

Model Predictive Control of Hybrid Systems

Alberto Bemporad

Dept. Information Engineering - University of Siena, Italy

<http://www.dii.unisi.it/~bemporad>



University of Siena
(founded in 1240)



European Embedded
Control Institute



2nd HYCON PhD School on Hybrid Systems (Siena, 18/7/2007)

Outline

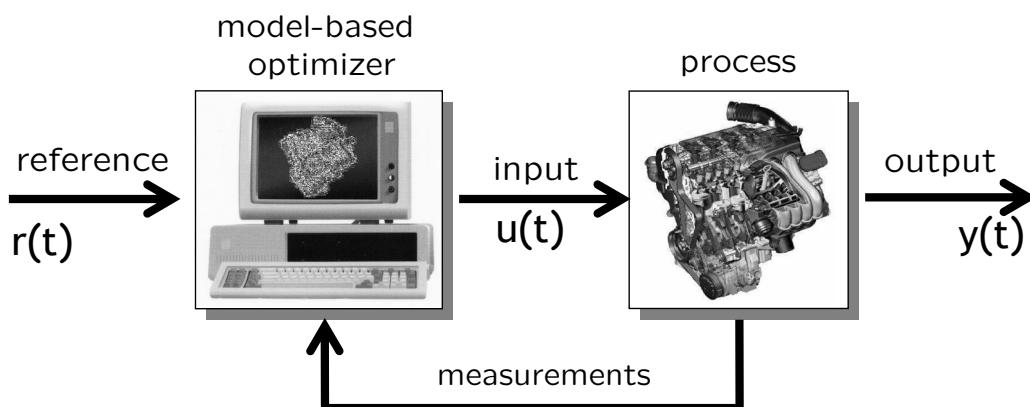
- Basics of Model Predictive Control (MPC)
- Hybrid models for MPC
- MPC of hybrid systems
- Explicit MPC (multiparametric programming)
- Optimization-based reachability analysis
- Examples

Model Predictive Control

- MPC concepts
- Linear MPC
- Matlab tools for linear MPC

3/150

Model Predictive Control



- MODEL: a model of the plant is needed to predict the future behavior of the plant
- PREDICTIVE: optimization is based on the predicted future evolution of the plant
- CONTROL: control complex constrained multivariable plants

4/150

Receding Horizon Philosophy

- At time t :

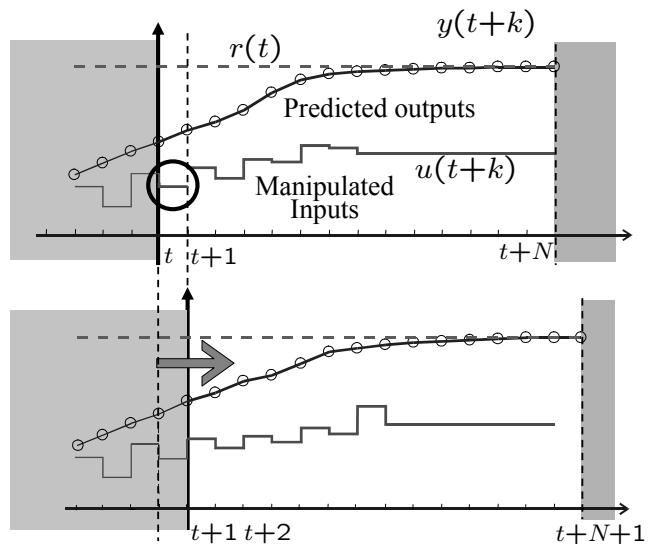
Solve an optimal control problem over a finite future horizon N

- minimize $f(|y - r|, |u|)$

- subject to constraints

$$u_{\min} \leq u \leq u_{\max}$$

$$y_{\min} \leq y \leq y_{\max}$$



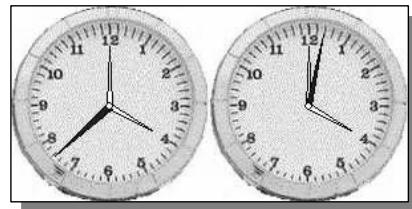
- Only apply the first optimal move $u^*(t)$
- Get new measurements, and repeat the optimization at time $t+1$

Advantage of on-line optimization: **FEEDBACK!**

5/150

Receding Horizon - Example

- MPC is like playing chess !



6/150

Receding Horizon – GPS Navigation

- **prediction model**

how vehicle moves on the map



- **constraints**

drive on roads, respect one-way roads, etc.

- **disturbances**

road works, driver's inattention, etc.

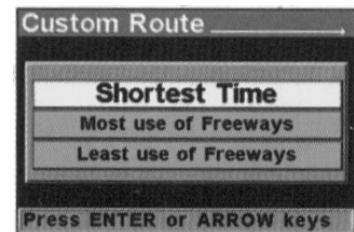
- **set point**

desired location

- **cost function:**

Ex: minimum time

Ex: penalty on highways



- **receding horizon mechanism**

event-based (optimal route re-planned when path is lost)

It's a **feedback** strategy !

7/150

Constrained Optimal Control

- Linear Model:
$$\begin{cases} x(t+1) = Ax(t) + Bu(t) \\ y(t) = Cx(t) \end{cases}$$
 $x \in \mathbb{R}^n, u \in \mathbb{R}^m$
 $y \in \mathbb{R}^p$

- Constraints:
$$\begin{cases} u_{\min} \leq u(t) \leq u_{\max} \\ y_{\min} \leq y(t) \leq y_{\max} \end{cases}$$

- Constrained optimal control problem (quadratic performance index):

$$\begin{aligned} & \min_{u(0), \dots, u(N-1)} \sum_{k=0}^{N-1} [x'(k)Qx(k) + u'(k)Ru(k)] + x'(N)Px(N) \\ \text{s.t. } & u_{\min} \leq u(k) \leq u_{\max}, \quad k = 0, \dots, N-1 \\ & y_{\min} \leq y(k) \leq y_{\max}, \quad k = 1, \dots, N \end{aligned}$$

$$Q = Q' \succeq 0, \quad R = R' \succ 0, \quad P \succeq 0$$

8/150

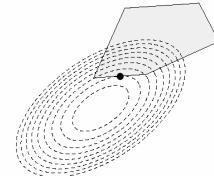
Constrained Optimal Control

- Optimization problem:

$$V(x(0)) = \frac{1}{2}x'(0)Yx(0) + \min_U \frac{1}{2}U'HU + x'(0)FU \quad (\text{quadratic})$$

$$\text{s.t. } GU \leq W + Sx(0) \quad (\text{linear})$$

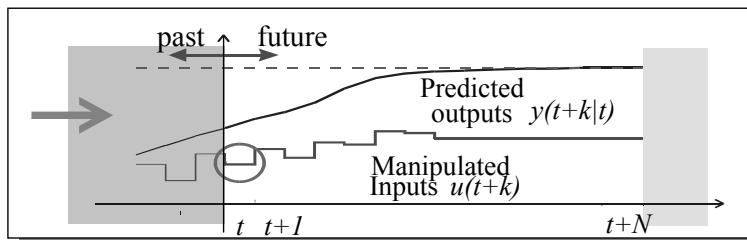
Convex QUADRATIC PROGRAM (QP)



- $U \triangleq [u'(0) \dots u'(N-1)]' \in \mathbb{R}^s$, $s \triangleq Nm$, is the optimization vector
- $H = H' \succ 0$, and H, F, Y, G, W, S depend on weights Q, R, P , upper and lower bounds $u_{\min}, u_{\max}, y_{\min}, y_{\max}$, and model matrices A, B, C

9/150

MPC of Linear Systems



At time t :

- Get/estimate the current state $x(t)$

- Solve the QP problem

$$\begin{aligned} & \min_U \frac{1}{2}U'HU + x'(t)FU \\ & \text{s.t. } GU \leq W + Sx(t) \end{aligned}$$

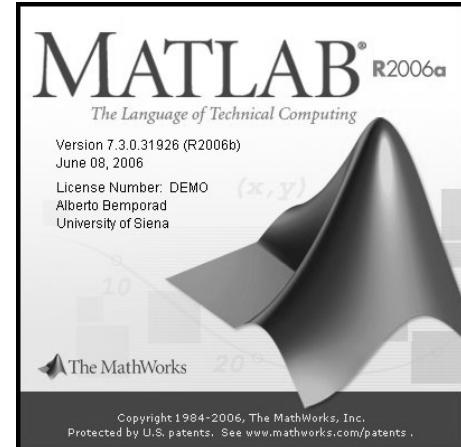
and let $U = \{u^*(0), \dots, u^*(N-1)\}$ be the solution
(=finite-horizon constrained open-loop optimal control)

- Apply only $u(t) = u^*(0)$ and discard the remaining optimal inputs
- Go to time $t+1$

10/150

Model Predictive Control Toolbox

- MPC Toolbox 2.0 (Bemporad, Ricker, Morari, 1998-2007):
 - Object-oriented implementation (MPC object)
 - MPC Simulink Library
 - MPC Graphical User Interface
 - RTW extension (code generation)
 - Linked to OPC Toolbox v2.0.1



Only linear models are handled

<http://www.mathworks.com/products/mpc/>

11/150

Example: AFTI-16

• Linearized model:

$$\left\{ \begin{array}{l} \dot{x} = \begin{bmatrix} -.0151 & -60.5651 & 0 & -32.174 \\ -.0001 & -1.3411 & .9929 & 0 \\ .00018 & 43.2541 & -.86939 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} x + \begin{bmatrix} -2.516 & -13.136 \\ -.1689 & -.2514 \\ -17.251 & -1.5766 \\ 0 & 0 \end{bmatrix} u \\ y = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} x, \end{array} \right.$$



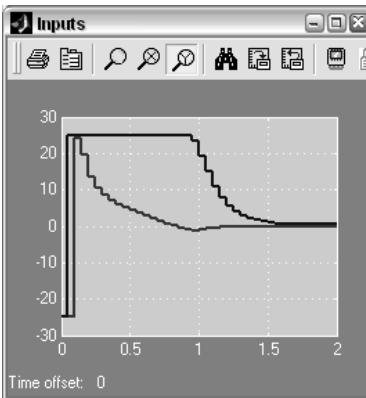
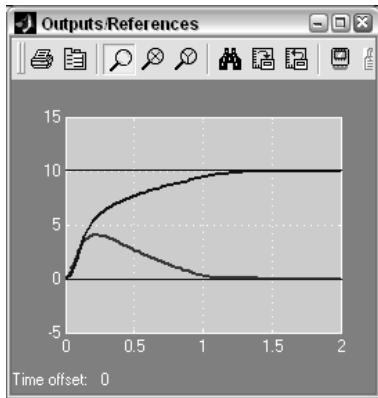
- Inputs: elevator and flaperon angle
- Outputs: attack and pitch angle
- Sampling time: $T_s = .05$ s (+ zero-order hold)
- Constraints: max 25° on both angles
- Open-loop response: unstable
(open-loop poles: $-7.6636, -0.0075 \pm 0.0556j, 5.4530$)

go to demo **afti16.m**

(MPC-Tbx)

12/150

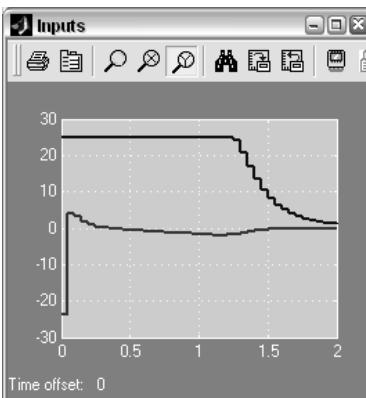
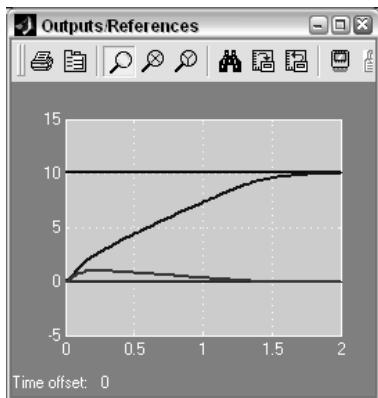
Example: AFTI-16



$$N_y = 10, N_u = 3,$$

$$w_y = \{10, 10\}, w_{\delta u} = \{.01, .01\},$$

$$u_{\min} = -25^\circ, u_{\max} = 25^\circ$$



$$N_y = 10, N_u = 3,$$

$$w_y = \{100, 10\}, w_{\delta u} = \{.01, .01\},$$

$$u_{\min} = -25^\circ, u_{\max} = 25^\circ$$

13/150

Convergence Result

Theorem 1 Consider the linear system

$$\begin{cases} x(t+1) = Ax(t) + Bu(t) \\ y(t) = Cx(t) \end{cases}$$

and the MPC control law based on

$$\begin{aligned} \min_U J(U, x(t)) &= \sum_{k=0}^{N-1} \left\{ y'(t+k|t) Q y(t+k|t) + u'(t+k) R u(t+k) \right\} \\ \text{subj. to} &\quad y_{\min} \leq y(t+k) \leq y_{\max} \\ &\quad u_{\min} \leq u(t+k) \leq u_{\max} \end{aligned}$$

Assume that the optimization problem is feasible at time $t = 0$. Then, for either $N \rightarrow \infty$ or with the extra constraint $x(t+N|t) = 0$, for all $R > 0$, $Q \geq 0$,

$$\begin{aligned} \lim_{t \rightarrow \infty} y(t) &= 0, \\ \lim_{t \rightarrow \infty} u(t) &= 0, \end{aligned}$$

while fulfilling the constraints. Moreover, provided that $(Q^{\frac{1}{2}}C, A)$ is a detectable pair, $\lim_{t \rightarrow \infty} x(t) = 0$.

(Keerthi and Gilbert, 1988)(Bemporad *et al.*, 1994)

Proof: Use value function as Lyapunov function

14/150

Convergence Proof

- Assume we set the terminal constraint $x(t + N|t) = 0$
- Let \mathcal{U}_t^* denote the optimal control sequence $\text{@} t \{u_t^*(0), \dots, u_t^*(N-1)\}$
- Let $V(t) \triangleq J(\mathcal{U}_t^*, x(t))$ = value function \rightarrow Lyapunov function
- By construction, $\mathcal{U}_1 = \{u_t^*(1), \dots, u_t^*(N-1), 0\}$ is feasible $\text{@} t+1$, and hence

$$V(t+1) = J(\mathcal{U}_{t+1}^*, x(t+1)) \leq J(\mathcal{U}_1, x(t+1)) = V(t) - y'(t)Qy(t) - u'(t)Ru(t)$$

- $V(t)$ is decreasing and lower-bounded by 0 $\Rightarrow \exists V_\infty = \lim_{t \rightarrow \infty} V(t) \Rightarrow V(t+1) - V(t) \rightarrow 0$, which implies $y'(t)Qy(t), u'(t)Ru(t) \rightarrow 0$
- Since $R > 0$, $u(t) \rightarrow 0$
- Assume for simplicity that $C = I$ (i.e., $y(t) = x(t)$) and $Q > 0$. Then also $x(t) \rightarrow 0$.

Global optimum is not needed to prove convergence !

15/150

Convergence Proof

If $Q \geq 0$ and/or $C \neq I$:

- For all $\forall k = 0, \dots, n-1$, we have

$$\lim_{t \rightarrow \infty} y'(t+k)Qy(t+k) = \lim_{t \rightarrow \infty} \|Q^{\frac{1}{2}}C(A^kx(t) + \sum_{j=0}^{k-1} A^jBu(t+k-1-j))\|^2 = 0$$

- As $u(t) \rightarrow 0$, also $Q^{\frac{1}{2}}CA^kx(t) \rightarrow 0$, and hence $\Theta x(t) \rightarrow 0$, where Θ is the observability matrix of $(Q^{\frac{1}{2}}C, A)$.
- If $(Q^{\frac{1}{2}}C, A)$ is observable, this also implies $x(t) \rightarrow 0$.
- If $(Q^{\frac{1}{2}}C, A)$ is only detectable, through a canonical decomposition one can observe that, as $u(t) \rightarrow 0$, unobservable modes go to zero spontaneously.

16/150

Convergence Proof

- Similar argument for infinite prediction horizon $N = \infty$:
 - Let \mathcal{U}_t^* denote the infinite optimal control sequence @ t $\{u_t^*(0), u_t^*(1) \dots\}$
 - Let $V(t) \triangleq J(\mathcal{U}_t^*, x(t))$ = value function (\Rightarrow Lyapunov function)
 - Because constraints were checked up to $t+k = \infty$, $\mathcal{U}_1 = \{u_t^*(1), u_t^*(2), \dots\}$ is feasible @ $t+1$ by construction.
 - Hence

$$\begin{aligned} V(t+1) &= J(\mathcal{U}_{t+1}^*, x(t+1)) \leq J(\mathcal{U}_1, x(t+1)) = \\ &= V(t) - y'(t)Qy(t) - u'(t)Ru(t) \end{aligned}$$

- Repeat same arguments as before

17/150

MPC and LQR

- Consider the MPC control law:

$$\begin{aligned} \min_U J(U, t) &= x'(t+T|t)Px(t+T|t) + \\ &\quad \sum_{k=0}^{N-1} \left\{ x'(t+k|t)Qx(t+k|t) + u'(t+k)Ru(t+k) \right\} \\ \text{subj. to } & y_{min} \leq y(t+k|t) \leq y_{max}, \quad k = 1, \dots, N \\ & u_{min} \leq u(t+k) \leq u_{max}, \quad k = 0, \dots, N-1 \\ & u(t+k) = Kx(t+k|t), \quad k = N_u, \dots, N-1 \end{aligned}$$



Jacopo Francesco
Riccati (1676 - 1754)

$R = R' > 0$, $Q = Q' \geq 0$, and P , K satisfy the Riccati equation

$$\boxed{\begin{aligned} K &= -(R + B'PB)^{-1}B'PA \\ P &= (A + BK)'P(A + BK) + K'RK + Q \end{aligned}}$$

- In a polyhedral region around the origin the MPC control law is equivalent to the constrained LQR controller with weights Q, R .
 (Chmielewski, Manousiouthakis, 1996)

(Scokaert and Rawlings, 1998)

MPC = constrained LQR

- The larger the horizon, the larger the region where MPC=LQR

18/150

Outline

- ✓ Basics of Model Predictive Control (MPC)
 - Hybrid models for MPC
 - MPC of hybrid systems
 - Explicit MPC (multiparametric programming)
 - Optimization-based reachability analysis
 - Examples

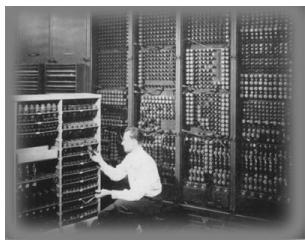
19/150

Hybrid Models for MPC

- Discrete Hybrid Automata (DHA)
- Mixed Logical Dynamical (MLD) Systems
- Piecewise Affine (PWA) Systems

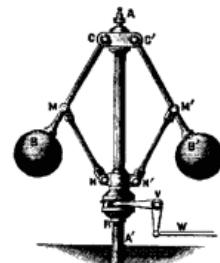
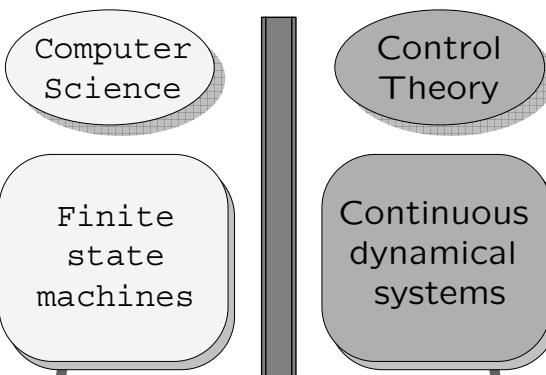
20/150

Hybrid Systems



$$x \in \{1, 2, 3, 4, 5\}$$

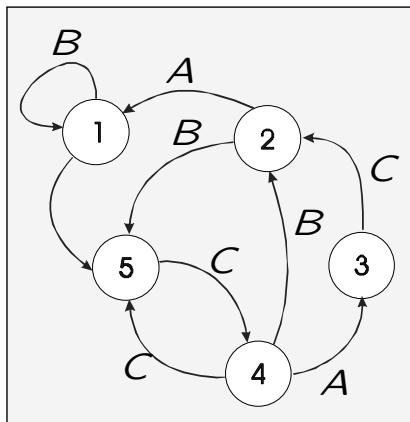
$$u \in \{A, B, C\}$$



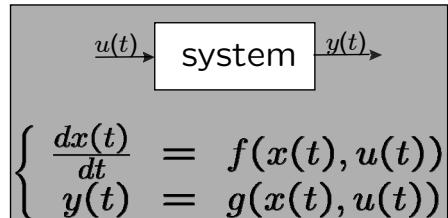
$$x \in \mathbb{R}^n$$

$$u \in \mathbb{R}^m$$

$$y \in \mathbb{R}^p$$



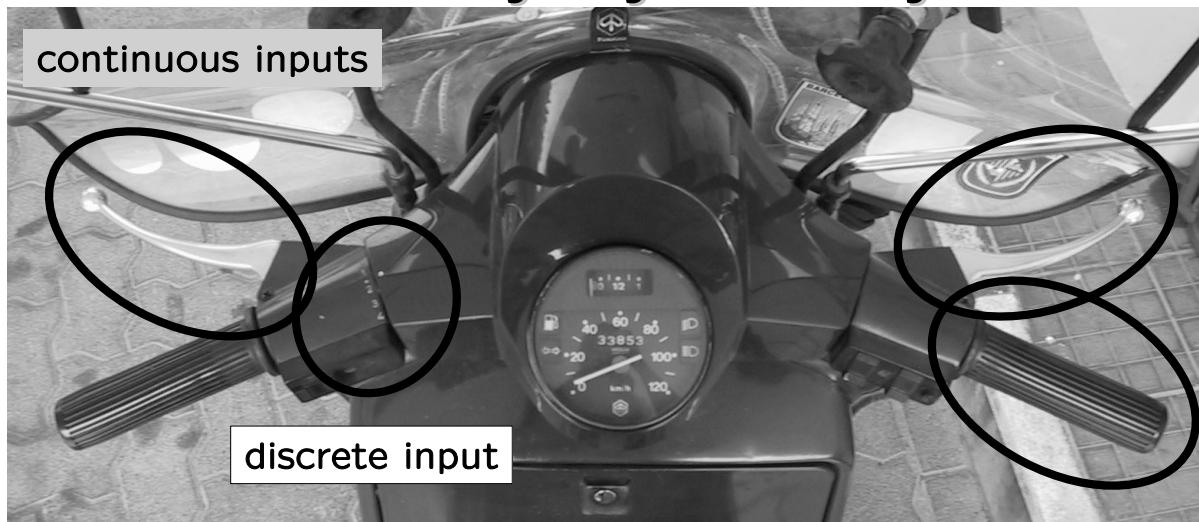
Hybrid systems



$$\begin{cases} x(k+1) = f(x(k), u(k)) \\ y(k) = g(x(k), u(k)) \end{cases}$$

21/150

"Intrinsically Hybrid" Systems



Discrete input
(gear 1,2,3,4,N)

Continuous inputs
(brakes, gas, clutch)

Continuous
dynamical states
(velocities, torques,
air-flows, fuel level)

Vespa

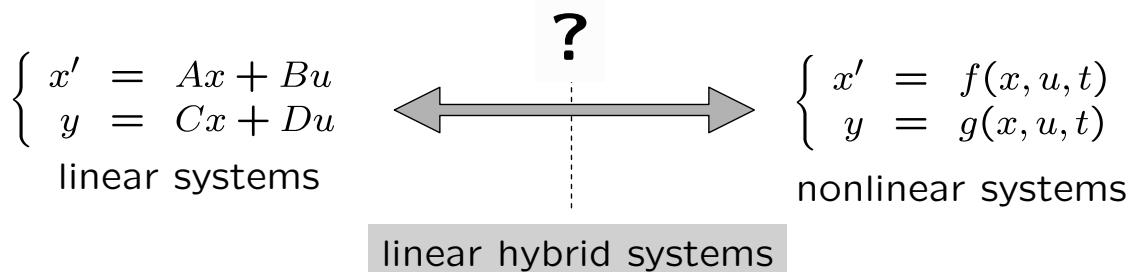


22/150

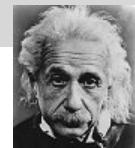
Key Requirements for Hybrid Models

- **Descriptive** enough to capture the behavior of the system
 - continuous dynamics (physical laws)
 - logic components (switches, automata, software code)
 - interconnection between logic and dynamics

