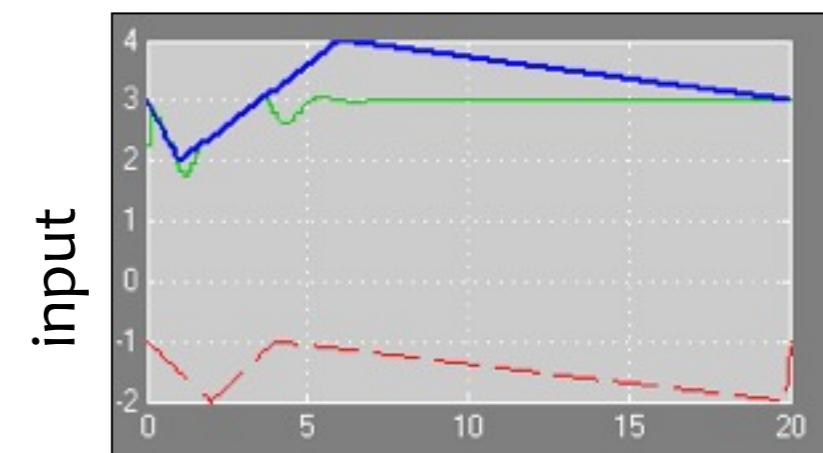
$$\begin{cases} x_u(k+1) = x_u(k) \\ y_u(k) = x_u(k) - u(k) \end{cases}$$

2. Add output constraint

$$y_u(k) \geq 0, \quad k = 0, 1, \dots, N$$

3. On-line implementation: feed the state back to the controller

$$x_u(t) = u_{\max}(t)$$



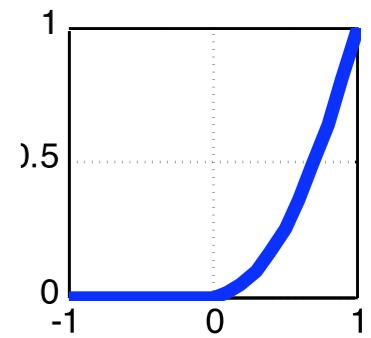
Note: same trick applies to linear MPC

`/demos/linear/varbounds.m`

# Asymmetric Weights

- Say you want to weight a variable  $u(k)$  only if  $u(k) \geq 0$
- One way is to introduce a binary variable  $[\delta=1] \leftrightarrow [u \geq 0]$ , a continuous variable  $z_u = u$  if  $\delta=1$ ,  $z_u = 0$  otherwise, and weight  $z_u$
- Better solution: avoid  $\delta$  and set:
- In Hysdel:

```
INTERFACE {
    INPUT{
        REAL u      [-100,100];
        REAL zu     [-1, 1e3];
        ...
    }
    IMPLEMENTATION{
        MUST{
            zu >= u;
            zu >= 0;
        }
    }
}
```



$$\begin{cases} \min & (\dots) + \sum z_u^2(k) \\ \text{s.t.} & \epsilon(k) \geq u(k) \\ & \epsilon(k) \geq 0 \end{cases}$$

- When  $\infty$ -norms are used, one can do the same trick:  
(better way: if the MILP problem constructor can be accessed, avoid introducing  $z_u$  and just remove the constraint  $\epsilon_u(k) \geq -[R]^i u(k)$  used to minimize  $|Ru(k)|$ , with constraint  $\epsilon_u(k) \geq 0$ )

$$\begin{cases} \min & (\dots) + |z_u(k)| \\ \text{s.t.} & z_u(k) \geq u(k) \\ & z_u(k) \geq 0 \end{cases}$$

Note: same trick applies to linear MPC

# General Remarks About MIP Modeling

The complexity of solving a mixed-integer program largely depends on the number of integer (binary) variables involved in the problem.

Henceforth, when creating a hybrid model one has to

**Be thrifty with integer variables !**

**Adding logical constraints usually helps ...**

Generally speaking:

**Modeling is an art**

(a unifying general theory does not exist)