- **Simple** enough for solving *analysis* and *synthesis* problems



“Make everything as simple as possible, but not simpler.”
— Albert Einstein



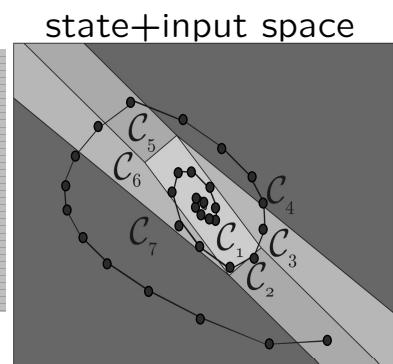
Piecewise Affine Systems

$$\begin{aligned} x(k+1) &= A_{i(k)}x(k) + B_{i(k)}u(k) + f_{i(k)} \\ y(k) &= C_{i(k)}x(k) + D_{i(k)}u(k) + g_{i(k)} \end{aligned}$$

$$i(k) \quad \text{s.t.} \quad H_{i(k)}x(k) + J_{i(k)}u(k) \leq K_{i(k)}$$

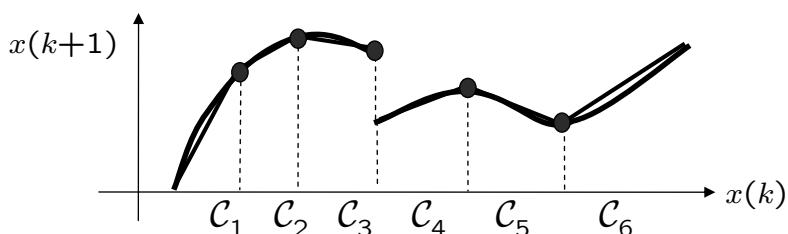
$$x \in \mathbb{R}^n, u \in \mathbb{R}^m, y \in \mathbb{R}^p$$

$$i(k) \in \{1, \dots, s\}$$



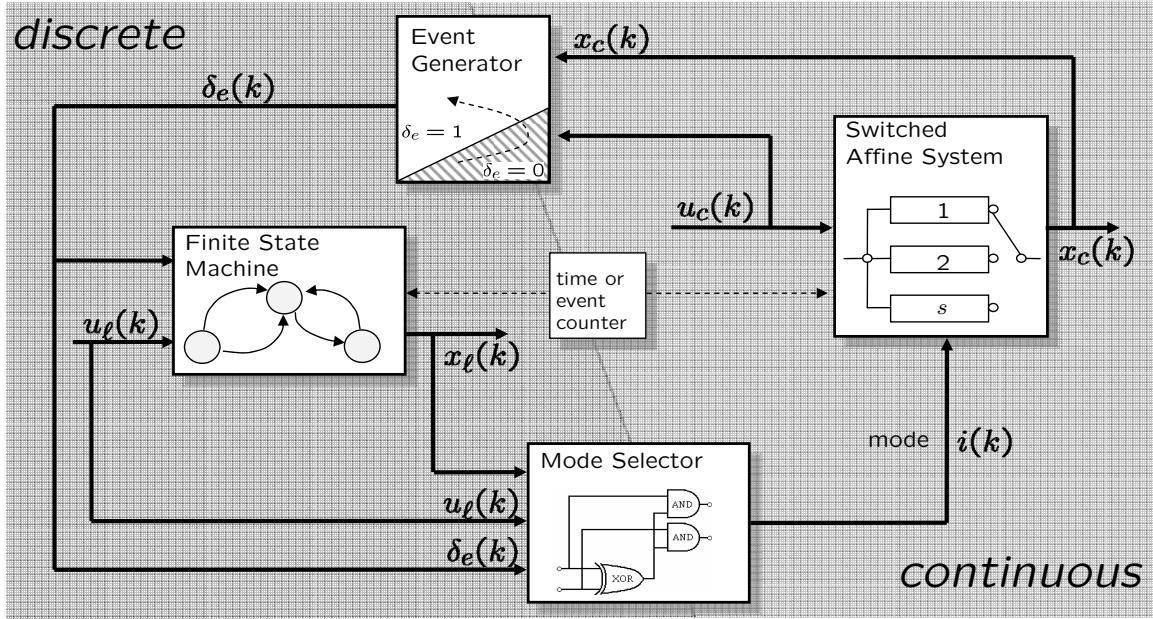
(Sontag 1981)

Can approximate nonlinear and/or discontinuous dynamics arbitrarily well



Discrete Hybrid Automaton

(Torrisi, Bemporad, 2004)



$$x_\ell \in \{0, 1\}^{n_b} = \text{binary states}$$

$$u_\ell \in \{0, 1\}^{m_b} = \text{binary inputs}$$

$$\delta_e \in \{0, 1\}^{n_e} = \text{event variables}$$

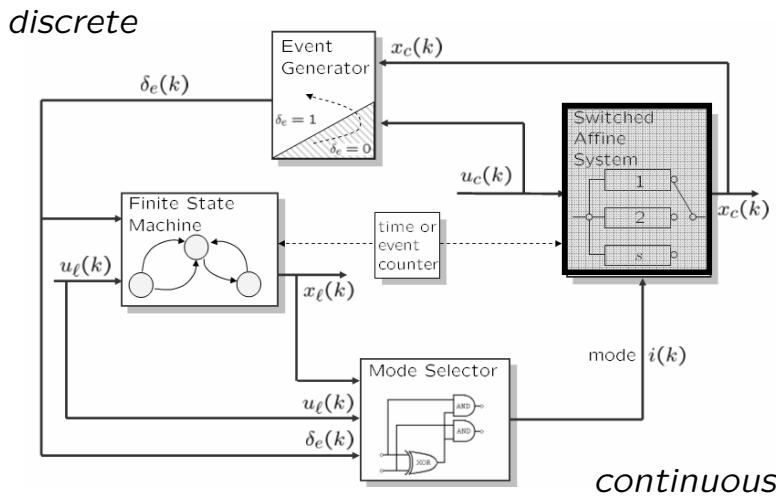
$$x_c \in \mathbb{R}^{n_c} = \text{continuous states}$$

$$u_c \in \mathbb{R}^{m_c} = \text{continuous inputs}$$

$$i \in \{1, 2, \dots, s\} = \text{current mode}$$

25/150

Switched Affine System



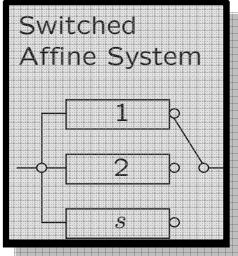
The affine dynamics depend on the current mode $i(k)$:

$$x_c(k+1) = A_{i(k)}x_c(k) + B_{i(k)}u_c(k) + f_{i(k)}$$

$$x_c \in \mathbb{R}^{n_c}, u_c \in \mathbb{R}^{m_c}$$

26/150

Switched Affine System



The state-update equation can be rewritten as a difference equation + *if-then-else* conditions:

$$z_1(t) = \begin{cases} A_1 x_c(t) + B_1 u_c(t) + f_1, & \text{if } (i(t) = 1) \\ 0, & \text{otherwise,} \end{cases}$$

$$\vdots$$

$$z_s(t) = \begin{cases} A_s x_c(t) + B_s u_c(t) + f_s, & \text{if } (i(t) = s), \\ 0, & \text{otherwise,} \end{cases}$$

$$x_c(t+1) = \sum_{i=1}^s z_i(t)$$

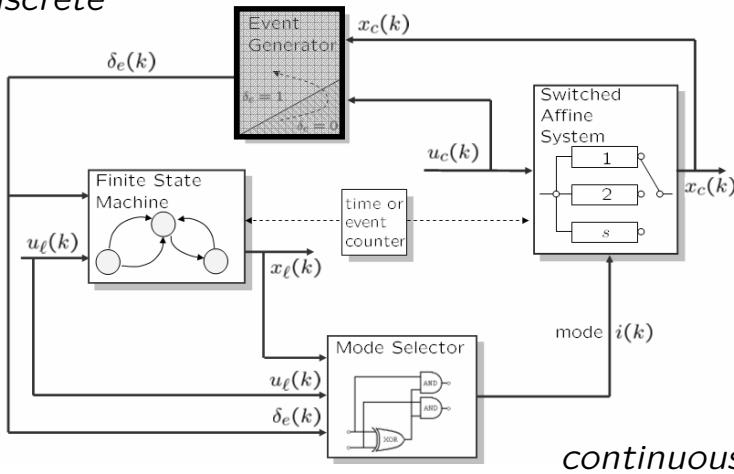
where $z_i(t) \in \mathbb{R}^{n_c}, i = 1, \dots, s$.

Output equations $y_c(t) = C_i x_c(t) + D_i u_c(t) + g_i$ admit a similar transformation.

27/150

Event Generator

discrete



continuous

Event variables are generated by linear threshold conditions over continuous states, continuous inputs, and time:

$$[\delta_e^i(k) = 1] \leftrightarrow [H^i x_c(k) + K^i u_c(k) \leq W^i]$$

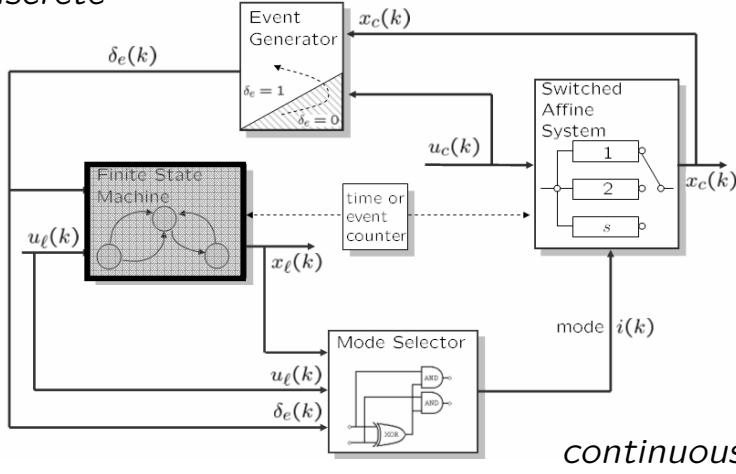
$$x_c \in \mathbb{R}^{n_c}, u_c \in \mathbb{R}^{m_c}, \delta_e \in \{0, 1\}^{n_e}$$

Example: $[\delta_e(k)=1] \leftrightarrow [x_c(k) \geq 0]$

28/150

Finite State Machine

discrete



continuous

The binary state of the finite state machine evolves according to a Boolean state update function:

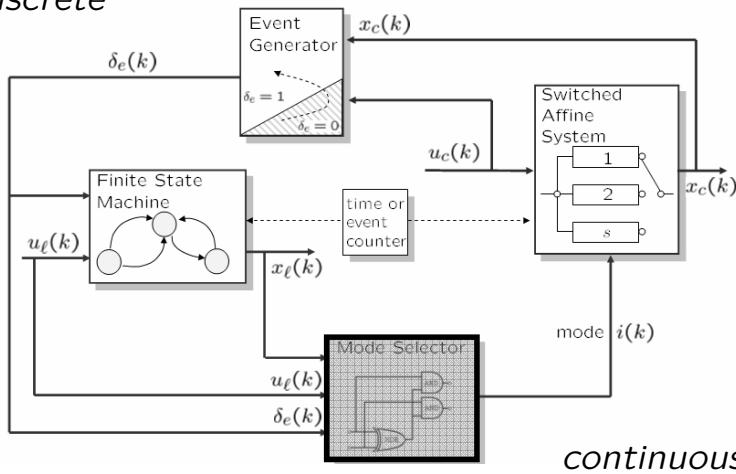
$$x_\ell(k+1) = f_B(x_\ell(k), u_\ell(k), \delta_e(k)) \quad x_\ell \in \{0, 1\}^{n_\ell}, u_\ell \in \{0, 1\}^{m_\ell}, \delta_e \in \{0, 1\}^{n_e}$$

Example: $x_\ell(k+1) = \neg \delta_e(k) \vee (x_\ell(k) \wedge u_\ell(k))$

29/150

Mode Selector

discrete



The mode selector can be seen as the output function of the discrete dynamics

continuous

The active mode $i(k)$ is selected by a Boolean function of the current binary states, binary inputs, and event variables:

$$i(k) = f_M(x_\ell(k), u_\ell(k), \delta_e(k)) \quad x_\ell \in \{0, 1\}^{n_\ell}, u_\ell \in \{0, 1\}^{m_\ell}, \delta_e \in \{0, 1\}^{n_e}$$

Example:

$$i(k) = \begin{bmatrix} \neg u_\ell(k) \vee x_\ell(k) \\ u_\ell(k) \wedge x_\ell(k) \end{bmatrix}$$

u_ℓ/x_ℓ	0	1
0	$i = \begin{bmatrix} 1 \\ 0 \end{bmatrix}$	$i = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$
1	$i = \begin{bmatrix} 0 \\ 0 \end{bmatrix}$	$i = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$

the system has 3 modes

30/150

Logic and Inequalities

$$X_1 \vee X_2 = \text{TRUE}$$

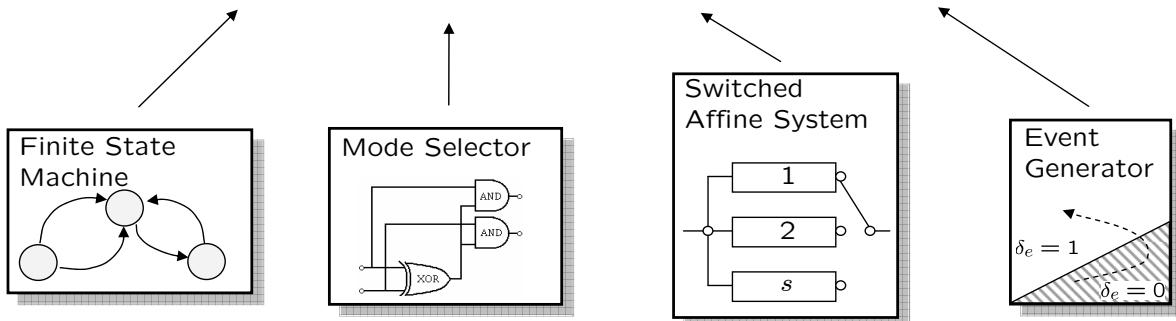
$$\longrightarrow \delta_1 + \delta_2 \geq 1, \quad \delta_1, \delta_2 \in \{0, 1\}$$

(Glover 1975,
Williams 1977,
Hooker 2000)

Any logic formula involving **Boolean variables** and **linear combinations of continuous variables** can be translated into a set of (**mixed-**)integer linear (in)equalities



All the DHA blocks can be translated into a set of (**mixed-**)integer linear equalities and inequalities



31/150

Logic and Inequalities

Any logic statement

$$f(X) = \text{TRUE}$$

$$\bigwedge_{j=1}^m (\bigvee_{i \in P_j} X_i \vee \bigvee_{i \in N_j} \neg X_i) \quad (\text{CNF})$$

$N_j, P_j \subseteq \{1, \dots, n\}$

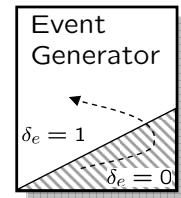
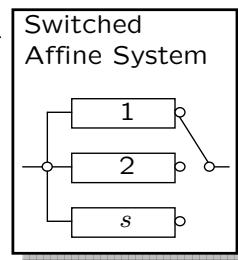
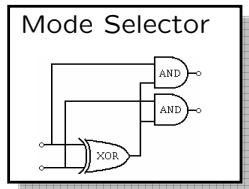
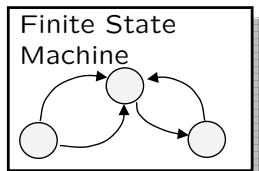
$$\left\{ \begin{array}{l} 1 \leq \sum_{i \in P_1} \delta_i + \sum_{i \in N_1} (1 - \delta_i) \\ \vdots \\ 1 \leq \sum_{i \in P_m} \delta_i + \sum_{i \in N_m} (1 - \delta_i) \end{array} \right.$$

$$[\delta_e^i(k) = 1] \leftrightarrow [H^i x_c(k) \leq W^i]$$

$$\left\{ \begin{array}{l} H^i x_c(k) - W^i \leq M^i(1 - \delta_e^i) \\ H^i x_c(k) - W^i > m^i \delta_e^i \end{array} \right.$$

$$\begin{array}{ll} \text{IF } [\delta = 1] \text{ THEN } z = a_1^T x + b_1^T u + f_1 \\ \text{ELSE } z = a_2^T x + b_2^T u + f_2 \end{array}$$

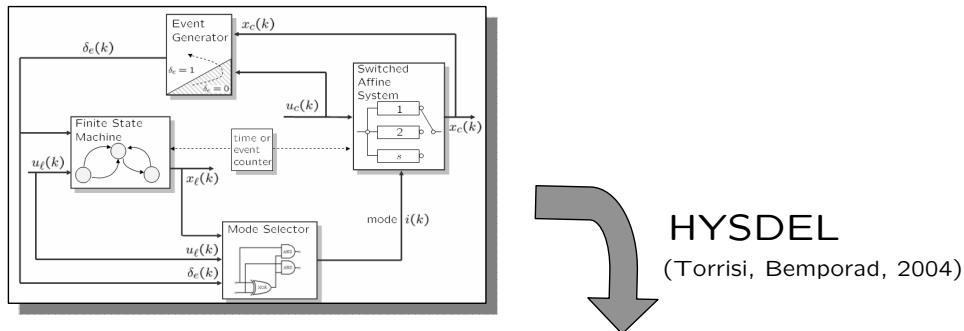
$$\left\{ \begin{array}{l} (m_2 - M_1)\delta + z \leq a_2 x + b_2 u + f_2 \\ (m_1 - M_2)\delta - z \leq -a_2 x - b_2 u - f_2 \\ (m_1 - M_2)(1 - \delta) + z \leq a_1 x + b_1 u + f_1 \\ (m_2 - M_1)(1 - \delta) - z \leq -a_1 x - b_1 u - f_1 \end{array} \right.$$



32/150

Mixed Logical Dynamical Systems

Discrete Hybrid Automaton



Mixed Logical Dynamical (MLD) Systems

(Bemporad, Morari 1999)

$$\begin{aligned}
 x(t+1) &= Ax(t) + B_1u(t) + B_2\delta(t) + B_3z(t) + B_5 \\
 y(t) &= Cx(t) + D_1u(t) + D_2\delta(t) + D_3z(t) + D_5 \\
 E_2\delta(t) + E_3z(t) &\leq E_4x(t) + E_1u(t) + E_5
 \end{aligned}$$

Continuous and binary variables

$$x \in \mathbb{R}^{nr} \times \{0,1\}^{nb}, u \in \mathbb{R}^{mr} \times \{0,1\}^{mb}$$

$$y \in \mathbb{R}^{pr} \times \{0,1\}^{pb}, \delta \in \{0,1\}^{rb}, z \in \mathbb{R}^{rr}$$

- Computationally oriented (mixed-integer programming)
- Suitable for controller synthesis, verification, ...

33/150

Mixed-Integer Models in OR

Translation of logical relations into linear inequalities is heavily used in operations research (OR) for solving complex decision problems by using mixed-integer programming (MIP)

Example: Optimal investments for quality of supply improvement in electrical energy distribution networks

(Bemporad, Muñoz, Piazzesi, 2006)



Example: Timetable generation (for demanding professors ...)

	8	9	10	11	12	13	14	15	16	17	18	19	
lun	Sistemi Operativi (18)					Misure per la Automazione (7)			Ingegneria del Software (18)				
mar	Basi di Dati (3)					Sistemi Operativi (3)			Robotica ed Automazione di Processo (18)				
mer	Robotica ed Automazione di Processo (8)			Misure per la Automazione (7)			Basi di Dati (18)			Laboratorio di Robotica e Realtà Virtuale (15)			
gio	Ingegneria del Software (18)					Sistemi Operativi (5)			Basi di Dati (3)				
ven	Laboratorio di Robotica e Realtà Virtuale (15)					Sistemi Operativi (5)			Robotica ed Automazione di Processo (8)			Misure per la Automazione (7)	
sab	Ingegneria del Software (18)					Sistemi Operativi (5)			Robotica ed Automazione di Processo (8)			Misure per la Automazione (7)	



CPU time: 0.2 s

34/150

A Simple Example

- System:

$$x(k+1) = \begin{cases} 0.8x(k) + u(k) & \text{if } x(k) \geq 0 \\ -0.8x(k) + u(k) & \text{if } x(k) < 0 \end{cases}$$

$$-10 \leq x(k) \leq 10$$

- Associate $[\delta(k) = 1] \leftrightarrow [x(k) \geq 0]$ and transform



$$\begin{aligned} x(k) &\geq m(1 - \delta(k)) & M = -m = 10 \\ x(k) &\leq -\epsilon + (M + \epsilon)\delta(k) & \epsilon > 0 \text{ "small"} \end{aligned}$$

- Then $x(k+1) = 1.6\delta(k)x(k) - 0.8x(k) + u(k)$

$$\begin{aligned} z(k) &\leq M\delta(k) & \delta(k) \in \{0, 1\} \\ z(k) = \delta(k)x(k) &\rightarrow z(k) \geq m\delta(k) \\ z(k) &\leq x(k) - m(1 - \delta(k)) \\ z(k) &\geq x(k) - M(1 - \delta(k)) \end{aligned}$$

- Rewrite as a linear equation



$$x(k+1) = 1.6z(k) - 0.8x(k) + u(k)$$

(Note: nonlinear system $x(k+1) = 0.8|x(k)| + u(k)$)

35/150

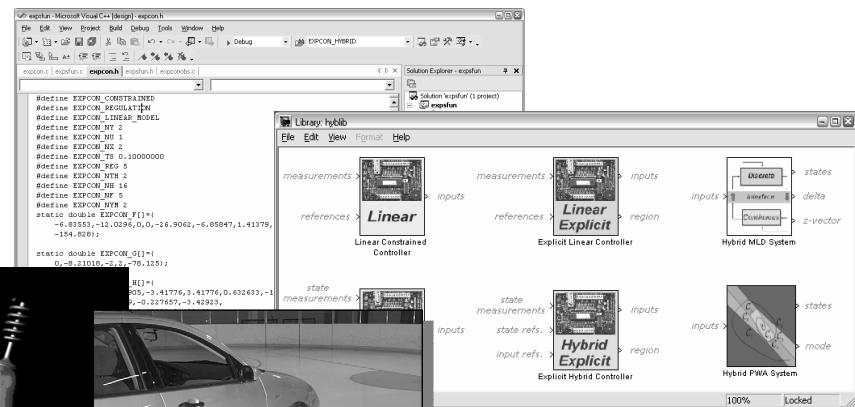
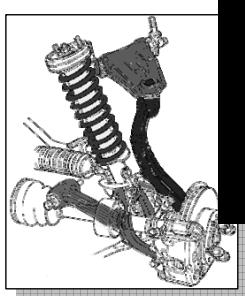
Hybrid Toolbox for Matlab

Features:

(Bemporad, 2003-today)

- Hybrid model (MLD and PWA) design, simulation, verification
- Control design for linear systems w/ constraints and hybrid systems (on-line optimization via QP/MILP/MIQP)
- Explicit control (via multiparametric programming)
- C-code generation
- Simulink

Support:



>1300 downloads in 2 years

<http://www.dii.unisi.it/hybrid/toolbox>

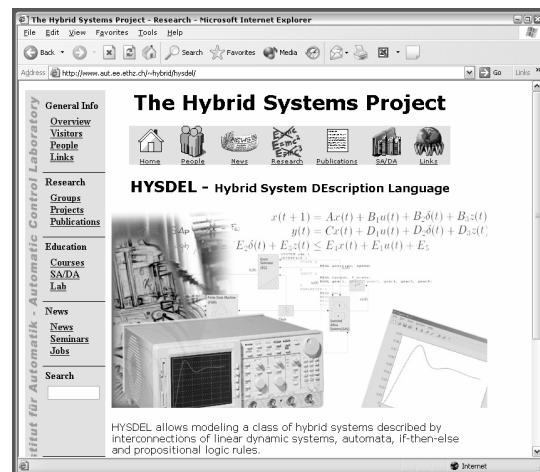
6/150

HYSDEL

(HYbrid Systems DEscription Language)

- Describe *hybrid systems*:

- Automata
- Logic
- Lin. Dynamics
- Interfaces
- Constraints



(Torrisi, Bemporad, 2004)

- Automatically generate MLD models in Matlab

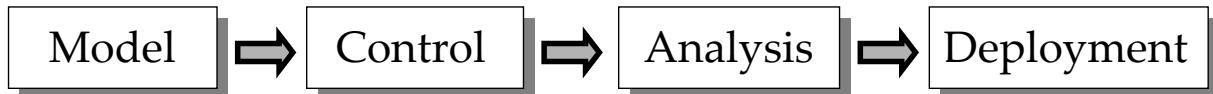
Download: <http://www.dii.unisi.it/hybrid/toolbox>

Reference: <http://control.ethz.ch/~hybrid/hysdel>

37/150

The Multi-Parametric Toolbox (MPT)

(Kvasnica, Grieder, Baotic)



- MPT is a repository of hybrid systems design tools from international experts utilizing state-of-the-art optimization packages
- Free, open-source, GNU licensed toolbox
- Special emphasis:
 - Low controller complexity
 - Numerical robustness (e.g. handling of degenerate problems)
 - Generation of real-time executable code

MPT, HYSDEL
Hybrid ID Tbx
YALMIP, CDD
Ellipsoidal Tbx
Proj. algorithms
SeDuMi

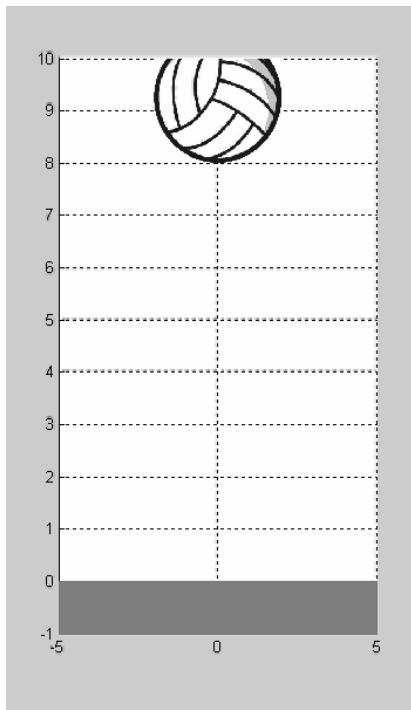
ETH
Eidgenössische Technische Hochschule Zürich
Swiss Federal Institute of Technology Zurich



<http://control.ee.ethz.ch/~mpt/>

38/150

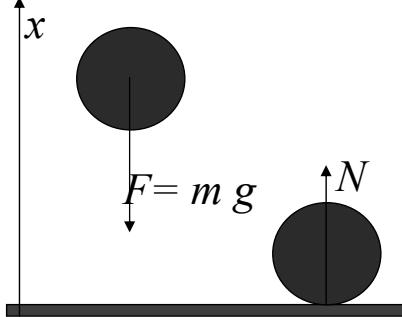
Example: Bouncing Ball



$$\ddot{x} = -g$$

$$x \leq 0 \Rightarrow \dot{x}(t^+) = -(1 - \alpha)\dot{x}(t^-)$$

$$\alpha \in [0, 1]$$



How to model the bouncing ball
as a discrete-time hybrid system ?

39/150

Bouncing Ball – Time Discretization

$$\frac{x(k) > 0 :}{\begin{array}{c} x \\ \downarrow -mg \end{array}} \quad v(k) \approx \frac{x(k) - x(k-1)}{T_s} \quad -g = \dot{x}(k) \approx \frac{v(k) - v(k-1)}{T_s}$$

$$\rightarrow \left\{ \begin{array}{l} v(k+1) = v(k) - T_s g \\ x(k+1) = x(k) + T_s v(k+1) \\ = x(k) + T_s v(k) - T_s^2 g \end{array} \right.$$

$$\frac{x(k) \leq 0 :}{\begin{array}{c} x \\ \downarrow \end{array}} \quad v(k) = -(1 - \alpha)v(k-1) \quad \begin{array}{l} x(k+1) = x(k-1) \\ = x(k) - T_s v(k) \end{array}$$

$$\rightarrow \left\{ \begin{array}{l} v(k+1) = -(1 - \alpha)v(k) \\ x(k+1) = x(k) - T_s v(k) \end{array} \right.$$

40/150

HYSDEL - Bouncing Ball

```

SYSTEM bouncing_ball {
INTERFACE {
/* Description of variables and constants */
    STATE { REAL height [-10,10];
              REAL velocity [-100,100]; }

PARAMETER {
    REAL g;
    REAL alpha; /* 0=elastic, 1=完全ly anelastic */
    REAL Ts; }

IMPLEMENTATION {
AUX { REAL z1;
      REAL z2;
      BOOL negative; }

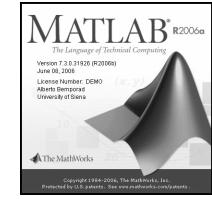
AD { negative = height <= 0; }

DA { z1 = { IF negative THEN height-Ts*velocity
            ELSE height+Ts*velocity-Ts*Tg};
     z2 = { IF negative THEN -(1-alpha)*velocity
            ELSE velocity-Ts*g}; }

CONTINUOUS {
height = z1;
velocity=z2; }
}
}

```

go to demo /demos/hybrid/bball.m



41/150

Bouncing Ball

```

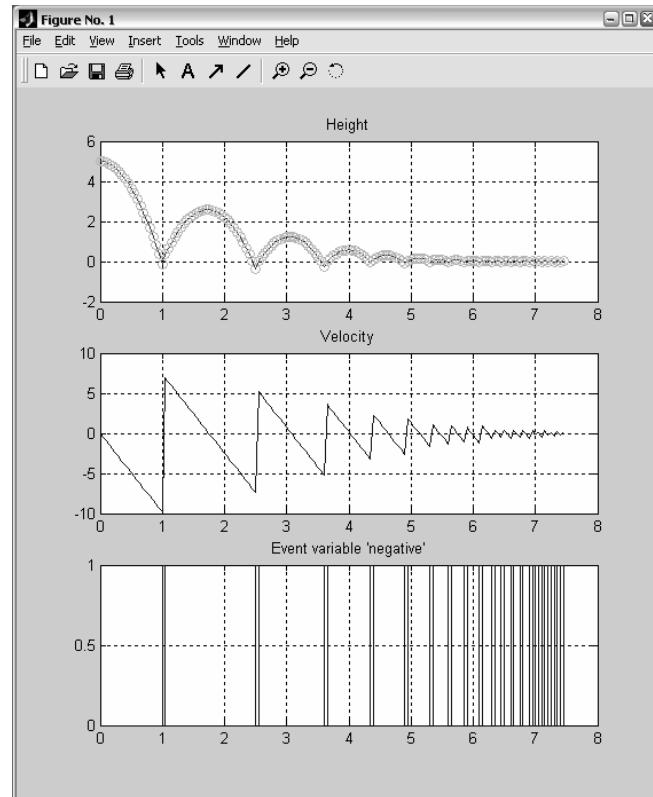
>>Ts=0.05;
>>g=9.8;
>>alpha=0.3;

>>S=mld('bouncing_ball',Ts);

>>N=150;
>>U=zeros(N,0);
>>x0=[5 0]';

>>[X,T,D]=sim(S,x0,U);

```

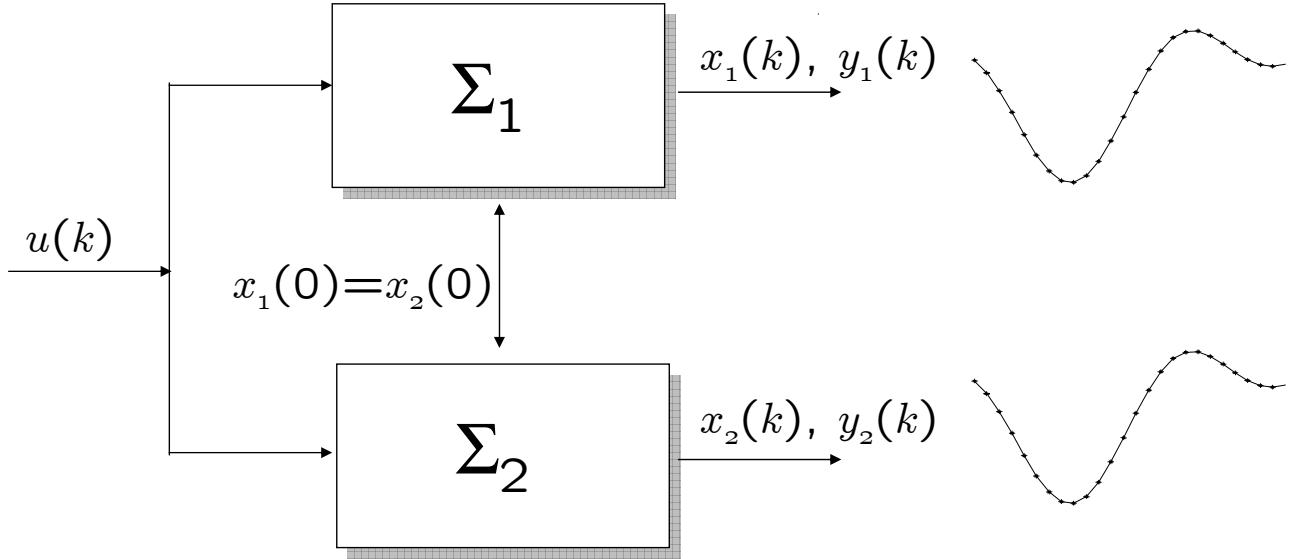


Note: no Zeno effects
in discrete time !

42/150

Equivalences of Hybrid Models

Definition 1 Two hybrid systems Σ_1, Σ_2 are equivalent if for all initial conditions $x_1(0) = x_2(0)$ and input $\{u_1(k)\}_{k \in \mathbb{Z}_+} = \{u_2(k)\}_{k \in \mathbb{Z}_+}$ then $x_1(k) = x_2(k)$ and $y_1(k) = y_2(k)$, for all $k \in \mathbb{Z}_+$.



43/150

MLD and PWA Systems

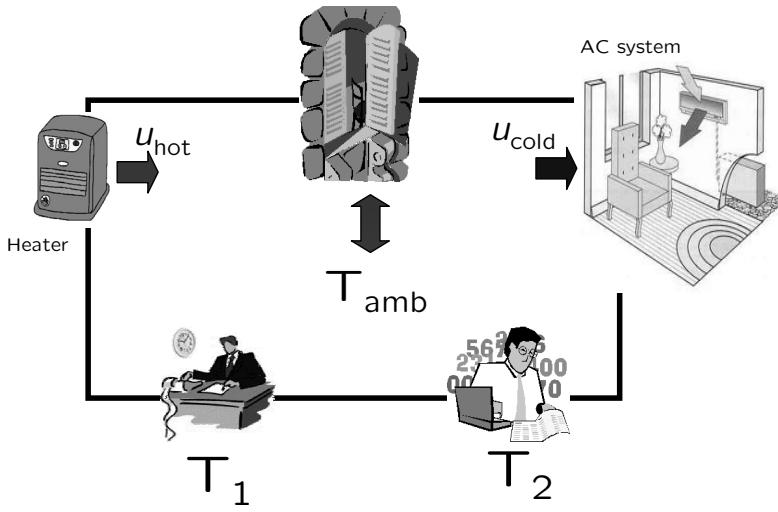
Theorem MLD systems and PWA systems are equivalent

(Bemporad, Ferrari-Trecate, Morari, IEEE TAC, 2000)

- Proof is constructive: given an MLD system it returns its equivalent PWA form
- Drawback: it needs the enumeration of all possible combinations of binary states, binary inputs, and δ variables
- Most of such combinations lead to empty regions
- Efficient algorithms are available for converting MLD models into PWA models avoiding such an enumeration:
 - A. Bemporad, "Efficient Algorithms for Converting Mixed Logical Dynamical Systems into an Equivalent Piecewise Affine Form", IEEE Trans. Autom. Contr., 2004.
 - T. Geyer, F.D. Torrisi and M. Morari, "Efficient Mode Enumeration of Compositional Hybrid Models", HSCC'03

44/150

Example: Room Temperature



Hybrid dynamics

- #1 turns the heater (air conditioning) on whenever he is cold (hot)
- If #2 is cold he turns the heater on, unless #1 is hot
- If #2 is hot he turns the air conditioning on, unless #2 is cold
- Otherwise, heater and air conditioning are off

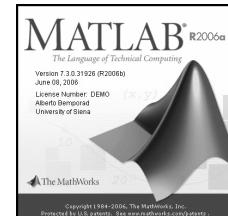
- $\dot{T}_1 = -\alpha_1(T_1 - T_{\text{amb}}) + k_1(u_{\text{hot}} - u_{\text{cold}})$ (body temperature dynamics of #1)
- $\dot{T}_2 = -\alpha_2(T_2 - T_{\text{amb}}) + k_2(u_{\text{hot}} - u_{\text{cold}})$ (body temperature dynamics of #2)

go to demo /demos/hybrid/heatcool.m

45/150

HYSDEL Model

```
SYSTEM heatcool {
  INTERFACE {
    STATE { REAL T1 [-10,50];
            REAL T2 [-10,50];
    }
    INPUT { REAL Tamb [-10,50];
    }
    PARAMETER {
      REAL Ts, alphai, alpha2, k1, k2;
      REAL Thot1, Tcold1, Thot2, Tcold2, Uc, Uh;
    }
  }
  IMPLEMENTATION {
    AUX { REAL uhot, ucold;
          BOOL hot1, hot2, cold1, cold2;
    }
    AD { hot1 = T1>=Thot1;
         hot2 = T2>=Thot2;
         cold1 = T1<=Tcold1;
         cold2 = T2<=Tcold2;
    }
    DA { uhot = (IF cold1 | (cold2 & ~hot1) THEN Uh ELSE 0);
         ucold = (IF hot1 | (hot2 & ~cold1) THEN Uc ELSE 0);
    }
    CONTINUOUS { T1 = T1+Ts*(-alphai*(T1-Tamb)+k1*(uhot-ucold));
                 T2 = T2+Ts*(-alpha2*(T2-Tamb)+k2*(uhot-ucold));
    }
  }
}
```



Hybrid Toolbox
for Matlab

<http://www.dii.unisi.it/hybrid/toolbox>

>>S=mld('heatcoolmodel',Ts)

get the MLD model in Matlab

>> [XX,TT]=sim(S,x0,U);

simulate the MLD model

46/150

Hybrid MLD Model

- MLD model

$$\begin{aligned} x(t+1) &= Ax(t) + B_1u(t) + B_2\delta(t) + B_3z(t) \\ y(t) &= Cx(t) + D_1u(t) + D_2\delta(t) + D_3z(t) \\ E_2\delta(t) + E_3z(t) &\leq E_1u(t) + E_4x(t) + E_5 \end{aligned}$$

- 2 continuous states: (temperatures T_1, T_2)
- 1 continuous input: (room temperature T_{amb})
- 2 auxiliary continuous vars: (power flows $u_{\text{hot}}, u_{\text{cold}}$)
- 6 auxiliary binary vars: (4 thresholds + 2 for OR condition)
- 20 mixed-integer inequalities

Possible combination of integer variables: $2^6 = 64$

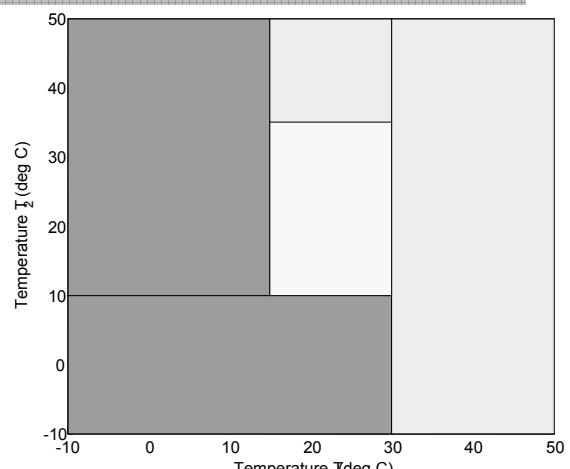
47/150

Hybrid PWA Model

- PWA model

$$\begin{aligned} x(k+1) &= A_{i(k)}x(k) + B_{i(k)}u(k) + f_{i(k)} \\ y(k) &= C_{i(k)}x(k) + D_{i(k)}u(k) + g_{i(k)} \\ i(k) \text{ s.t. } H_{i(k)}x(k) + J_{i(k)}u(k) &\leq K_{i(k)} \end{aligned}$$

- 2 continuous states: (temperatures T_1, T_2)
- 1 continuous input: (room temperature T_{amb})
- 5 polyhedral regions
(partition does not depend on input)



>> P=pwa(S);

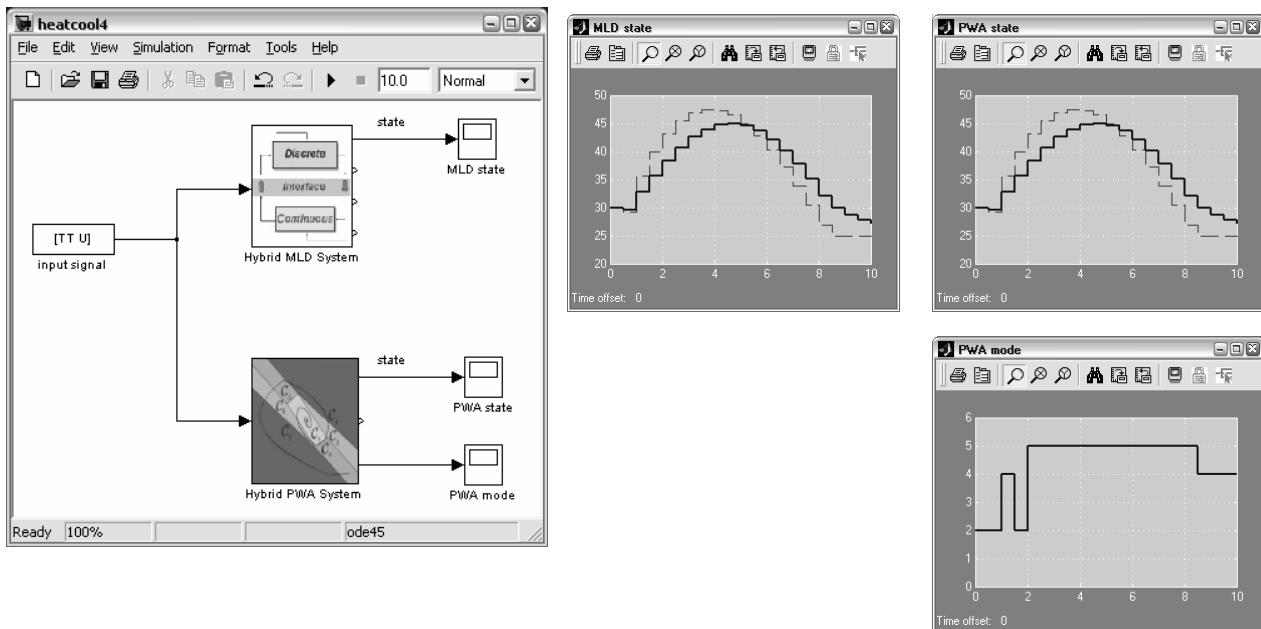
$$\begin{cases} u_{\text{hot}} = 0 \\ u_{\text{cold}} = 0 \end{cases}$$

$$\begin{cases} u_{\text{hot}} = 0 \\ u_{\text{cold}} = \bar{U}_C \end{cases}$$

$$\begin{cases} u_{\text{hot}} = \bar{U}_H \\ u_{\text{cold}} = 0 \end{cases}$$

48/150

Simulation in Simulink



MLD and PWA models are equivalent

49/150

Other Existing Hybrid Models

- Linear complementarity (LC) systems (Heemels, 1999)

$$\begin{aligned} x(t+1) &= Ax(t) + B_1u(t) + B_2w(t) \\ y(t) &= Cx(t) + D_1u(t) + D_2w(t) \\ v(t) &= E_1x(t) + E_2u(t) + E_3w(t) + e_4 \\ 0 \leq v(t) \perp w(t) &\geq 0 \end{aligned}$$

Ex: mechanical systems
circuits with diodes etc.

- Extended linear complementarity (ELC) systems

Generalization of LC systems

(De Schutter, De Moor, 2000)

- Min-max-plus-scaling (MMPS) systems

(De Schutter, Van den Boom, 2000)

$$\begin{aligned} x(t+1) &= M_x(x(t), u(t), w(t)) \\ y(t) &= M_x(x(t), u(t), w(t)) \\ 0 \geq M_c(x(t), u(t), w(t)) & \geq 0 \end{aligned}$$

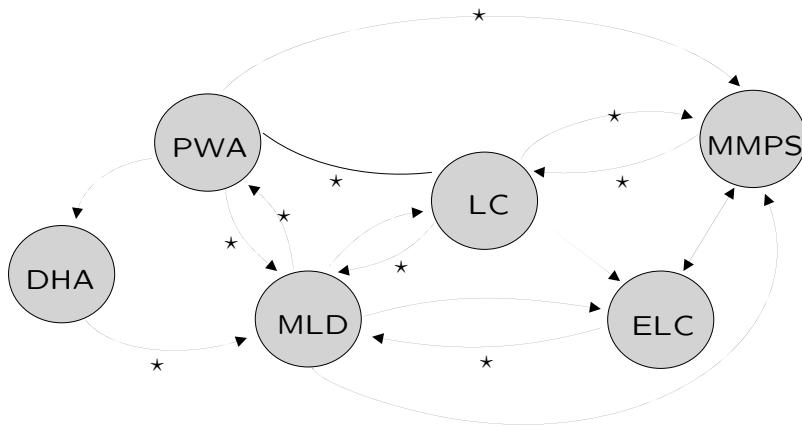
MMPS function: defined by the grammar
 $M := x_i | \alpha | \max(M_1, M_2) | \min(M_1, M_2) | M_1 + M_2 | \beta M_1$

Example: $x(t+1) = 2 \max(x(t), 0) + \min(-\frac{1}{2}u(t), 1)$

Used for modeling discrete-event systems (t=event counter)

50/150

Equivalence Results



Theorem All the above six classes of discrete-time hybrid models are equivalent (possibly under some additional assumptions, such as boundedness of input and state variables)

(Heemels, De Schutter, Bemporad, *Automatica*, 2001)

(Torrisi, Bemporad, *IEEE CST*, 2003)

(Bemporad and Morari, *Automatica*, 1999)

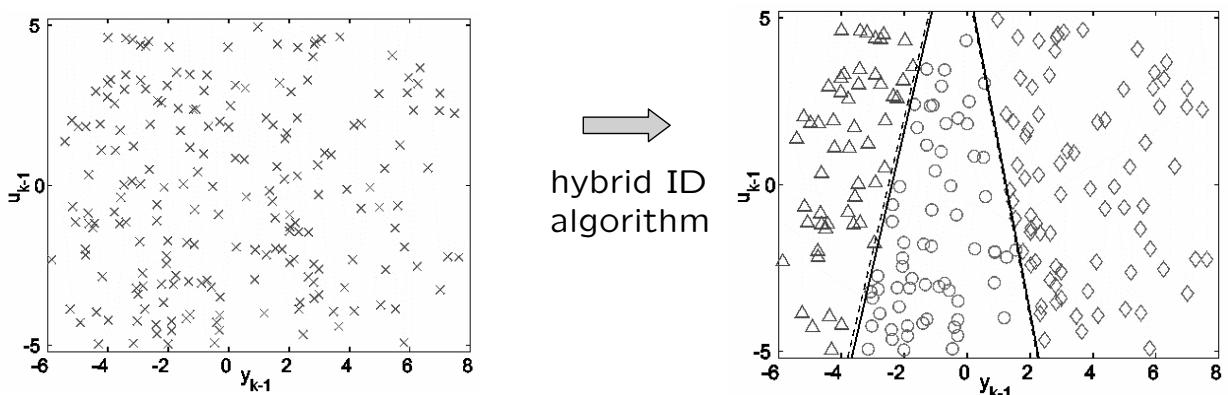
(Bemporad, Ferrari-T.,Morari, *IEEETAC*, 2000)

Theoretical properties and analysis/synthesis tools can be transferred from one class to another

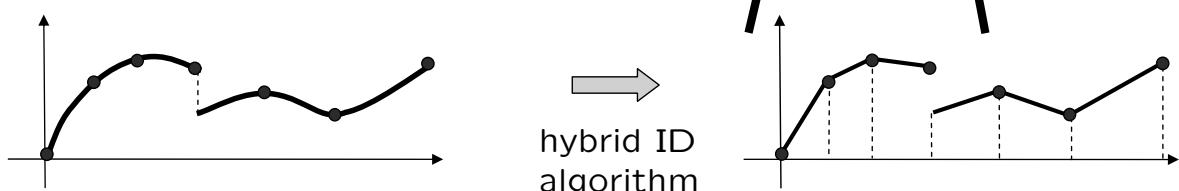
51/150

Hybrid Systems Identification

Given I/O data, estimate the parameters of the affine submodels **and** the partition of the PWA map



Other scenario: “hybridization” of (known) nonlinear models:



52/150

PWA Identification Problem

A. Known Guardlines (partition known, parameters unknown):
least-squares problem
(Ljung's ID TBX) **EASY PROBLEM**

B. Unknown Guardlines (partition *and* parameters unknown):
Generally non-convex, local minima **HARD PROBLEM!**

Some recent approaches to Hybrid ID:

- K-means clustering in a feature space (Ferrari-Trecate, Muselli, Liberati, Morari, 2003)
- Bayesian approach (Juloski, Heemels, Weiland, 2004)
- Mixed-integer programming (Roll, Bemporad, Ljung, 2004)
- Bounded error (partition of infeasible set of inequalities) (Bemporad, Garulli, Paoletti, Vicino, 2003)
- Algebraic geometric approach (Vidal, Soatto, Sastry, 2003)
- Hyperplane clustering in data space (Münz, Krebs, 2002)

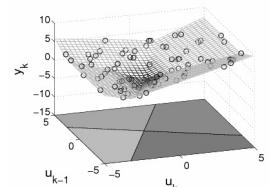
53/150

Hybrid Identification Toolboxes

- PWAID – Piecewise Affine Identification Toolbox (Paoletti, Roll, 2007)

<http://www.control.isy.liu.se/~roll/PWAID/>

- Mixed-integer approach (hinging-hyperplane models)
- Bounded error approach



-
- HIT – Hybrid Identification Toolbox (Ferrari-Trecate, 2006)

http://www-rocq.inria.fr/who/Giancarlo.Ferrari-Trecate/HIT_toolbox.html

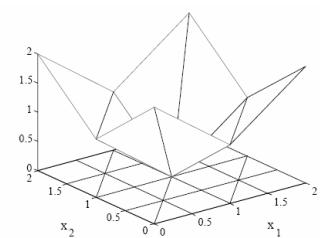


- K-means clustering in a feature space
-

- PWL Toolbox (Julian, 2000)

<http://www.pedrojulian.com>

- High level canonical PWL representations



54/150

Outline

- ✓ Basics of Model Predictive Control (MPC)
- ✓ Hybrid models for MPC
 - MPC of hybrid systems
 - Explicit MPC (multiparametric programming)
 - Optimization-based reachability analysis
 - Examples

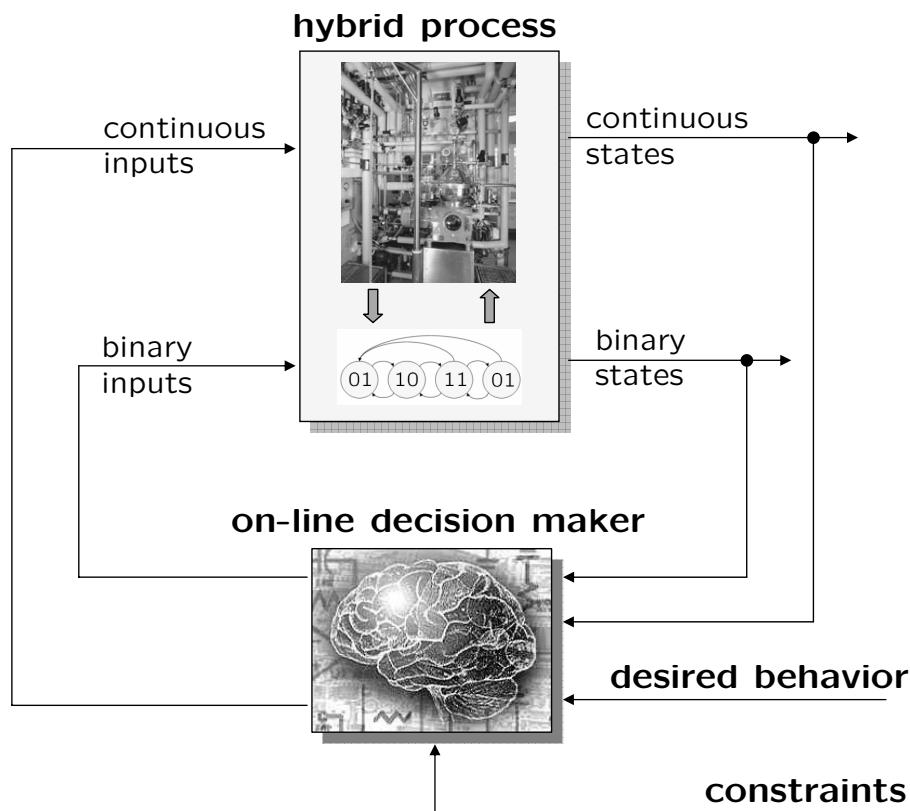
55/150

MPC of Hybrid Systems

- Problem setup
- Convergence properties
- Computational aspects

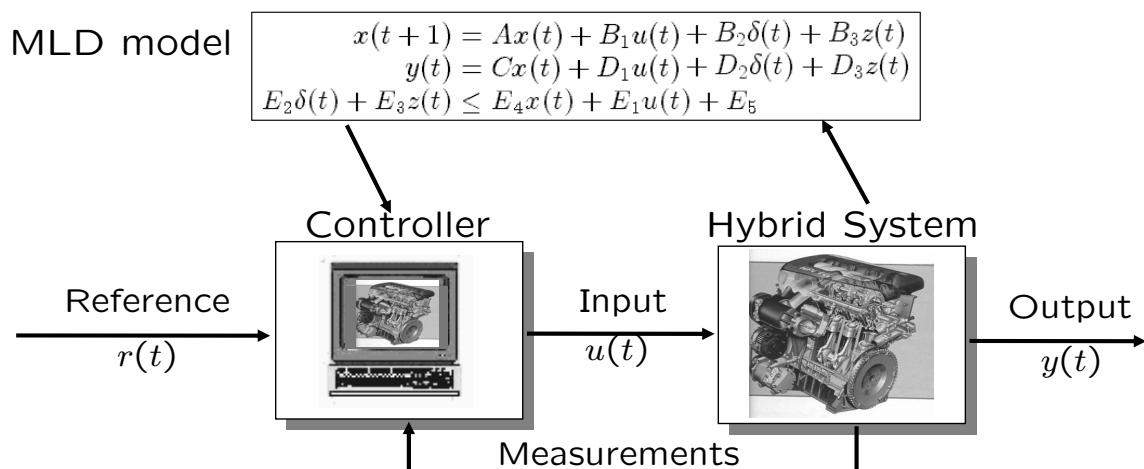
56/150

Hybrid Control Problem



57/150

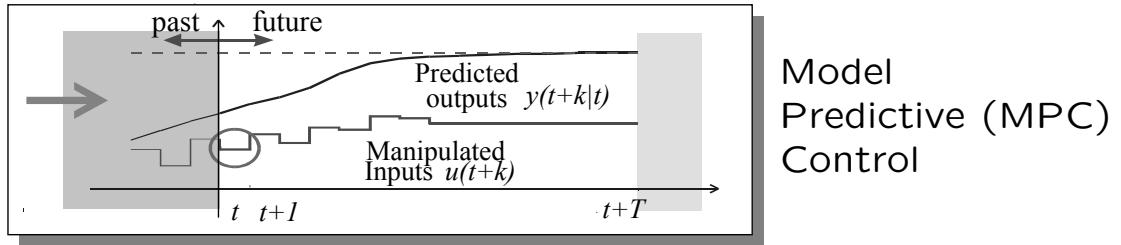
Model Predictive Control of Hybrid Systems



- MODEL: use an MLD (or PWA) model of the plant to predict the future behavior of the hybrid system
- PREDICTIVE: optimization is still based on the predicted future evolution of the hybrid system
- CONTROL: the goal is to control the hybrid system

58/150

MPC for Hybrid Systems



- At time t solve with respect to $U \triangleq \{u(t), u(t+1), \dots, u(t+T-1)\}$ the finite-horizon open-loop, optimal control problem:

$$\min_{u(t), \dots, u(t+T-1)} \sum_{k=0}^{T-1} \|y(t+k|t) - r(t)\| + \rho \|u(t+k) - u_r\| \\ + \sigma (\|\delta(t+k) - \delta_r\| + \|z(t+k) - z_r\| + \|x(t+k|t) - x_r\|)$$

subject to MLD model

$$x(t|t) = x(t)$$

$$x(t+T|t) = x_r$$

- Apply only $u(t)=u_t^*$ (discard the remaining optimal inputs)
- At time $t+1$: get new measurements, repeat optimization

59/150

Closed-Loop Convergence

Theorem 1 Let $(x_r, u_r, \delta_r, z_r)$ be the equilibrium values corresponding to the set point r , and assume $x(0)$ is such that the MPC problem is feasible at time $t = 0$. Then $\forall Q, R \succ 0, \forall \sigma > 0$

$$\lim_{t \rightarrow \infty} y(t) = r \\ \lim_{t \rightarrow \infty} u(t) = u_r$$

$\lim_{t \rightarrow \infty} x(t) = x_r, \lim_{t \rightarrow \infty} \delta(t) = \delta_r, \lim_{t \rightarrow \infty} z(t) = z_r$, and all constraints are fulfilled.

(Bemporad, Morari 1999)

Proof: Easily follows from standard Lyapunov arguments

More stability results: see (Lazar, Heemels, Weiland, Bemporad, 2006)

60/150

Convergence Proof

- Assume we set the terminal constraint $x(t+T|t) = x_r$ in the optimal control problem
- Let \mathcal{U}_t^* denote the optimal control sequence $\{u_t^*(0), \dots, u_t^*(T-1)\}$
- Let $V(t) \triangleq J(\mathcal{U}_t^*, x(t))$ = value function \rightarrow Lyapunov function
- By construction, $\mathcal{U}_1 = \{u_1^*(1), \dots, u_1^*(T-1), u_r\}$ is feasible @ $t+1$
- Hence,

$$V(t+1) \leq J(\mathcal{U}_1, x(t+1)) = V(t) - \|y(t) - r\|_Q - \|u(t) - u_r\|_R - \sigma(\|\delta(t) - \delta_r\| - \|z(t) - z_r\| - \|x(t) - x_r\|)$$

- Hence $V(t)$ is decreasing and lower-bounded by 0 $\Rightarrow \exists V_\infty = \lim_{t \rightarrow \infty} V(t)$
 $\Rightarrow V(t+1) - V(t) \rightarrow 0$
- Hence, $\|y(t) - r\|_Q \rightarrow 0, \|u(t) - u_r\|_R \rightarrow 0, \dots, \|x(t) - x_r\| \rightarrow 0$

Note: Global optimum not needed for convergence !

61/150

Hybrid MPC - Example

PWA system:

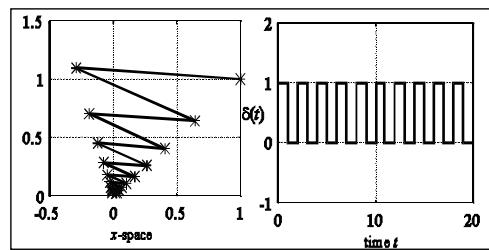
$$x(t+1) = 0.8 \begin{bmatrix} \cos \alpha(t) & -\sin \alpha(t) \\ \sin \alpha(t) & \cos \alpha(t) \end{bmatrix} x(t) + \begin{bmatrix} 0 \\ 1 \end{bmatrix} u(t)$$

$$y(t) = x_2(t)$$

$$\alpha(t) = \begin{cases} \frac{\pi}{3} & \text{if } x_1(t) > 0 \\ -\frac{\pi}{3} & \text{if } x_1(t) \leq 0 \end{cases}$$

Constraint: $-1 \leq u(t) \leq 1$

Open loop behavior



go to demo /demos/hybrid/bm99sim.m

62/150

Hybrid MPC - Example

HYSDEL
model

```
/*
  2x2 PWA system - Example from the paper
  A. Bemporad and M. Morari, ``Control of systems integrating logic, dynamics,
  and constraints,''
  Automatica, vol. 35, no. 3, pp. 407-427, 1999.
  (C) 2003 by A. Bemporad, 2003 */

SYSTEM pwa {
    INTERFACE {
        STATE { REAL x1 [-10,10];
                 REAL x2 [-10,10]; }

        INPUT { REAL u [-1.1,1.1]; }

        OUTPUT{ REAL y; }
    }

    PARAMETER {
        REAL alpha = 1.0472; /* 60 deg in radians */
        REAL C = cos(alpha);
        REAL S = sin(alpha); }

    IMPLEMENTATION {
        AUX { REAL z1,z2;
              BOOL sign; }

        AD { sign = x1<=0; }

        DA { z1 = {IF sign THEN 0.8*(C*x1+S*x2)
                   ELSE 0.8*(C*x1-S*x2) };
             z2 = {IF sign THEN 0.8*(-S*x1+C*x2)
                   ELSE 0.8*(S*x1+C*x2) }; }

        CONTINUOUS { x1 = z1;
                     x2 = z2+u; }

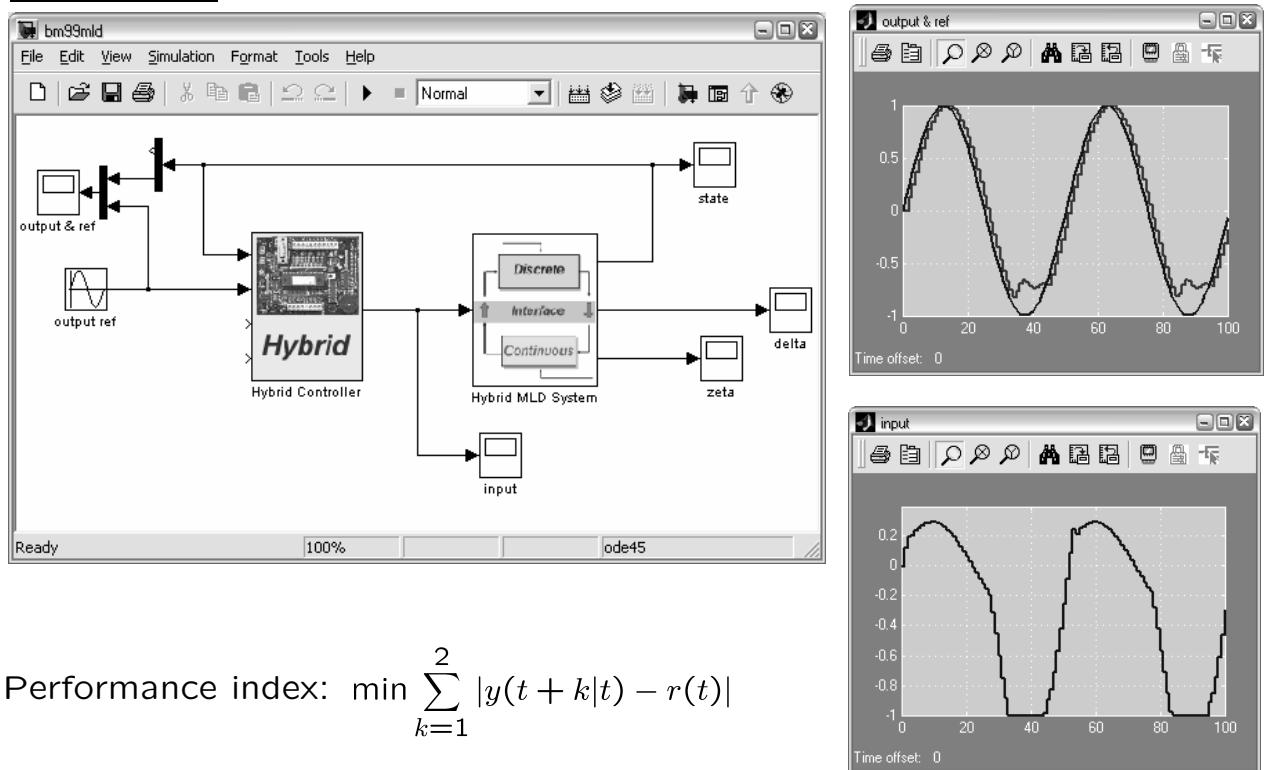
        OUTPUT { y = x2; }
    }
}
```

/demos/hybrid/bm99.hys

63/150

Hybrid MPC - Example

Closed loop:



64/150

Hybrid MPC – Temperature Control

```
>>refs.x=2; % just weight state #2
>>Q.x=1;
>>Q.rho=Inf; % hard constraints
>>Q.norm=2; % quadratic costs
>>N=2; % optimization horizon
>>limits.xmin=[25;-Inf];
```

```
>>C=hybcon(S,Q,N,limits,refs);
```

```
>> C
Hybrid controller based on MLD model S <heatcoolmodel.hys>

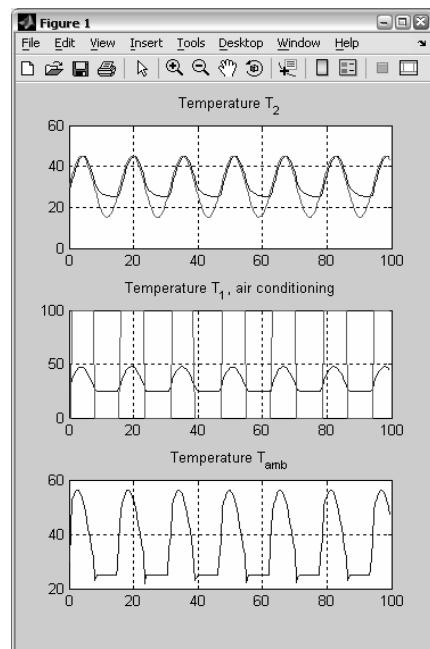
2 state measurement(s)
0 output reference(s)
0 input reference(s)
1 state reference(s)
0 reference(s) on auxiliary continuous z-variables

20 optimization variable(s) (8 continuous, 12 binary)
46 mixed-integer linear inequalities
sampling time = 0.5, MILP solver = 'glpk'

Type "struct(C)" for more details.
>>
```

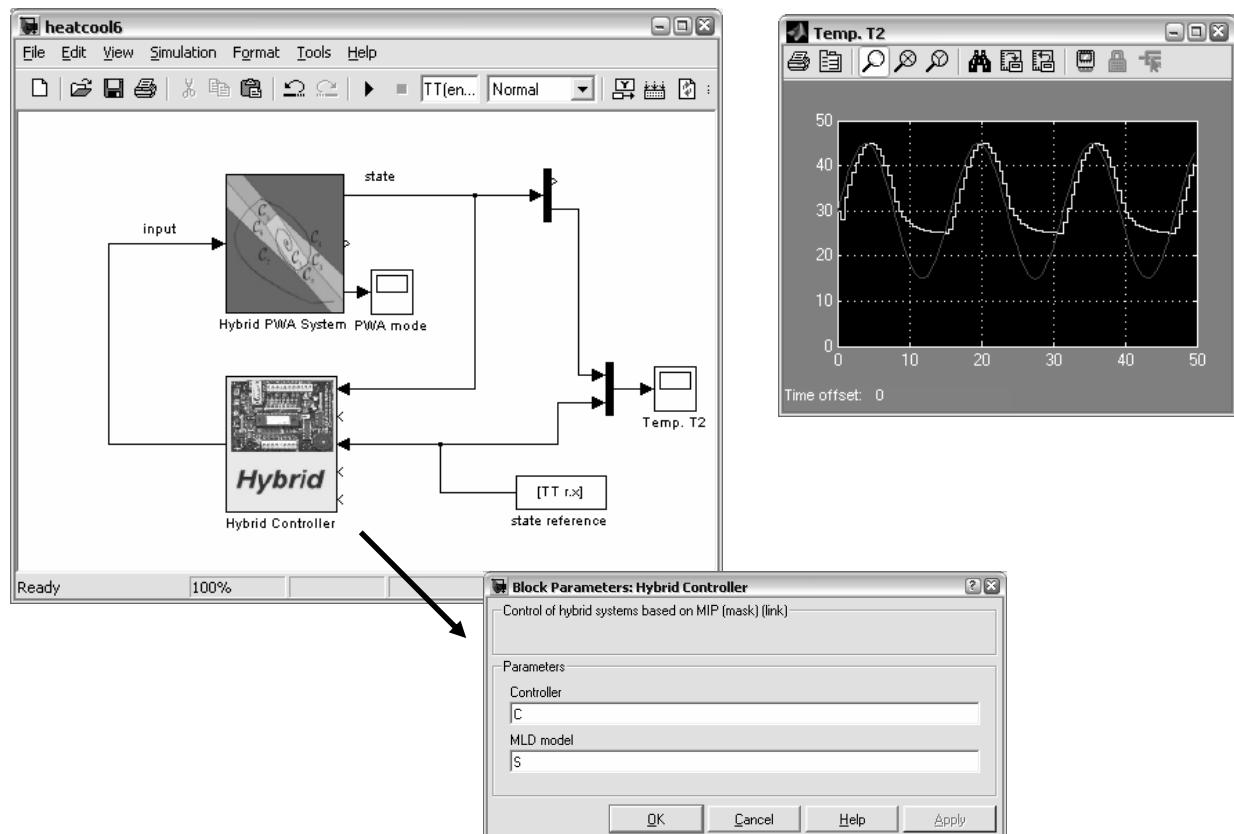
```
>> [XX,UU,DD,ZZ,TT]=sim(C,S,r,x0,Tstop);
```

$$\begin{aligned} \text{min } & \sum_{k=1}^2 (x_2(k) - r)^2 \\ \text{s.t. } & x_1(k) \geq 25 \quad k = 1, 2 \\ & \text{MLD model} \end{aligned}$$



65/150

Hybrid MPC – Temperature Control



66/150

Optimal Control of Hybrid Systems: Computational Aspects



67/150

MIQP Formulation of MPC

(Bemporad, Morari, 1999)

$$\begin{aligned} \min_{\xi} J(\xi, x(0)) &= \sum_{t=0}^{T-1} y'(t) Q y(t) + u'(t) R u(t) \\ \text{subject to } & \begin{cases} x(t+1) = Ax(t) + B_1 u(t) + B_2 \delta(t) + B_3 z(t) + B_5 \\ y(t) = Cx(t) + D_1 u(t) + D_2 \delta(t) + D_3 z(t) + D_5 \\ E_2 \delta(t) + E_3 z(t) \leq E_4 x(t) + E_1 u(t) + E_5 \end{cases} \end{aligned}$$

- Optimization vector:

$$\xi = [u(0), \dots, u(T-1), \delta(0), \dots, \delta(T-1), z(0), \dots, z(T-1)]'$$



$$\begin{aligned} & \min_{\xi} \frac{1}{2} \xi' H \xi + x(0)' F \xi + \cancel{\frac{1}{2} x'(0)' Y x(0)} \\ & \text{subj. to } G \xi \leq W + S x(t) \end{aligned}$$

**Mixed Integer
Quadratic
Program
(MIQP)**

$$u \in \mathbb{R}^{n_u}, \delta \in \{0, 1\}^{n_\delta}, z \in \mathbb{R}^{n_z}$$



$$\xi \in \mathbb{R}^{(n_u+n_z)T} \times \{0, 1\}^{n_\delta T}$$

ξ has both real and $\{0, 1\}$ components

68/150

MILP Formulation of MPC

(Bemporad, Borrelli, Morari, 2000)

$$\begin{aligned} \min_{\xi} \quad & J(\xi, x(0)) = \sum_{t=0}^{T-1} \|Qy(t)\|_{\infty} + \|Ru(t)\|_{\infty} \\ \text{subject to} \quad & \text{MLD model} \end{aligned}$$

- Basic trick: introduce slack variables

$$\begin{array}{ccc} \min_x |x| & \rightarrow & \min_{x, \epsilon} \epsilon \\ \text{s.t. } \epsilon \geq x \\ \epsilon \geq -x \end{array}$$

$$\begin{array}{ccc} \left\{ \begin{array}{l} \epsilon_k^x \geq \|Qy(t+k|t)\|_{\infty} \\ \epsilon_k^u \geq \|Ru(t+k)\|_{\infty} \end{array} \right. & \rightarrow & \left\{ \begin{array}{ll} \epsilon_k^x \geq [Qy(t+k|t)]_i & i = 1, \dots, p \quad k = 1, \dots, T-1 \\ \epsilon_k^x \geq -[Qy(t+k|t)]_i & i = 1, \dots, p \quad k = 1, \dots, T-1 \\ \epsilon_k^u \geq [Ru(t+k)]_i & i = 1, \dots, m \quad k = 0, \dots, T-1 \\ \epsilon_k^u \geq -[Ru(t+k)]_i & i = 1, \dots, m \quad k = 0, \dots, T-1 \end{array} \right. \end{array}$$

- Optimization vector:

$$\xi = [\epsilon_1^x, \dots, \epsilon_{T-1}^x, \epsilon_0^u, \dots, \epsilon_{T-1}^u, u(0), \dots, u(T-1), \delta(0), \dots, \delta(T-1), z(0), \dots, z(T-1)]'$$

$$\begin{array}{ccc} \rightarrow & \begin{array}{c} \min_{\xi} J(\xi, x(0)) = \sum_{k=0}^{T-1} \epsilon_k^x + \epsilon_k^u \\ \text{s.t. } G\xi \leq W + Sx(0) \end{array} & \text{Mixed Integer Linear Program (MILP)} \end{array}$$

ξ has both real and $\{0, 1\}$ components

69/150

Mixed-Integer Program Solvers

- Mixed-integer programming is NP -hard

BUT

- Extremely rich literature in operations research (still very active)

MILP/MIQP is nowadays a technology (CPLEX, Xpress-MP, BARON, GLPK, see e.g. <http://plato.la.asu.edu/bench.html> for a comparison)

- No need to reach the global optimum for stability of MPC (see proof of the theorem), although performance deteriorates

(Ingimundarson, Ocampo-Martinez, Bemporad, 2007)

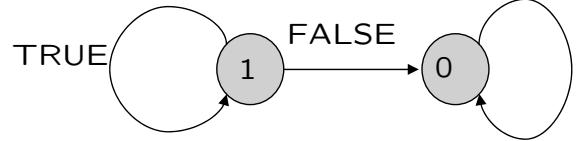
70/150

“Hybrid” Solvers

Main drawback of Mixed-Integer Programming for solving hybrid MPC problems

Loss of the original Boolean structure

$$\begin{cases} x_\ell(k+1) + (1 - x_\ell(k)) + (1 - u_\ell(k)) \geq 1 \\ x_\ell(k) + (1 - x_\ell(k+1)) \geq 1 \\ u_\ell(k) + (1 - x_\ell(k+1)) \geq 1 \end{cases}$$

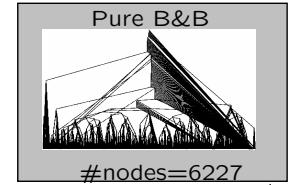
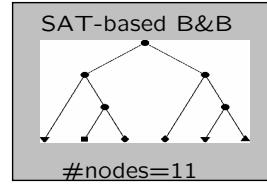


$$x_\ell(k+1) = x_\ell(k) \wedge u_\ell(k)$$

- Possibility of combining symbolic + numerical solvers
Example: SAT + linear programming

N. Vars	N. Cons	Sat instances		Unsat instances	
		zCHAFF	CPLEX	zCHAFF	CPLEX
20	91	0	0.036	-	-
50	218	0	0.343	0	0.453
75	325	0	0.203	0	3.671
100	430	0	23.328	0	33.921
125	538	0.016	15.171	0.031	209.766
150	645	0.031	20.625	0.281	4949.58
175	753	0.031	> 1500	0.891	> 5000

(Bemporad, Giorgetti, IEEE TAC 2006)



71/150

Main Drawbacks of On-line Combinatorial Optimization

- Computation time may be too long: good for large sampling times (>1s), but not for fast sampling applications (e.g. 1ms)
- Requires relatively expensive hardware (e.g. a PC, not suitable on inexpensive µ-controllers)
- Software complexity: MIP solver code difficult to certify (not good for safety critical applications)
- Worst-case computation time is often hard to estimate (over-estimation is usually exaggerated)

Any way to get rid of on-line solvers ?

Outline

- ✓ Basics of Model Predictive Control (MPC)
- ✓ Hybrid models for MPC
- ✓ MPC of hybrid systems
- Explicit MPC (multiparametric programming)
- Optimization-based reachability analysis
- Examples

73/150

Explicit Model Predictive Control

- Multiparametric quadratic programming
- Explicit linear MPC
- Explicit hybrid MPC

74/150

On-Line vs. Off-Line Optimization

$$\begin{aligned} \min_U & \quad \frac{1}{2} U' H U + x'(t) F' U + \frac{1}{2} x'(t) Y x(t) \\ \text{subj. to} & \quad G U \leq W + S x(t), \end{aligned}$$

- **On-line** optimization: given $x(t)$ solve the problem at each time step t (the control law $u=u(x)$ is **implicitly** defined by the QP solver)

→ Quadratic Program (QP)

- **Off-line** optimization: solve the QP **for all** $x(t)$ to find the control law $u=u(x)$ **explicitly**

→ multi-parametric Quadratic Program (mp-QP)

75/150

Multiparametric Quadratic Programming

(Bemporad et al., 2002)

$$\begin{aligned} \min_U & \quad \frac{1}{2} U' H U + x' F' U + \frac{1}{2} x' \cancel{Y} x \\ \text{subj. to} & \quad G U \leq W + S x, \end{aligned}$$

$$U \triangleq [u'_0 \ \dots \ u'_{N-1}]'$$

$$U \in \mathbb{R}^r, \ r \triangleq m N_u$$

$$x \in \mathbb{R}^n$$

- Objective: solve the QP **for all** $x \in X \subseteq \mathbb{R}^n$

- Assumptions: $\begin{bmatrix} H & F \\ F' & Y \end{bmatrix} \succeq 0$ (always satisfied if QP problem originates from optimal control problem)

$$H \succ 0 \quad (\text{can be easily satisfied, e.g. by choosing input weight matrix } > 0)$$

76/150

Linearity of the Solution

- $x_0 \in X$
- solve QP to find $U^*(x_0), \lambda^*(x_0)$
 - identify active constraints at $U^*(x_0)$
 - form matrices $\tilde{G}, \tilde{W}, \tilde{S}$ by collecting active constraints: $\tilde{G}U^*(x_0) - \tilde{W} - \tilde{S}x_0 = 0$

KKT optimality conditions:

$$\begin{aligned} (1) \quad & HU + Fx + G'\lambda = 0, & (2) \quad & \tilde{G}U - \tilde{W} - \tilde{S}x = 0 \\ (3) \quad & \lambda_i(G^i U - W^i - S^i x) = 0, & (4) \quad & \hat{G}U \leq \hat{W} + \hat{S}x \\ (5) \quad & \tilde{\lambda}_i \geq 0, \quad \tilde{\lambda}_i = 0 \end{aligned}$$

From (1) : $U = -H^{-1}(Fx + \tilde{G}'\tilde{\lambda})$

\hat{G} =rows of G not in \tilde{G}
(inactive constraints)

From (2) : $\tilde{\lambda}(x) = -(\tilde{G}H^{-1}\tilde{G}')^{-1}(\tilde{W} + (\tilde{S} + \tilde{G}H^{-1}F)x).$

$$U(x) = H^{-1}[\tilde{G}'(\tilde{G}H^{-1}\tilde{G}')^{-1}(\tilde{W} + (\tilde{S} + \tilde{G}H^{-1}F)x) - Fx]$$

- In some neighborhood of x_0 , λ and U are explicit affine functions of x !

77/150

Determining a Critical Region

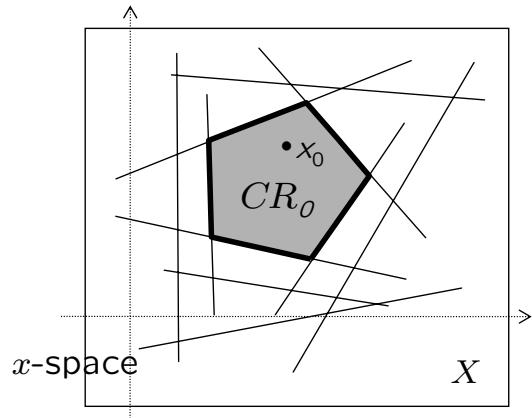
- Impose primal and dual feasibility:
- linear inequalities in x !

$$\begin{aligned} \hat{G}U(x) &\leq \hat{W} + \hat{S}x \\ \tilde{\lambda}(x) &\geq 0 \end{aligned}$$

- Remove redundant constraints:
(this requires solving LP's)

→ critical region CR_0

$$CR_0 = \{x \in X : Ax \leq \mathcal{B}\}$$

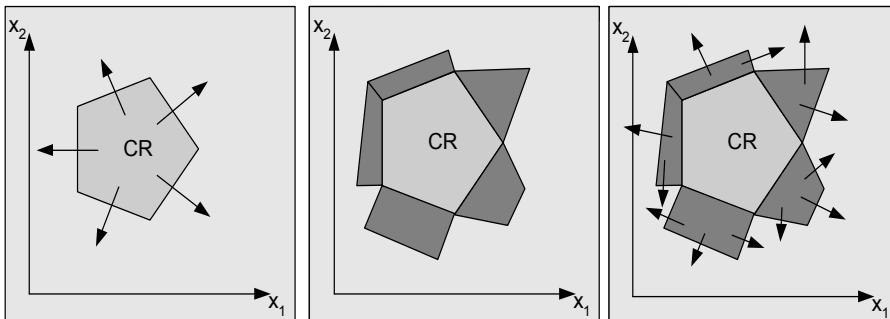


- CR_0 is the set of all and only parameters x for which $\tilde{G}, \tilde{W}, \tilde{S}$ is the optimal combination of active constraints at the optimizer

78/150

Getting All Critical Regions

(Tøndel, Johansen, Bemporad, 2003)



The active set of a neighboring region is found by using the active set of the current region + knowledge of the type of hyperplane we are crossing:

$\hat{G}^i U(x) \leq \hat{W}^i + \hat{S}^i x \Rightarrow$ The corresponding constraint is **added** to the active set

$\tilde{\lambda}_j(x) \geq 0 \Rightarrow$ The corresponding constraint is **withdrawn** from the active set

79/150

Properties of multiparametric-QP

Theorem 1 Consider a multi-parametric quadratic program with $H \succ 0$, $\begin{bmatrix} H & F \\ F' & Y \end{bmatrix} \succeq 0$. The set X^* of parameters x for which the problem is feasible is a polyhedral set, the value function $V^* : X^* \mapsto \mathbb{R}$ is piecewise quadratic, convex and continuous and the optimizer $U^* : X^* \mapsto \mathbb{R}^r$ is piecewise affine and continuous.

$$U^*(x) = \arg \min_U \frac{1}{2} U' H U + x' F' U \quad \text{continuous, piecewise affine}$$

subj. to $GU \leq W + Sx$

$$V^*(x) = \frac{1}{2} x' Y x + \min_U \frac{1}{2} U' H U + x' F' U \quad \text{convex, continuous, piecewise quadratic, } C^1 \text{ (if no degeneracy)}$$

subj. to $GU \leq W + Sx$

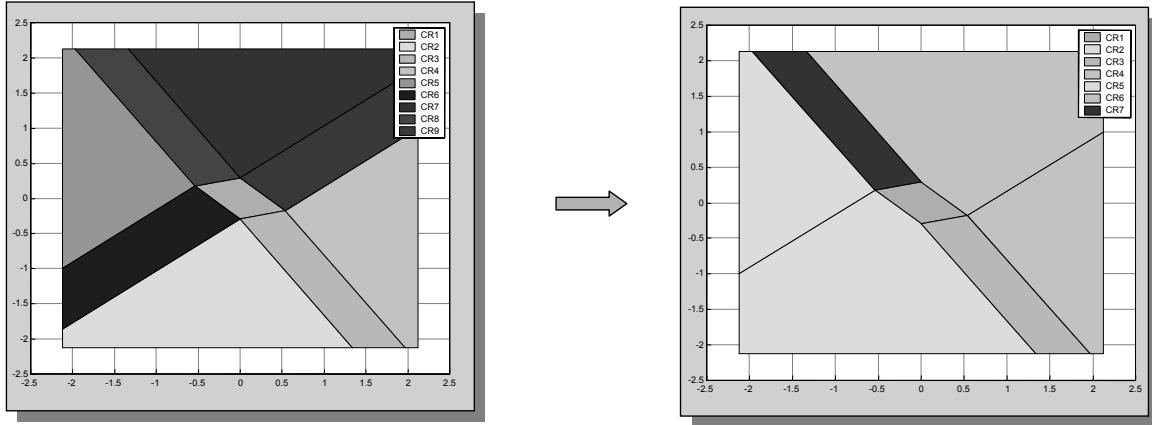
Corollary: The linear MPC controller is a continuous piecewise affine function of the state

$$u(x) = \begin{cases} F_1 x + g_1 & \text{if } H_1 x \leq K_1 \\ \vdots & \vdots \\ F_M x + g_M & \text{if } H_M x \leq K_M \end{cases}$$



80/150

Complexity Reduction



$$U(x) \triangleq [u'_0(x) \ u'_1(x) \ \dots \ u'_{N-1}(x)]'$$

Regions where the first component of the solution is the same can be joined (when their union is convex).

(Bemporad, Fukuda, Torrisi, *Computational Geometry*, 2001)

81/150

Double Integrator Example

- System: $y(t) = \frac{1}{s^2}u(t)$ \rightarrow $x(t+1) = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}x(t) + \begin{bmatrix} 0 \\ 1 \end{bmatrix}u(t)$
sampling + ZOH $T_s=1$ s $y(t) = \begin{bmatrix} 1 & 0 \\ 1 & 0 \end{bmatrix}x(t)$

- Constraints: $-1 \leq u(t) \leq 1$

- Control objective: minimize $\sum_{t=0}^{\infty} y^2(t) + \frac{1}{100}u^2(t)$
 $u(t+k) = K_{LQ}x(t+k|t), \forall k \geq N_u$

- Optimization problem: for $N_u=2$

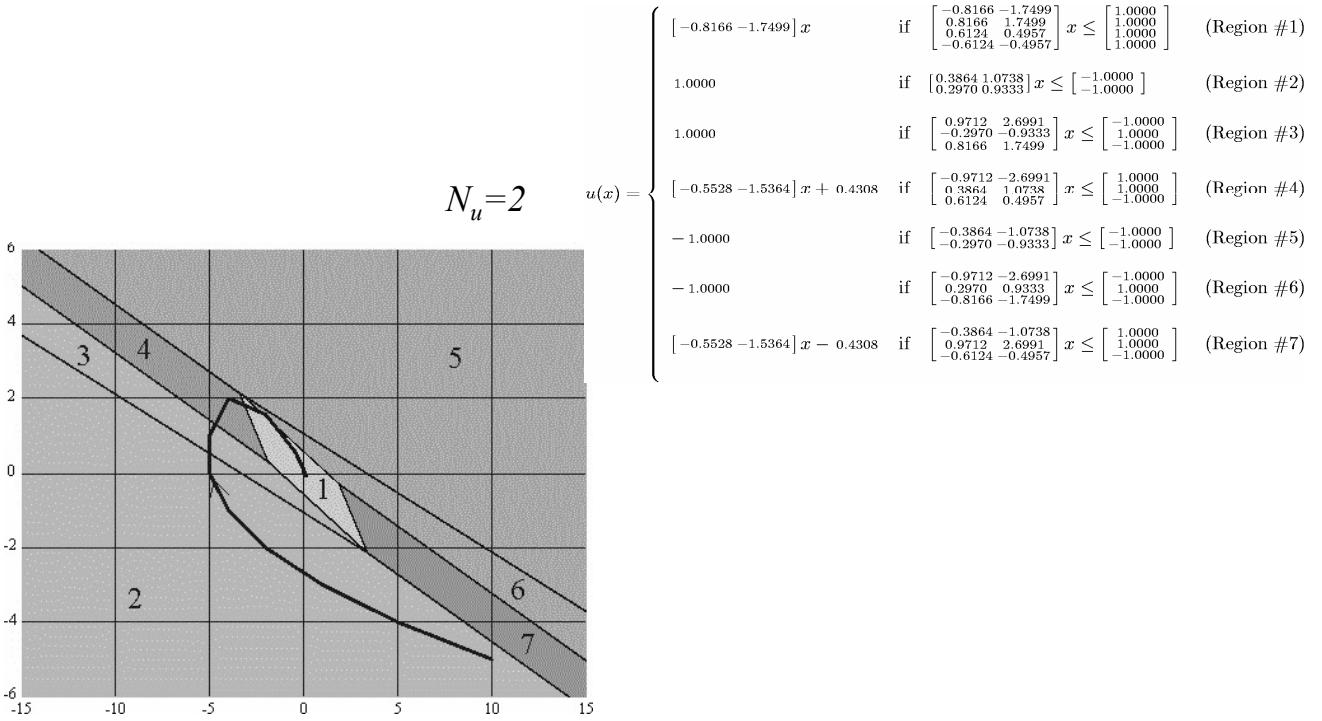
$$H = \begin{bmatrix} 0.8365 & 0.3603 \\ 0.3603 & 0.2059 \end{bmatrix}, \quad F = \begin{bmatrix} 0.4624 & 1.2852 \\ 0.1682 & 0.5285 \end{bmatrix}$$

(cost function is normalized by max svd(H))

$$G = \begin{bmatrix} 1 & 0 \\ -1 & 0 \\ 0 & 1 \\ 0 & -1 \end{bmatrix}, \quad W = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix}, \quad S = \begin{bmatrix} 0 & 0 \\ 0 & 0 \\ 0 & 0 \\ 0 & 0 \end{bmatrix}$$

82/150

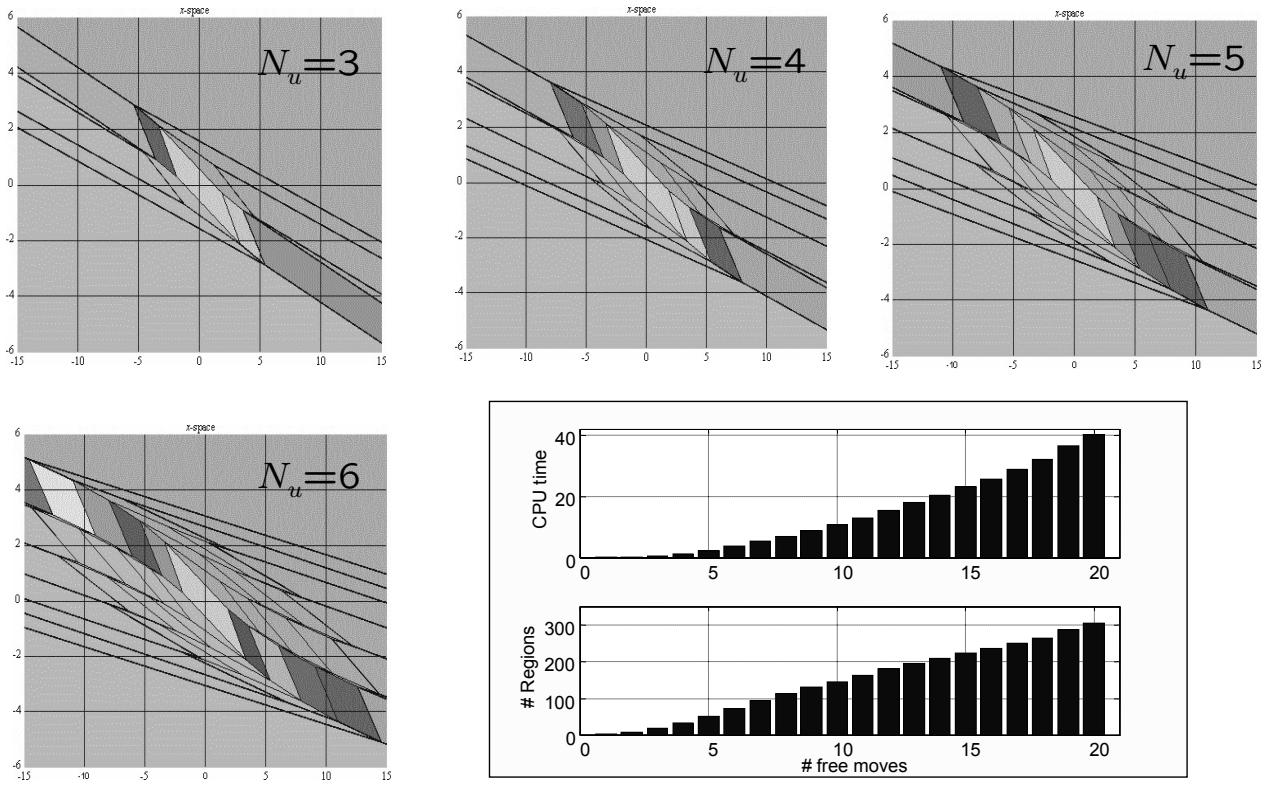
mp-QP solution



go to demo [/demos/linear/doubleintexp.m](#) (Hyb-Tbx)

83/150

Complexity



(is the number of regions finite for $N_u \rightarrow \infty$?)

84/150

Complexity

- Worst-case complexity analysis:

$$M \triangleq \sum_{\ell=0}^q \binom{q}{\ell} = 2^q \quad \text{combinations of active constraints}$$

- Usually the number of regions is much smaller, as many combinations of active constraints are never feasible and optimal at any parameter vector x
- Strongest dependence on the number q of constraints
- Strong dependence on the number N_u of free moves
- Weak dependence on the number n of parameters x

- Example:

states\horizon	$N = 1$	$N = 2$	$N = 3$	$N = 4$	$N = 5$
$n=2$	3	6.7	13.5	21.4	19.3
$n=3$	3	6.9	17	37.3	77
$n=4$	3	7	21.65	56	114.2
$n=5$	3	7	22	61.5	132.7
$n=6$	3	7	23.1	71.2	196.3
$n=7$	3	6.95	23.2	71.4	182.3
$n=8$	3	7	23	70.2	207.9

Data averaged over 20 randomly generated single-input single-output systems subject to input saturation ($q=2N$)

85/150

Reference Tracking, MIMO System

QP-based vs. Explicit MPC:

$2N$	QP (ms) average	worst	explicit (ms) average	worst	regions	[storage kb]
4	1.1	1.5	0.005	0.1	25	16
8	1.3	1.9	0.023	1.1	175	78
20	2.5	2.6	0.038	3.3	1767	811
30	5.3	7.2	0.069	4.4	5162	2465
40	10.9	13.0	0.239	15.6	11519	5598

(Intel Centrino 1.4 GHz)

Average time on 100 random [states,references, previous inputs] Worst-case time on 100 random [states,references, previous inputs]

Organization of explicit solution on a binary tree can improve CPU time (worst & average) from n_r to about $\log_2 n_r$ (n_r =number of regions)

(Tøndel, Johansen, Bemporad, *Automatica*, 2003)

86/150

Example: AFTI-16

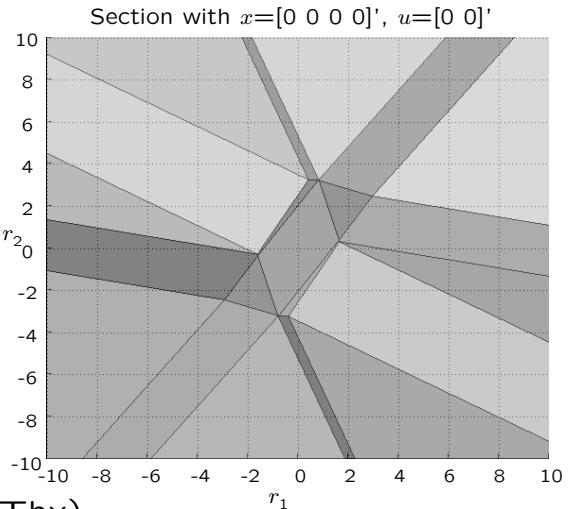
- Linearized model:

$$\left\{ \begin{array}{l} \dot{x} = \begin{bmatrix} -.0151 & -60.5651 & 0 & -32.174 \\ -.0001 & -1.3411 & .9929 & 0 \\ .00018 & 43.2541 & -.86939 & 0 \\ 0 & 0 & 1 & 0 \end{bmatrix} x + \begin{bmatrix} -2.516 & -13.136 \\ -.1689 & -.2514 \\ -17.251 & -1.5766 \\ 0 & 0 \end{bmatrix} u \\ y = \begin{bmatrix} 0 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{bmatrix} x, \end{array} \right.$$



- Inputs: elevator and flaperon angle
- Outputs: attack and pitch angle
- Sampling time: $T_s = .05$ s (+ zero-order hold)
- Constraints: max 25° on both angles
- Open-loop response: unstable
(open-loop poles: $-7.6636, -0.0075 \pm 0.0556j, 5.4530$)

Explicit controller: 8 parameters, 51 regions



go to demo `linear/afti16.m` (Hyb-Tbx)

87/150

Explicit Hybrid MPC (MLD)

$$\min_{\xi} J(\xi, \underline{x}(t)) = \sum_{k=0}^{T-1} \|Q(y_k - \underline{r}(t))\|_{\infty} + \|Ru_k\|_{\infty}$$

subject to

$$\begin{cases} x_{k+1} = Ax_k + B_1u_k + B_2\delta_k + B_3z_k + B_5 \\ y_k = Cx_k + D_1u_k + D_2\delta_k + D_3z_k + D_5 \\ E_2\delta_k + E_3z_k \leq E_4x_k + E_1u_k + E_5 \\ x_0 = \underline{x}(t) \end{cases}$$

- On-line optimization: solve the problem *for each* given $x(t)$

Mixed-Integer Linear Program (MILP)

- Off-line optimization: solve the MILP **for all** $x(t)$ in advance

$$\begin{aligned} \min_{\xi} \quad & \sum_{k=0}^{T-1} \epsilon_i^x + \epsilon_i^u \\ \text{s.t.} \quad & G\xi \leq W + S \begin{bmatrix} \underline{x}(t) \\ \underline{r}(t) \end{bmatrix} \end{aligned}$$

multi-parametric Mixed Integer Linear Program (mp-MILP)

88/150

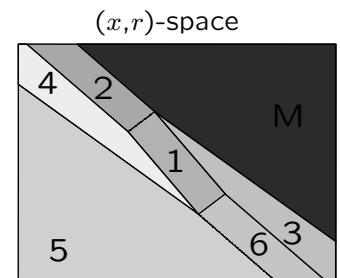
Multiparametric MILP

$$\begin{aligned} \min_{\xi=\{\xi_c, \xi_d\}} \quad & f' \xi_c + d' \xi_d \\ \text{s.t.} \quad & G \xi_c + E \xi_d \leq W + Fx \end{aligned}$$

$\xi_c \in \mathbb{R}^n$
 $\xi_d \in \{0, 1\}^m$

- mp-MILP can be solved (by alternating MILPs and mp-LPs)
(Dua, Pistikopoulos, 1999)
- **Theorem:** The multiparametric solution $\xi^*(x)$ is piecewise affine
- **The MPC controller is piecewise affine in x, r**

$$u(x, r) = \begin{cases} F_1 x + E_1 r + g_1 & \text{if } H_1 [\vec{x}] \leq K_1 \\ \vdots & \vdots \\ F_M x + E_M r + g_M & \text{if } H_M [\vec{x}] \leq K_M \end{cases}$$



89/150

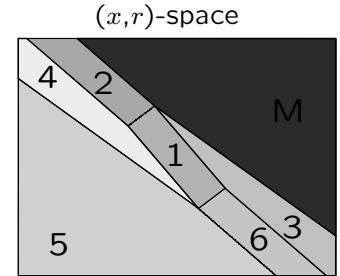
Explicit Hybrid MPC (PWA)

$$\begin{aligned} \min_U J(U, x, r) = & \sum_{k=0}^{T-1} \|R(y(k) - r)\|_p + \|Qu(k)\|_p \\ \text{subject to} & \begin{cases} \text{PWA model} \\ x(0) = x \end{cases} \end{aligned}$$

$p = 1, 2, \infty$
 $\|v\|_2 = v'v$
 $\|v\|_\infty = \max |v_i|$
 $\|v\|_1 = \sum v_i$

- **The MPC controller is piecewise affine in x, r**

$$u(x, r) = \begin{cases} F_1 x + E_1 r + g_1 & \text{if } H_1 [\vec{x}] \leq K_1 \\ \vdots & \vdots \\ F_M x + E_M r + g_M & \text{if } H_M [\vec{x}] \leq K_M \end{cases}$$



Note: in the 2-norm case the partition may not be fully polyhedral

90/150

Computation of Explicit Hybrid MPC (PWA)

Method A: (Borrelli, Baotic, Bemporad, Morari, *Automatica*, 2005)

Use a combination of DP (dynamic programming) and mpLP
(1-norm, ∞ -norm), or mpQP (quadratic forms)

Method B: (Bemporad, *Hybrid Toolbox*, 2003) (Alessio, Bemporad, ADHS 2006)(Mayne, ECC 2001)

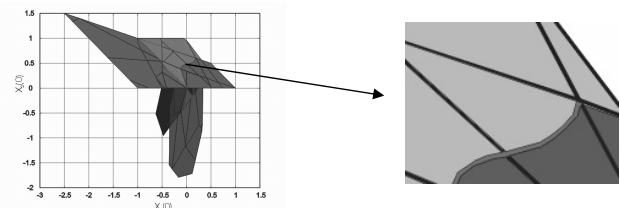
1 - Use backwards (=DP) reachability analysis for enumerating all feasible mode sequences $I = \{i(0), i(1), \dots, i(T-1)\}$;

2 - For each fixed sequence I , solve the explicit finite-time optimal control problem for the corresponding linear time-varying system (mpQP or mpLP);

3 - Case $1/\infty$ -norm: Compare value functions and split regions.

Quadratic case: keep overlapping regions (possibly eliminate overlaps that are never optimal) and compare on-line (if needed).

Note: in the 2-norm case, the fully explicit partition may not be polyhedral



91/150

Hybrid Control Examples (Revisited)

92/150

Hybrid Control Example

PWA system:

$$x(t+1) = 0.8 \begin{bmatrix} \cos \alpha(t) & -\sin \alpha(t) \\ \sin \alpha(t) & \cos \alpha(t) \end{bmatrix} x(t) + \begin{bmatrix} 0 \\ 1 \end{bmatrix} u(t)$$

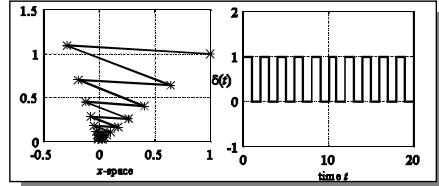
$$y(t) = x_2(t)$$

$$\alpha(t) = \begin{cases} \frac{\pi}{3} & \text{if } x_1(t) \geq 0 \\ -\frac{\pi}{3} & \text{if } x_1(t) < 0 \end{cases}$$

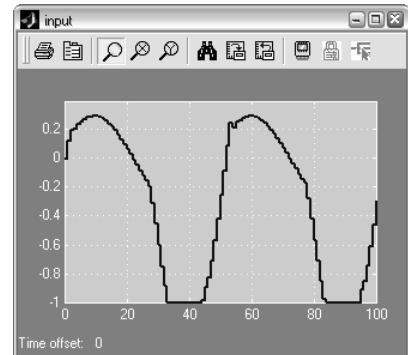
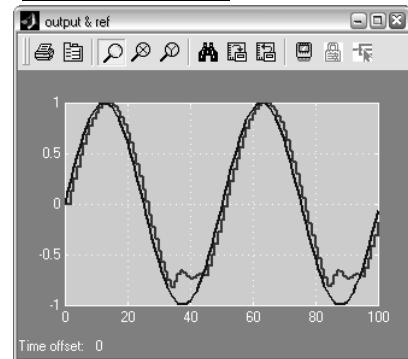
Constraints: $-1 \leq u(t) \leq 1$

Objective: $\min \sum_{k=1}^2 |y(t+k|t) - r(t)|$

Open loop behavior:



Closed loop:

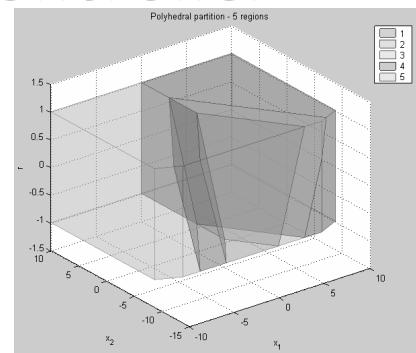


HybTbx: /demos/hybrid/bm99sim.m

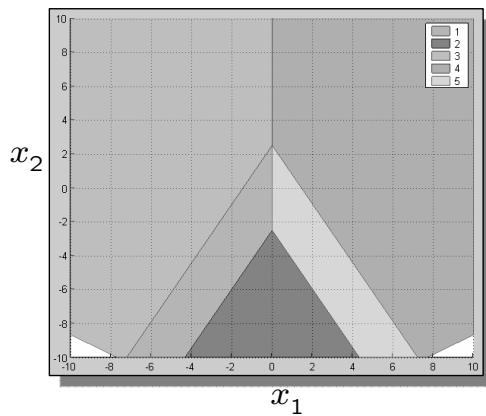
93/150

Explicit PWA Controller

$$u(x, r) = \begin{cases} [0.6928 \ -0.4 \ 1] \begin{bmatrix} x \\ r \end{bmatrix} & \text{if } \begin{bmatrix} 0.6928 & -0.4 & 1 \\ -0.4 & -0.6928 & 0 \\ 0 & 0 & -1 \\ -0.6928 & 0.4 & -1 \\ 1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} 1 \\ 10 \\ 1 \\ 1 \\ 1 \\ 1e-006 \end{bmatrix} \\ & \text{(Region \#1)} \\ 1 & \text{if } \begin{bmatrix} -0.6928 & 0.4 & -1 \\ 0.6928 & 0.4 & -1 \\ 0 & 0 & 1 \\ 0 & 0 & -1 \\ 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} -1 \\ 1 \\ 1 \\ 10 \end{bmatrix} \\ & \text{(Region \#2)} \\ -1 & \text{if } \begin{bmatrix} -0.4 & -0.6928 & 0 \\ 0 & 0 & -1 \\ 1 & 0 & 0 \\ 0.6928 & -0.4 & 1 \\ 0 & 0 & 1 \\ -1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} 1e-006 \\ 10 \\ 10 \\ -1 \\ 1 \\ 10 \end{bmatrix} \\ & \text{(Region \#3)} \\ -1 & \text{if } \begin{bmatrix} -1 & 0 & 0 \\ 0 & -1 & 0 \\ 0.4 & -0.6928 & 0 \\ 0 & -1 & 0 \\ -0.6928 & -0.4 & 1 \\ 0 & 0 & -1 \\ 1 & 0 & 0 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} 0 \\ 10 \\ 10 \\ -1 \\ 1 \\ 10 \end{bmatrix} \\ & \text{(Region \#4)} \\ [-0.6928 \ -0.4 \ 1] \begin{bmatrix} x \\ r \end{bmatrix} & \text{if } \begin{bmatrix} -0.6928 & -0.4 & 1 \\ 0.4 & -0.6928 & 0 \\ 0 & 0 & -1 \\ 0.6928 & 0.4 & -1 \\ -1 & 0 & 0 \\ 0 & 0 & 1 \end{bmatrix} \begin{bmatrix} x \\ r \end{bmatrix} \leq \begin{bmatrix} 1 \\ 10 \\ 1 \\ 1 \\ 1 \\ 0 \end{bmatrix} \\ & \text{(Region \#5)} \end{cases}$$



Section with $r=0$



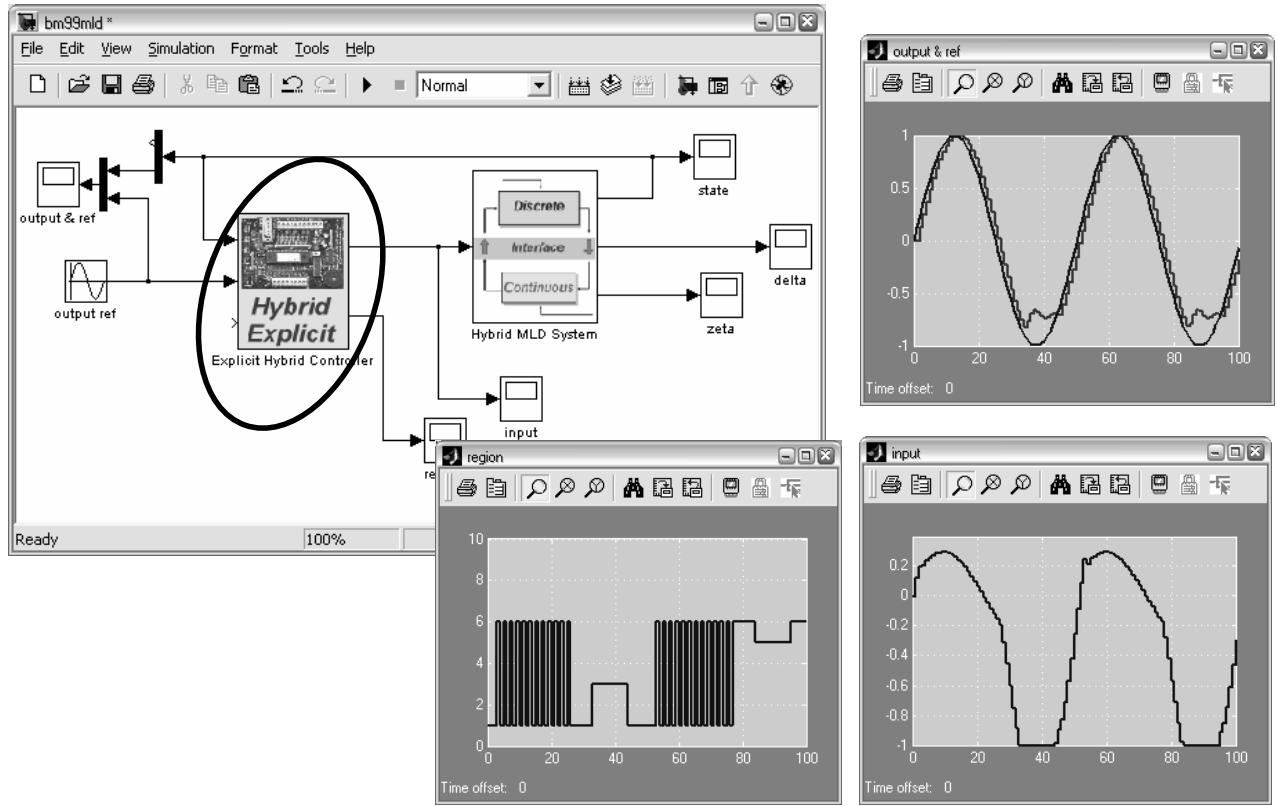
HybTbx: /demos/hybrid/bm99sim.m
(CPU time: 1.51 s, Pentium M 1.4GHz)

PWA law \equiv MPC law !

94/150

Hybrid Control Example

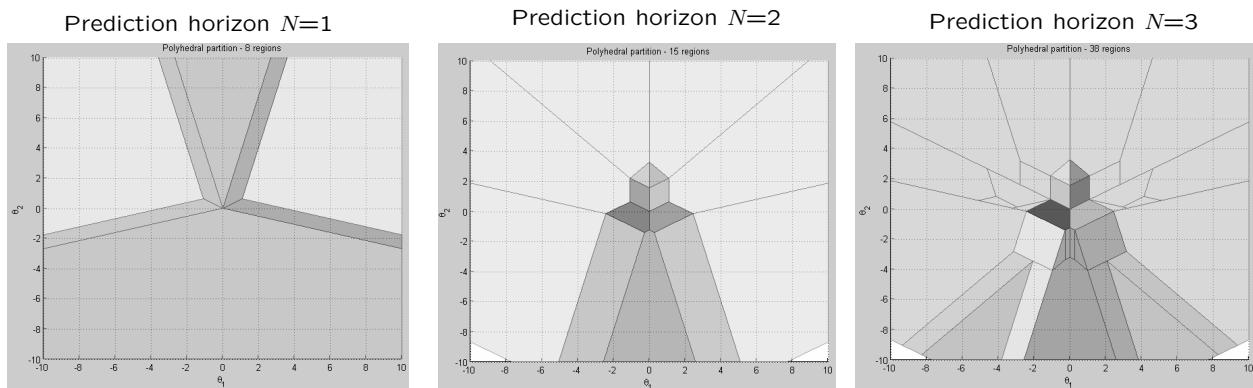
Closed loop:



95/150

Explicit PWA Regulator

$$\text{Objective: } \min \sum_{k=1}^N \|x(t+k|t)\|_\infty$$



HybTbx: [/demos/hybrid/bm99benchmark.m](#)

96/150

Explicit MPC – Temperature Control

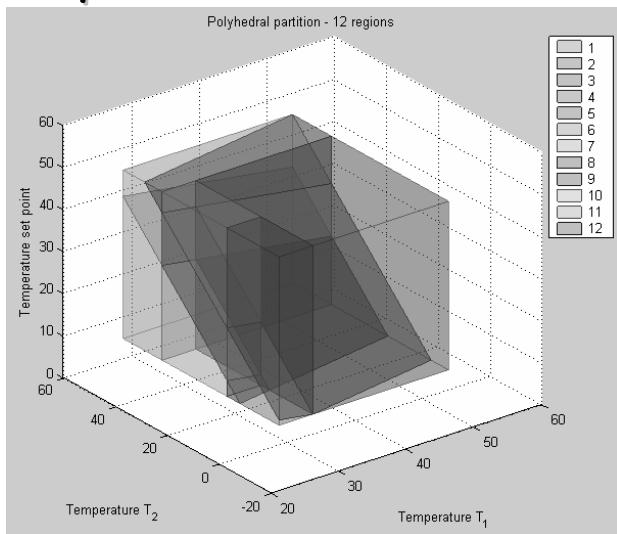
```
>>E=expcon(C,range,options);
```

```
>> E

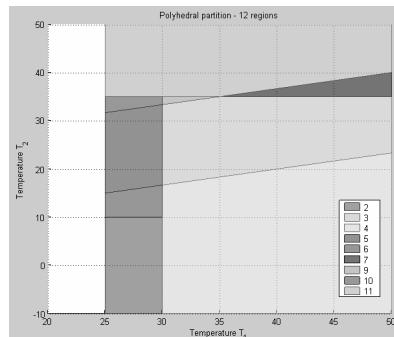
Explicit controller (based on hybrid controller C)
 3 parameter(s)
 1 input(s)
 12 partition(s)
 sampling time = 0.5

The controller is for hybrid systems (tracking)
This is a state-feedback controller.

Type "struct(E)" for more details.
>>
```



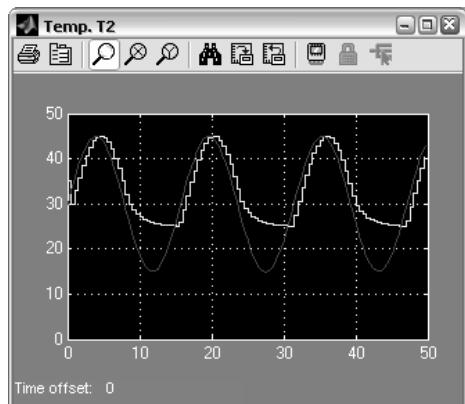
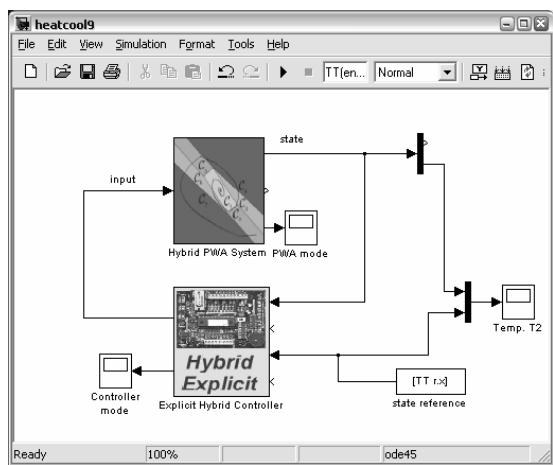
$$\begin{aligned} \min \quad & \sum_{k=1}^2 x_2^2(k) \\ \text{s.t. } & x_1(k) \geq 25 \quad k = 1, 2 \\ & \text{PWA model} \end{aligned}$$



Section in the (T_1, T_2) -space
for $T_{\text{ref}} = 30$

97/150

Explicit MPC – Temperature Control



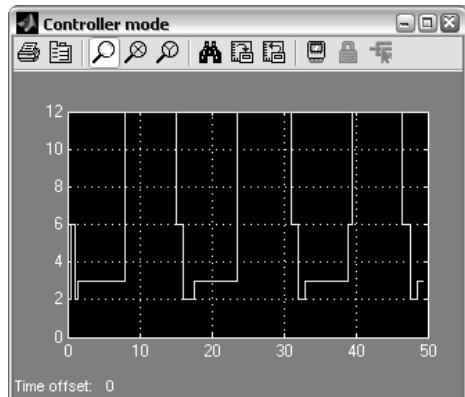
Generated
C-code

utils/expcon.h

```
#define EXPCON_REG 12
#define EXPCON_NTH 3
#define EXPCON_NYH 2
#define EXPCON_NH 72
#define EXPCON_NF 12
static double EXPCON_F[]={
    -1,0,0,0,-1,0,
    -1,-1,-1,-1,-1,0,-3,-3,
    -3,0,-3,0,0,0,0,0,
    0,0,4,4,4,0,4,0,0,
    0,0,0,0};

static double EXPCON_G[]={
    101.6,1.6,1.6,-1.6,98.4001,0,100,51.6,
    101.6,51.6,48.4,50};

static double EXPCON_H[]={
    0,0,0,-0.00999999,0,-0.0333333,
    0.02,0.00999999,-0.02,0,0,-0.0333333,0.02,0.00999999,
    0,0,-0.02,0.02,0,-1,0.00999999,0,
```



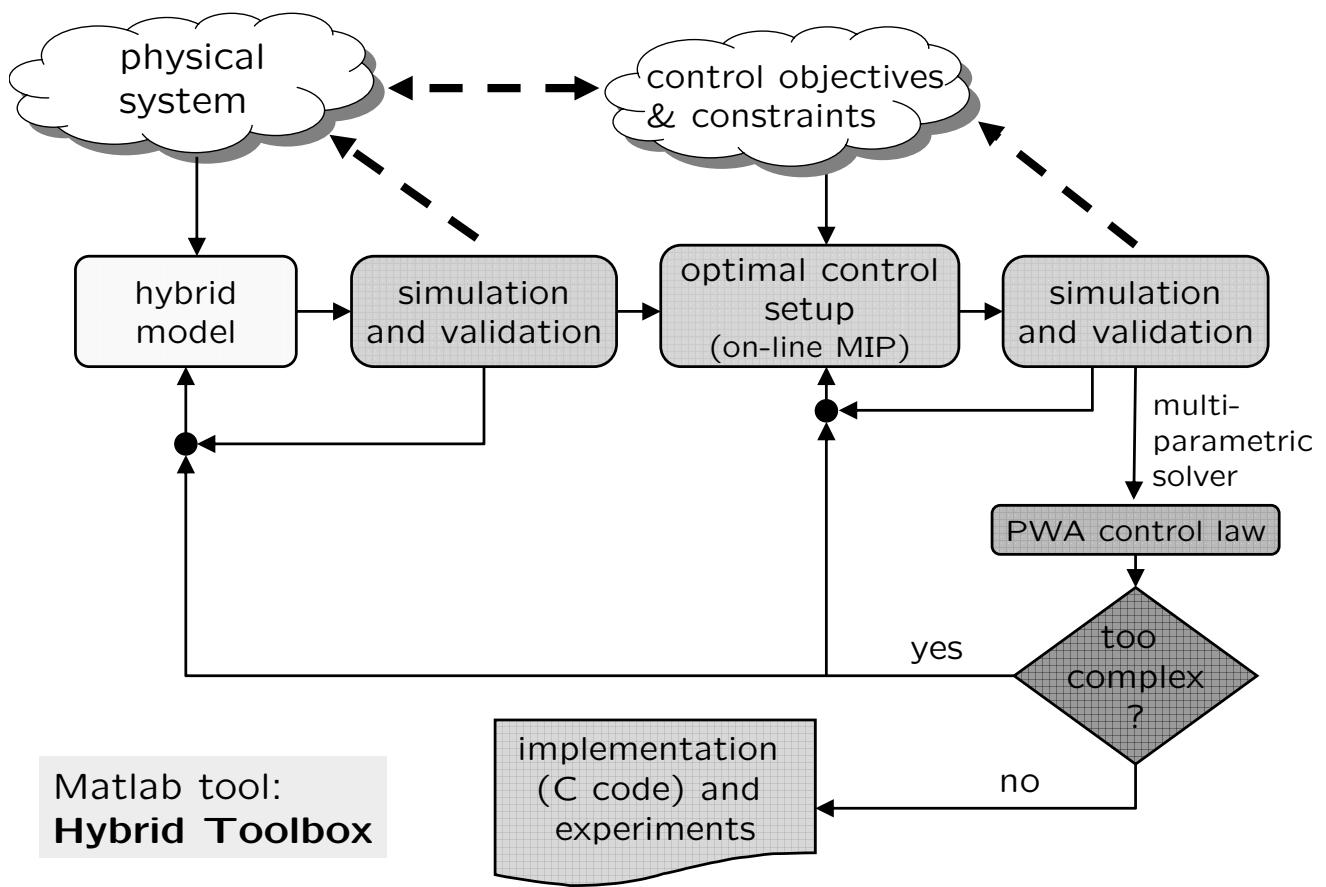
98/150

Implementation Aspects of Hybrid MPC

- **Alternatives:**
 - (1) solve MIP on-line
 - (2) evaluate a PWA function
- **Small problems** (short horizon $N=1,2$, one or two inputs): explicit PWA control law preferable
 - time to evaluate the control law is shorter than MIP
 - control code is simpler (no complex solver must be included in the control software !)
 - more insight in controller's behavior
- **Medium/large problems** (longer horizon, many inputs and binary variables): MIP preferable

99/150

Hybrid Control Design Flow



100/150

A Few Hybrid MPC Tricks ...

101/150

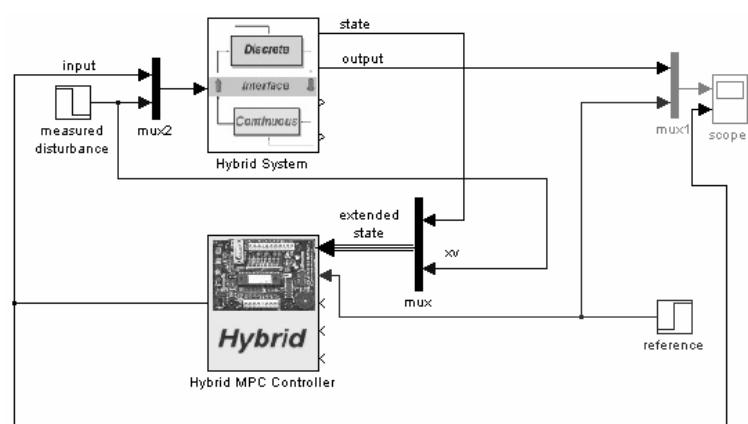
Measured Disturbances

- Disturbance $v(k)$ can be measured at time k
- Augment the hybrid prediction model with a constant state

$$x_v(k+1) = x_v(k)$$

- In Hysdel:

```
INTERFACE{
    STATE{
        REAL x      [-1e3, 1e3];
        REAL xv     [-1e3, 1e3];
    }
    ...
}
IMPLEMENTATION{
    CONTINUOUS{
        x = A*x + B*u + Bv*xv;
        xv = xv;
        ...
    }
}
```



/demos/hybrid/hyb_meas_dist.m

Note: same trick applies to linear MPC

102/150

Hybrid MPC - Tracking

- Optimal control problem (quadratic performance index):

$$\begin{aligned} \min_{\Delta U} \quad & \sum_{k=0}^{N-1} \|W^y(y(k+1) - r(t))\|^2 + \|W^{\Delta u}\Delta u(k)\|^2 \\ & [\Delta u(k) \triangleq u(k) - u(k-1)] \\ \text{subj. to } \quad & u_{\min} \leq u(k) \leq u_{\max}, \quad k = 0, \dots, N-1 \\ & \Delta u_{\min} \leq \Delta u(k) \leq \Delta u_{\max}, \quad k = 0, \dots, N-1 \\ & y_{\min} \leq y(k) \leq y_{\max}, \quad k = 1, \dots, N \end{aligned}$$

- Optimization problem:
(MIQP)

$$\begin{aligned} \min_{\Delta U} \quad & J(\Delta U, x(t)) = \frac{1}{2} \Delta U' H \Delta U + [x'(t) \ r'(t) \ u'(t-1)] F \Delta U \\ \text{s.t.} \quad & G \Delta U \leq W + K \begin{bmatrix} x(t) \\ r(t) \\ u(t-1) \end{bmatrix} \end{aligned}$$

Note: same trick as in linear MPC

103/150

Integral Action in Hybrid MPC

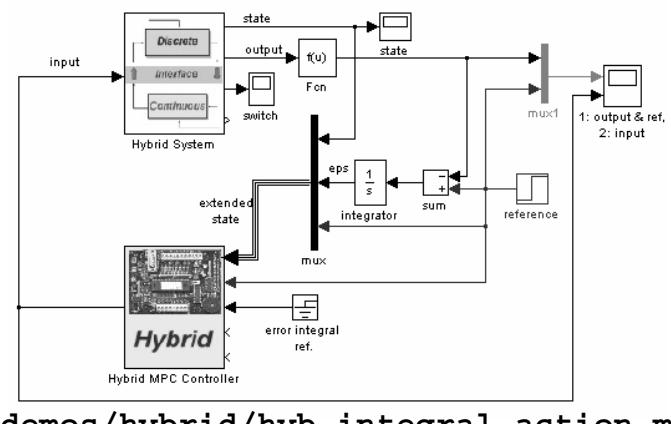
- Augment the hybrid prediction model with integrators of output errors as additional states:

$$\epsilon(k+1) = \epsilon(k) + T_s \cdot (r(k) - y(k)) \quad T_s = \text{sampling time}$$

- Treat $r(k)$ as a measured disturbance (=additional constant state)
- Add weight on $\epsilon(k)$ in cost function to make $\epsilon(k) \rightarrow 0$

- In Hysdel:

```
INTERFACE{
  STATE{
    REAL x           [-100,100];
    ...
    REAL epsilon     [-1e3, 1e3];
    REAL r           [0,      100];
  }
  OUTPUT {
    REAL y;
  }
}
IMPLEMENTATION{
  CONTINUOUS{
    epsilon=epsilon+Ts*(r-(c*x));
    r=r;
    ...
  }
  OUTPUT{
    y=c*x;
  }
}
```



/demos/hybrid/hyb_integral_action.m

Note: same trick applies to linear MPC

104/150

Reference/Disturbance Preview

Measured disturbance $v(t)$ is known M steps in advance

$$\left\{ \begin{array}{l} x_{v,M-1}(k+1) = x_{v,M-2}(k) \\ x_{v,M-2}(k+1) = x_{v,M-3}(k) \\ \vdots \\ x_{v,1}(k+1) = x_{v,0}(k) \\ x_{v,0}(k+1) = x_{v,0}(k) \\ v(k) = x_{v,M-1}(k), k = 0, \dots, N-1 \end{array} \right.$$

initial condition

Note: same trick applies to linear MPC

$$\left\{ \begin{array}{l} x_{v,M-1}(0) = v(t) \\ x_{v,M-2}(0) = v(t+1) \\ \vdots = \vdots \\ x_{v,1}(0) = v(t+M-2) \\ x_{v,0}(0) = v(t+M-1) \end{array} \right.$$

produces $v = \{v(t), v(t+1), \dots, v(t+M-1), v(t+M-1), \dots\}$

Preview of reference $r(t)$: similar.

105/150

Delays – Method 1

- Hybrid model w/ delays:

$$\begin{aligned} x(t+1) &= Ax(t) + B_1 u(t-\tau) + B_2 \delta(t) + B_3 z(t) \\ E_2 \delta(t) + E_3 z(t) &\leq E_1 u(t-\tau) + E_4 x(t) + E_5 \end{aligned}$$

- Map delays to poles in $z=0$:

$$x_k(t) \triangleq u(t-k) \Rightarrow x_k(t+1) = x_{k-1}(t) \quad k = 1, \dots, \tau$$

- Extend the state space of the MLD model:

$$\begin{bmatrix} x \\ x_\tau \\ x_{\tau-1} \\ \vdots \\ x_1 \end{bmatrix} (t+1) = \begin{bmatrix} A & B_1 & 0 & 0 & \dots & 0 \\ 0 & 0 & I_m & 0 & \dots & 0 \\ 0 & 0 & 0 & I_m & \dots & 0 \\ \vdots & \vdots & \vdots & \ddots & \vdots & \vdots \\ 0 & 0 & 0 & 0 & \dots & 0 \end{bmatrix} \begin{bmatrix} x \\ x_\tau \\ x_{\tau-1} \\ \vdots \\ x_1 \end{bmatrix} (t) + \begin{bmatrix} 0 \\ 0 \\ 0 \\ \vdots \\ I_m \end{bmatrix} u(t) + \begin{bmatrix} B_2 \\ 0 \\ 0 \\ \vdots \\ I_m \end{bmatrix} \delta(t) + \begin{bmatrix} B_3 \\ 0 \\ 0 \\ \vdots \\ I_m \end{bmatrix} z(t)$$

- Apply MPC to the extended MLD system

Note: same trick as in linear MPC

106/150

Delays – Method 2

- Delay-free MLD model: $\bar{x}(t+1) = A\bar{x}(t) + B_1u(t) + B_2\bar{\delta}(t) + B_3\bar{z}(t)$
 $E_2\bar{\delta}(t) + E_3\bar{z}(t) \leq E_1u(t) + E_4\bar{x}(t) + E_5$

$$\bar{x}(t) \triangleq x(t+\tau), \bar{\delta}(t) \triangleq \delta(t+\tau), \bar{z}(t) \triangleq z(t+\tau)$$

- Design MPC for delay-free model: $u(t) = f_{\text{MPC}}(\bar{x}(t))$

- Compute the predicted state:

$$\bar{x}(t) = A^\tau x(t) + \sum_{j=0}^{\tau-1} A^{\tau-1-j} (B_1 u(t-1-j) + B_2 \bar{\delta}(t+j) + B_3 \bar{z}(t+j))$$

where $\bar{\delta}(t+j), \bar{z}(t+j)$ are obtained from MLD ineq. (or HYSDEL model)

- Compute MPC action:

$$u(t) = f_{\text{MPC}}(\bar{x}(t))$$

For better closed-loop performance the model used for predicting the future hybrid state $x(t+\tau)$ may be more accurate than MLD model !

107/150

Prioritized Constraints

In optimization problems constraints may be classified in a hierarchy of priorities, where constraints at higher priority must be satisfied before those at lower priority.

Given the optimization problem:

$\min_x \quad f(x)$ s.t. $g(x) \leq 0$	$f : \mathbb{R}^{d_1} \times \{0, 1\}^{d_2} \rightarrow \mathbb{R}$ $g : \mathbb{R}^{d_1} \times \{0, 1\}^{d_2} \rightarrow \mathbb{R}^c$
---	--

1. Define the set of indices of the constraints:

$$\mathcal{C} = \{1, 2, \dots, c\}$$

2. There are r priority levels (i^{th} priority level $>$ $(i+1)^{\text{th}}$ priority level)

Let $\mathcal{P}_i \subseteq \mathcal{C}$ be the set of constraints at priority i .

with $\mathcal{P}_i \cap \mathcal{P}_j = \emptyset$, $i \neq j$ i.e. a constraint cannot be associated with more than one priority level.

108/150

Prioritized Constraints

3. Associate slack variables to each priority level:

$$g_k(x) \leq \epsilon_i, \quad \forall k \in \mathcal{P}_i \quad \epsilon_i \geq 0, \quad \forall i = 1, \dots, r$$

4. Associate binary variables to each priority level:

Each binary variable represents the possible violation of one or more constraints on the same level.

$$[\delta_i = 0] \rightarrow [\epsilon_i = 0] \quad \Leftrightarrow \quad 0 \leq \epsilon_i \leq M_i \delta_i$$

and add:

$$[\delta_j = 0] \rightarrow [\delta_i = 0], \quad \forall j > i \quad \Leftrightarrow \quad \delta_i \leq \delta_j, \quad i < j$$

i.e. if cons at lower priorities are not violated then cons at higher priorities cannot be violated

5. Add constraints in 3 and 4 to the optimization problem and minimize the following cost function:

$$\min f(x) + W' \delta \quad W \in \mathbb{R}^r \text{ "suitably" large}$$

109/150

Optimal Control Example

Consider the following optimal control problem:

$$\begin{aligned} \min \quad & \sum_{k=0}^{N-1} \|y(k) - r_y\|_\infty \\ \text{s.t.} \quad & \text{MLD dynamics} \\ & x_1(k) \geq 1 \\ & x_2(k) \leq 2 \\ & u(k) \geq -1 \end{aligned} \quad \begin{array}{c} \} \\ \} \end{array} \quad \begin{array}{l} \text{with:} \\ \text{Higher priority} \\ \text{Lower priority} \end{array}$$

Problem can be reformulated as follows:

$$\begin{aligned} \min \quad & \sum_{k=0}^{N-1} \|y(k) - r_y\|_\infty + W_1 \delta_1(k) + W_2 \delta_2(k) \\ \text{s.t.} \quad & \text{MLD dynamics} \\ & 1 - x_1(k) \leq \epsilon_1(k) \\ & x_2(k) - 2 \leq \epsilon_1(k) \\ & -u(k) - 1 \leq \epsilon_2(k) \\ & 0 \leq \epsilon_i(k) \leq M_i \delta_i(k), \quad i = 1, 2 \\ & \delta_1 \leq \delta_2 \end{aligned}$$

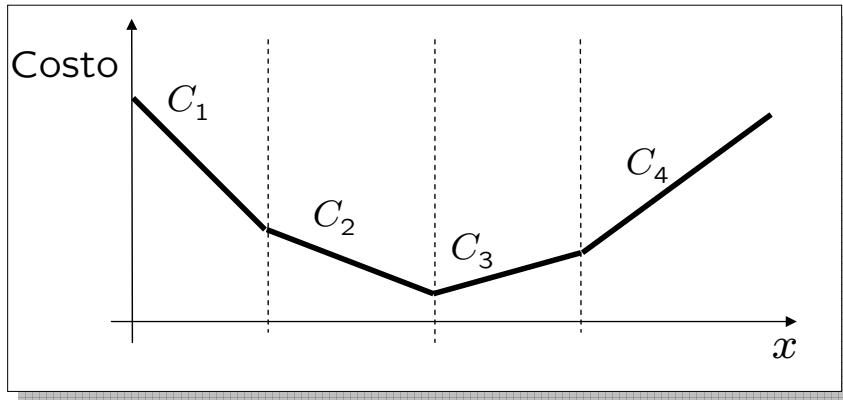
W_1, W_2 are suitable cost weights

```
...
IMPLEMENTATION{
AUX{
REAL eps1[0, M1], eps2[0, M2];
BOOL d1, d2;
}
...
MUST{
1-x1 <= eps1;
x2-2 <= eps1;
-u-1 <= eps2;
eps1 >= 0; eps1 <= M1*d1;
eps2 >= 0; eps2 <= M2*d2;
d1 <= d2;
}
```

In HYSDEL

110/150

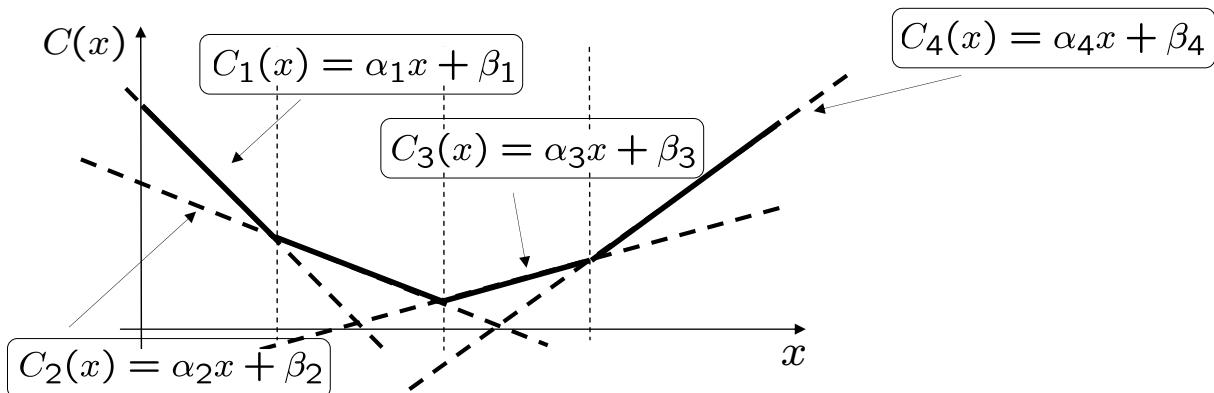
Piecewise Affine Cost Functions



Convex piecewise affine cost functions can be represented **without** introducing binary variables

111/150

Piecewise Affine Cost Functions



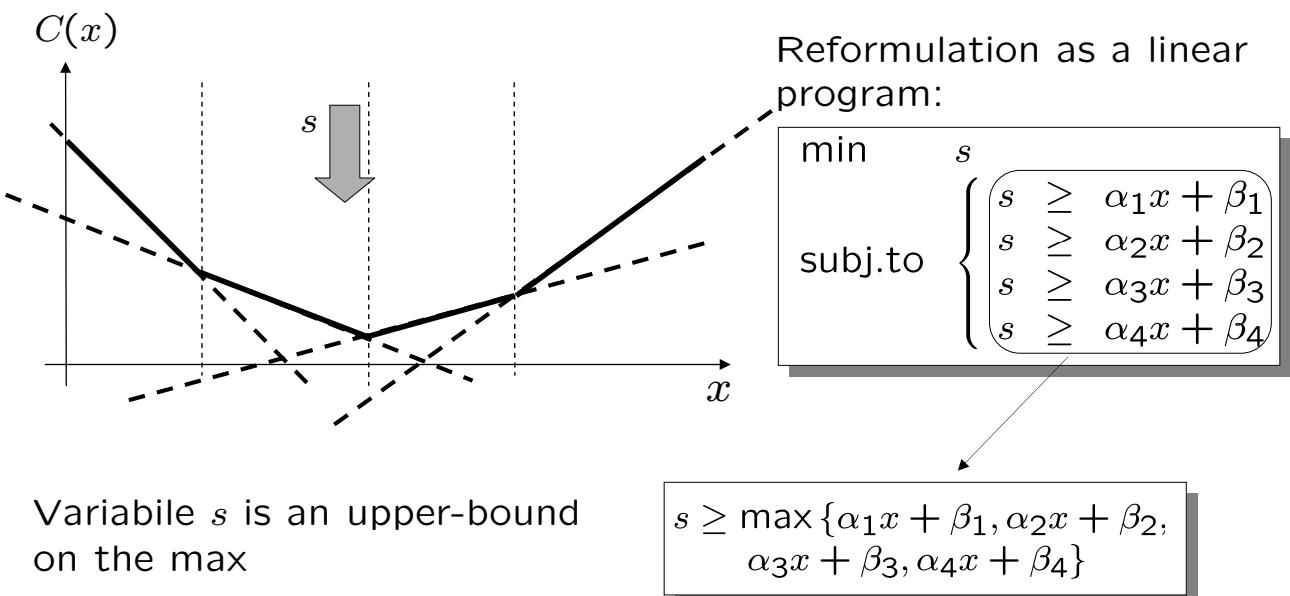
It is easy to see that:

$$C(x) = \max \{\alpha_1x + \beta_1, \alpha_2x + \beta_2, \alpha_3x + \beta_3, \alpha_4x + \beta_4\}$$

In general: Every convex piecewise affine function can be represented as the max of affine functions, and vice versa (Schechter, 1987)

112/150

Piecewise Affine Cost Functions



It is easy to show (by contradiction) that at optimality we have:

$$s = \max \{\alpha_1 x + \beta_1, \alpha_2 x + \beta_2, \alpha_3 x + \beta_3, \alpha_4 x + \beta_4\}$$

Piecewise affine convex constraints can be dealt with similarly

113/150

General Remarks About MIP Modeling

The complexity of solving a mixed-integer program largely depends on the number of integer (binary) variables involved in the problem.

Henceforth, when creating a hybrid model one has to

Be thrifty with integer variables !

Adding logical constraints usually helps ...

Generally speaking:

Modeling is art

(a unifying general theory does not exist)



114/150

Outline

- ✓ Basics of Model Predictive Control (MPC)
- ✓ Hybrid models for MPC
- ✓ MPC of hybrid systems
- ✓ Explicit MPC (multiparametric programming)
 - Optimization-based reachability analysis
 - Examples

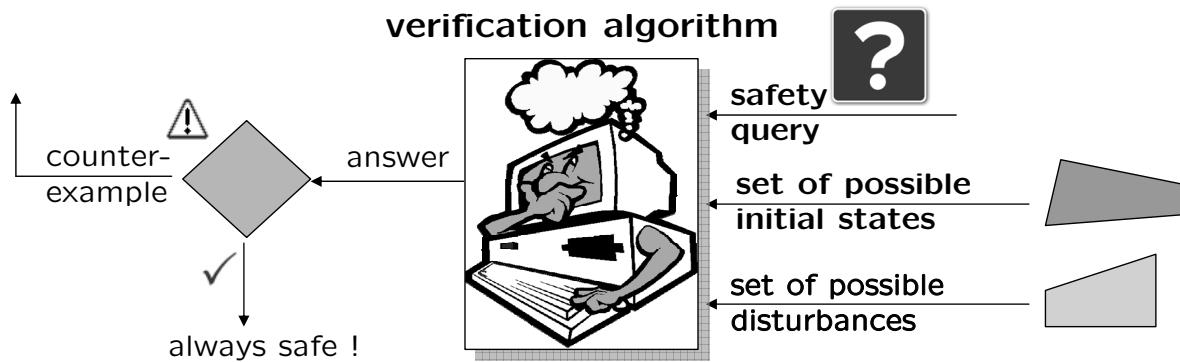
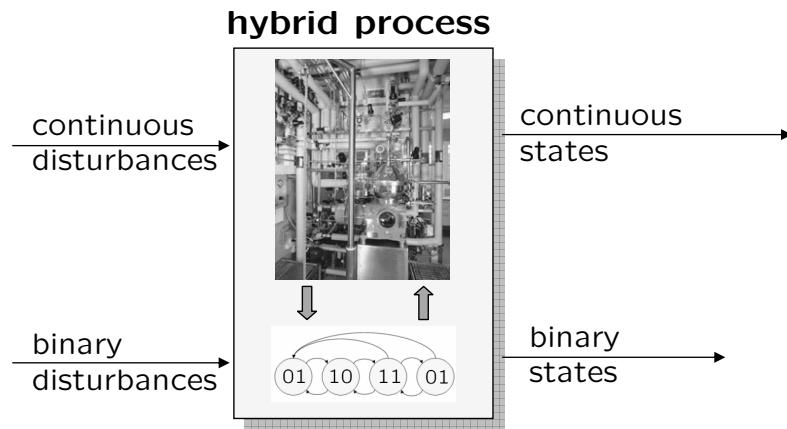
115/150

Optimization-based Reachability Analysis

(Verification of Safety Properties
via Mixed-integer Programming)

116/150

Hybrid Verification Problem



117/150

Verification Algorithm

- **QUERY:** Is the target set X_f reachable after N steps from some initial state $x_0 \in X_0$ for some input profile $u \in U$?
- Computation: Solve the mixed-integer linear program (MILP)

$$\begin{aligned} \min \quad & 0 \\ \text{s.t.} \quad & \left\{ \begin{array}{l} x(k+1) = Ax(k) + B_1u(k) + B_2\delta(k) + B_3z(k) \\ E_2\delta(k) + E_3z(k) \leq E_1u(k) + E_4x(k) + E_5 \\ Suu(k) \leq T_u \quad (u(k) \in U) \\ k = 0, 1, \dots, N-1 \\ S_0x(0) \leq T_0 \quad (x(0) \in X_0) \\ S_fx(N) \leq T_f \quad (x(N) \in X_f) \end{array} \right. \end{aligned}$$

with respect to $u(0), \delta(0), z(0), \dots, u(N-1), \delta(N-1), z(N-1), x(0)$

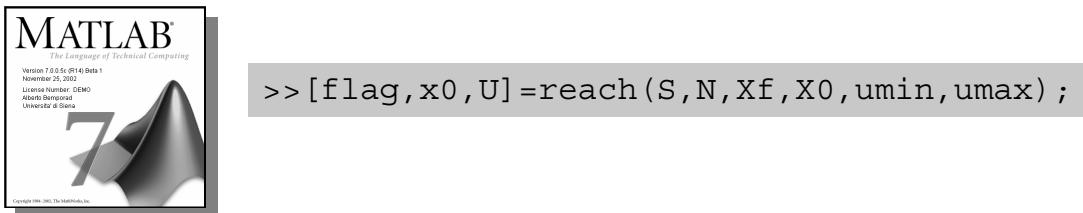
- Alternative solutions:

- Exploit the special structure of the problem and use polyhedral computation. (Torrisi, 2003)
- Use abstractions (LPs) + SAT solvers (Giorgetti, Pappas, Bemporad, 2005)

118/150

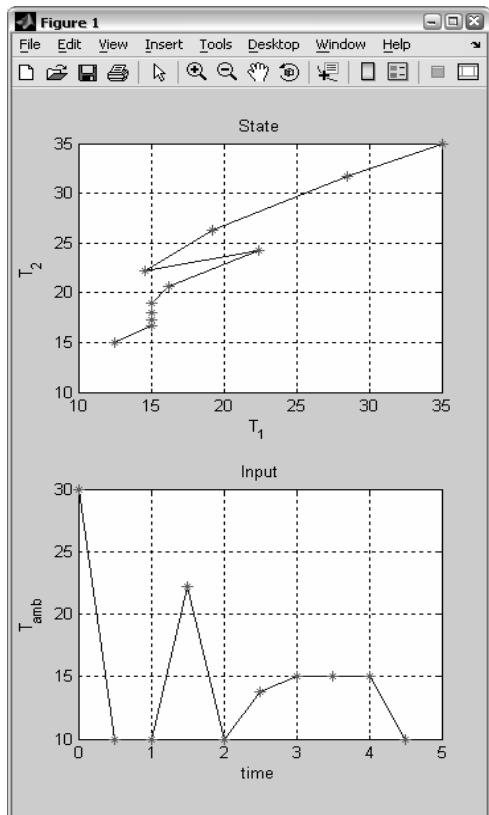
Verification Example 1

- MLD model: room temperature system
- $X_f = \left\{ \begin{bmatrix} T_1 \\ T_2 \end{bmatrix} : 10 \leq T_1, T_2 \leq 15 \right\}$ (set of unsafe states)
- $X_0 = \left\{ \begin{bmatrix} T_1 \\ T_2 \end{bmatrix} : 35 \leq T_1, T_2 \leq 40 \right\}$ (set of initial states)
- $U = \{T_{\text{amb}} : 10 \leq T_{\text{amb}} \leq 30\}$ (set of possible inputs)
- $N=10$ (time horizon)



119/150

Verification Example 1



$$U = \{T_{\text{amb}} : 10 \leq T_{\text{amb}} \leq 30\}$$

The MATLAB command window displays the following text:

```
No MIP objective value available. Exiting...
```

At the bottom, a message is circled in red:

Xf is not reachable from X0

$$U = \{T_{\text{amb}} : 20 \leq T_{\text{amb}} \leq 30\}$$

120/150

Verification Algorithm (2)

- **QUERY:** Is the target set X_f reachable **within** N steps from some initial state $x_0 \in X_0$ for some input profile $u \in U$?
- IDEA: Augment the MLD system to register the entrance of the target (unsafe) set $X_f = \{x : A_f x \leq b_f\}$:
 - Add a new variable $\delta_f(k)$, where $[\delta_f(k)=1] \rightarrow [x(k+1) \in X_f]$
 - $\implies A_f(Ax(k) + B_1 u(k) + B_2 \delta(k) + B_3 z(k)) \leq b_f + M(1 - \delta_f(k))$
big-M
 - Add the constraint $\sum_{k=0}^{N-1} \delta_f(k) \geq 1 \quad (x \in X_f \text{ for at least one } k)$
 - Solve MILP feasibility test

121/150

Verification Example 3

- States/inputs: $x_1, x_2, x_3 \in \mathbb{R}, x_4, x_5 \in \{0, 1\}$
 $u_1, u_2 \in \mathbb{R}, u_3 \in \{0, 1\}$
- Events: $[\delta_1 = 1] \leftrightarrow [x_1 \leq 0]$
 $[\delta_2 = 1] \leftrightarrow [x_2 \geq 1]$
 $[\delta_3 = 1] \leftrightarrow [x_3 - x_2 \leq 1]$
- Switched dynamics: $x_1(k+1) = \begin{cases} 0.1x_1 + 0.5x_2 & \text{if } \delta_1 \& \delta_2 | x_4 \\ -0.3x_3 - x_1 + u_1 & \text{otherwise} \end{cases}$
 $x_2(k+1) = \begin{cases} -0.8x_1 + 0.7x_3 - u_1 - u_2 & \text{if } \delta_3 | x_5 \\ -0.3x_1 - 2x_2 & \text{otherwise} \end{cases}$
 $x_3(k+1) = \begin{cases} -0.1x_3 + u_2 & \text{if } \delta_3 \& x_5 | x_4 \& \delta_1 \\ x_3 - 0.5x_1 - 2u_1 & \text{otherwise} \end{cases}$
- Automaton: $x_4(k+1) = x_4 \& \delta_1$
 $x_5(k+1) = ((x_4 | x_5) \& (\delta_1 | \delta_2)) | \delta_3 \& u_3$

Query: Starting from X_0 , is it possible that $x_i(k) \in X_f$ at some $k \leq N$, under the restriction that $x_3(k) + x_2(k) \leq 0$, $\delta_1(k) | \delta_2(k) | x_5(k)$ is true, and $\sim x_4(k) | x_5(k)$ is also true $\forall k \leq N$?

go to demo /demos/hybrid/reachtest.m

122/150

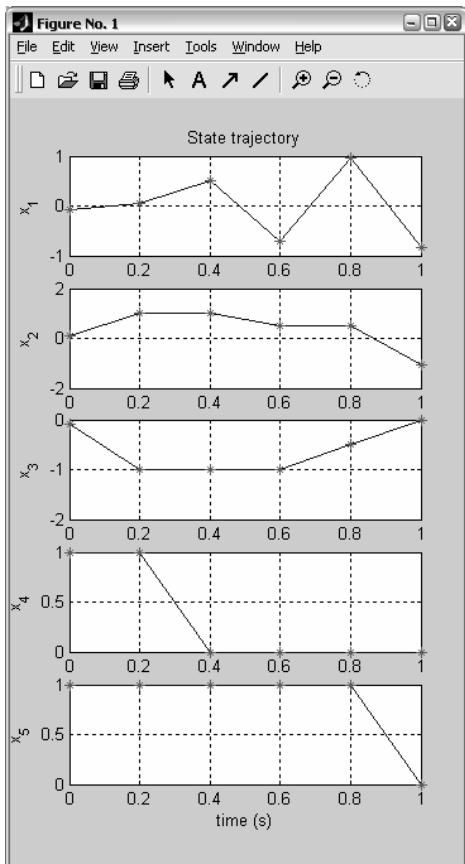
Verification Example 3

- $X_f = \left\{ \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{bmatrix} : -1 \leq x_1, x_3 \leq 1, 0.5 \leq x_2 \leq 1, x_4, x_5 \in \{0, 1\} \right\}$ (set of unsafe states)
- $X_0 = \left\{ \begin{bmatrix} x_1 \\ x_2 \\ x_3 \\ x_4 \\ x_5 \end{bmatrix} : -0.1 \leq x_1, x_3 \leq 0.1, x_2 = 0.1, x_4, x_5 \in \{0, 1\} \right\}$ (set of initial states)
- $U = \left\{ \begin{bmatrix} u_1 \\ u_2 \\ u_3 \end{bmatrix} : -1 \leq u_1 \leq 1, -2 \leq u_2 \leq 2, u_3 \in \{0, 1\} \right\}$ (set of possible inputs)
- $N=5$ (time horizon)

```
>> [flag,x0,U,xf,X,T,D,Z,Y,reachtime]=reach(S, [1 N], Xf, X0);
```

123/150

Verification Example 3



```

>> reachtest
Hybrid Toolbox v.1.0.11 [Sep 20, 2005] - (c) 200
elapsed_time =
0.2200

>> reachtime
reachtime =
2
3
4

>>

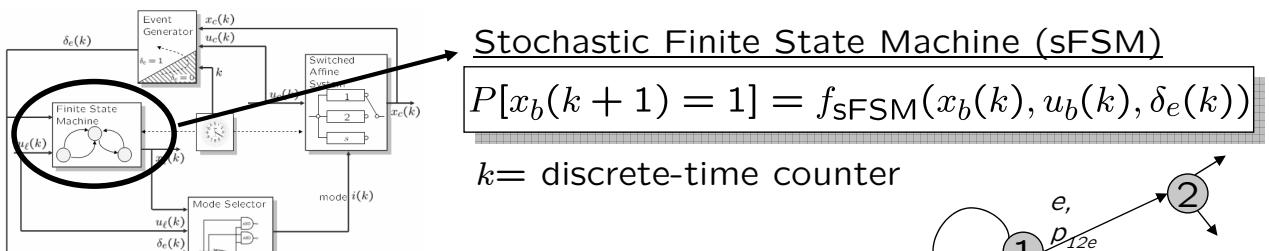
```

The set X_f is reached by $x(k)$ at times $k=2,3,4$

124/150

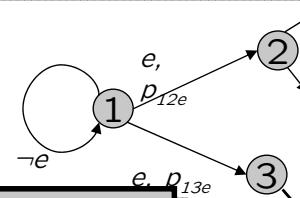
DHA - Extensions

Discrete-time Hybrid Stochastic Automaton (DHSA)

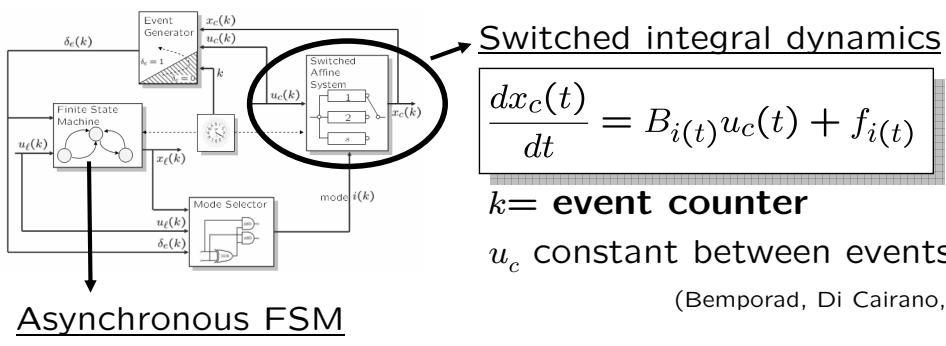


(Bemporad, Di Cairano, IEEE TAC, submitted)

All control/verification techniques developed for DHA can be extended to DHSA and icHA !



Event-based Continuous-time Hybrid Automaton (icHA)



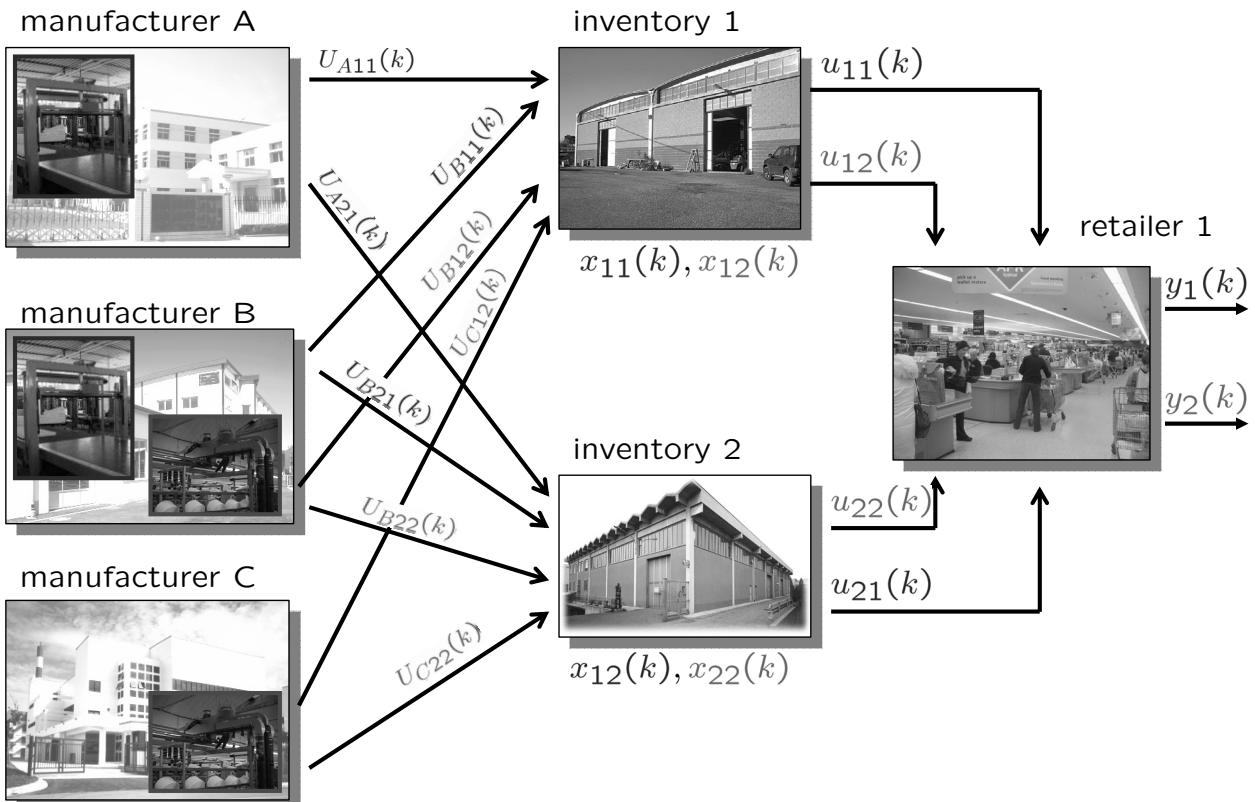
Asynchronous FSM

(Bemporad, Di Cairano, Julvez, CDC05 & HSCC-06)
125/150

Outline

- ✓ Basics of Model Predictive Control (MPC)
- ✓ Hybrid models for MPC
- ✓ MPC of hybrid systems
- ✓ Explicit MPC (multiparametric programming)
- ✓ Optimization-based reachability analysis
- Examples

A Simple Example in Supply Chain Management



127/150

System Variables

- continuous states:

$x_{ij}(k)$ = amount of j hold in inventory i at time k ($i=1,2$, $j=1,2$)

- continuous outputs:

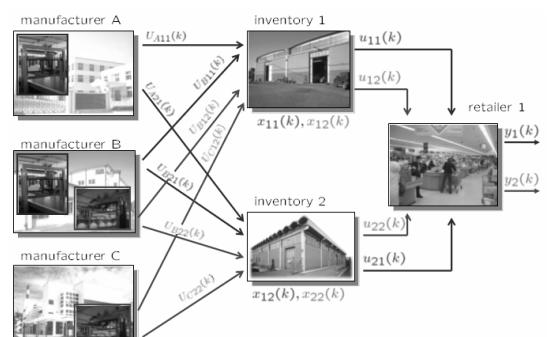
$y_j(k)$ = amount of j sold at time k ($i=1,2$)

- continuous inputs:

$u_{ij}(k)$ = amount of j taken from inventory i at time k ($i=1,2$, $j=1,2$)

- binary inputs:

$U_{Xij}(k) = 1$ if manufacturer X produces and send j to inventory i at time k



128/150

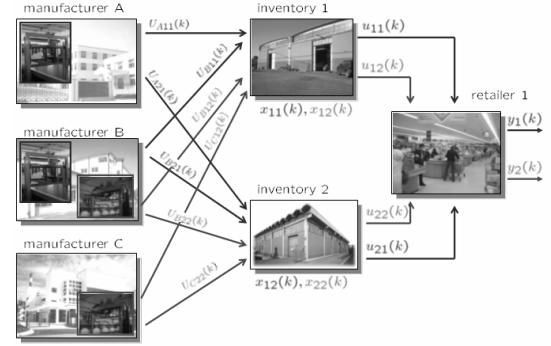
Constraints

- Max capacity of inventory i :

$$0 \leq \sum_j x_{ij}(k) \leq x_{Mi} \quad \text{Numerical values: } x_{M1}=10, x_{M2}=10$$

- Max transportation from inventories:

$$0 \leq u_{ij}(k) \leq u_M$$



- A product can only be sent to one inventory:

$UA11(k)$ and $UA21(k)$ cannot be =1 at the same time
 $UB11(k)$ and $UB21(k)$ cannot be =1 at the same time
 $UB12(k)$ and $UB22(k)$ cannot be =1 at the same time
 $UC12(k)$ and $UC22(k)$ cannot be =1 at the same time

- A manufacturer can only produce one type of product at one time:
 $[UB11(k)=1 \text{ or } UB21(k)=1]$ and $[UB12(k)=1 \text{ or } UB22(k)=1]$
cannot be true

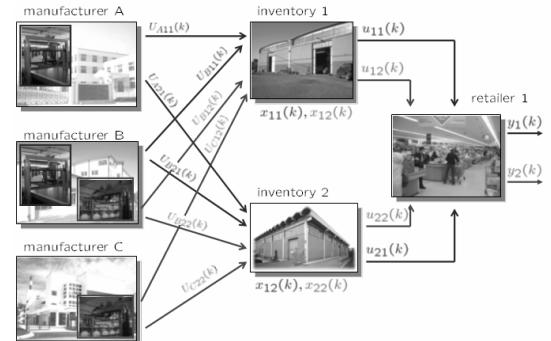
129/150

Dynamics

$P_{A1}, P_{B1}, P_{B2}, P_{C2}$ = amount of type 1(2) produced by A (B, C) in one time interval

Numerical values:

$$P_{A1}=4, P_{B1}=6, P_{B2}=7, P_{C2}=3$$



- Level of inventories:

$$\begin{cases} x_{11}(k+1) = x_{11}(k) + P_{A1}U_{A11}(k) + P_{B1}U_{B11}(k) - u_{11}(k) \\ x_{12}(k+1) = x_{12}(k) + P_{B2}U_{B12}(k) + P_{C2}U_{C12}(k) - u_{12}(k) \\ x_{21}(k+1) = x_{21}(k) + P_{A1}U_{A21}(k) + P_{B1}U_{B21}(k) - u_{21}(k) \\ x_{22}(k+1) = x_{22}(k) + P_{B2}U_{B22}(k) + P_{C2}U_{C22}(k) - u_{22}(k) \end{cases}$$

130/150

Hybrid Dynamical Model

```

SYSTEM supply_chain{
INTERFACE {
    STATE { REAL x11 [0,10];
             REAL x12 [0,10];
             REAL x21 [0,10];
             REAL x22 [0,10]; }

    INPUT { REAL u11 [0,10];
             REAL u12 [0,10];
             REAL u21 [0,10];
             REAL u22 [0,10];
             BOOL UA11,UA21,UB11,UB21,UB21,UC12,UC22; }

    OUTPUT {REAL y1,y2; }

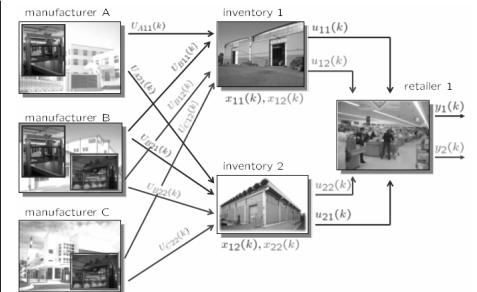
    PARAMETER { REAL PA1,PB1,PB2,PC2,xM1,xM2; }
}
IMPLEMENTATION {

    AUX { REAL zA11, zB11, zB12, zC12, zA21, zB21, zB22, zC22; }

    DA { zA11 = {IF UA11 THEN PA1 ELSE 0};
        zB11 = {IF UB11 THEN PB1 ELSE 0};
        zB12 = {IF UB12 THEN PB2 ELSE 0};
        zC12 = {IF UC12 THEN PC2 ELSE 0};
        zA21 = {IF UA21 THEN PA1 ELSE 0};
        zB21 = {IF UB21 THEN PB1 ELSE 0};
        zB22 = {IF UB22 THEN PB2 ELSE 0};
        zC22 = {IF UC22 THEN PC2 ELSE 0}; }
}

```

/demos/hybrid/supply_chain.m



```

CONTINUOUS {x11 = x11 + zA11 + zB11 - u11;
            x12 = x12 + zB12 + zC12 - u12;
            x21 = x21 + zA21 + zB21 - u21;
            x22 = x22 + zB22 + zC22 - u22; }

OUTPUT {      y1 = u11 + u21;
            y2 = u12 + u22; }

MUST {      ~(UA11 & UA21);
            ~(UC12 & UC22);
            ~((UB11 | UB21) & (UB12 | UB22));
            ~((UB11 & UB21);
            ~((UB12 & UB22));
            x11+x12 <= xM1;
            x11+x12 >=0;
            x21+x22 <= xM2;
            x21+x22 >=0; }
} }

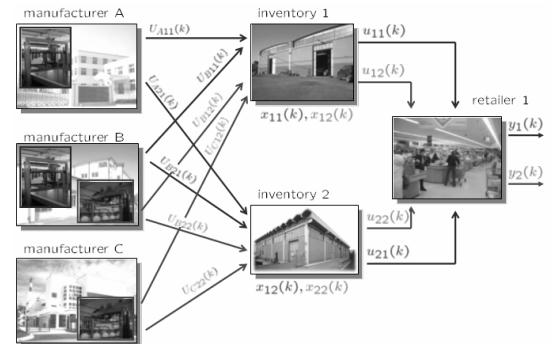
```

1/150

Objectives

- Meet customer demand as much as possible:

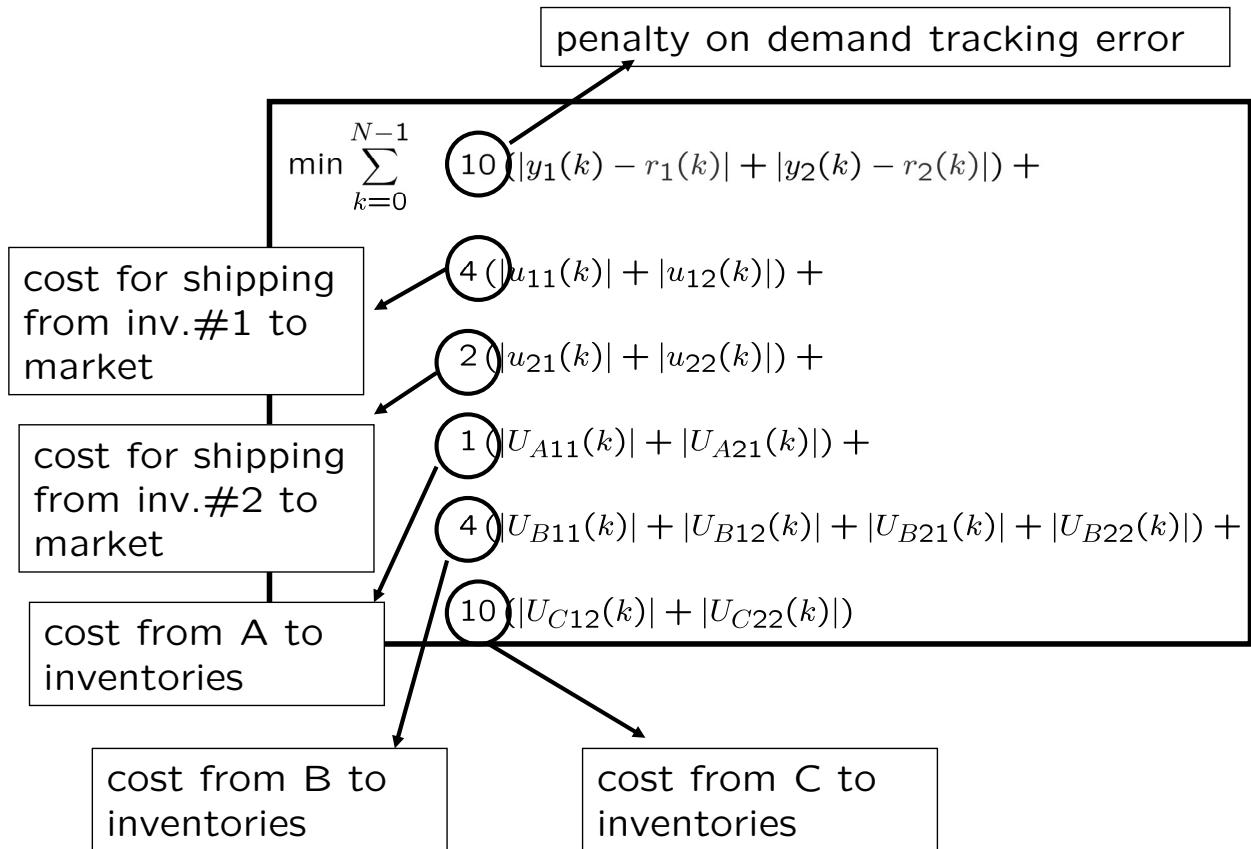
$$y_1 \approx r_1, y_2 \approx r_2$$



- Minimize transportation costs

- Fulfill all constraints

Performance Specs

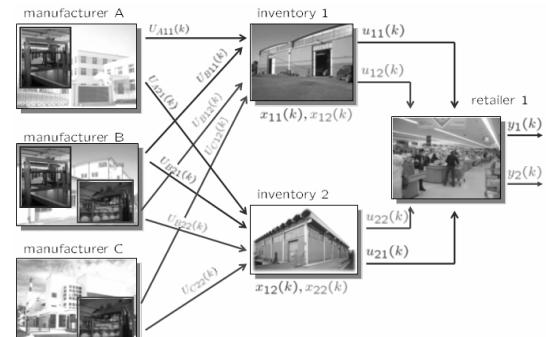


133/150

Hybrid MPC - Example

```
>>refs.y=[1 2]; % weights output2 #1,#2
>>Q.y=diag([10 10]); % output weights
...
>>Q.norm=Inf; % infinity norms
>>N=2; % optimization horizon
>>limits.umin=umin; % constraints
>>limits.umax=umax;
>>limits.xmin=xmin;
>>limits.xmax=xmax;

>>C=hybcon(S,Q,N,limits,refs);
```



```
>> C

Hybrid controller based on MLD model S <supply_chain.hys> [Inf-norm]

4 state measurement(s)
2 output reference(s)
12 input reference(s)
0 state reference(s)
0 reference(s) on auxiliary continuous z-variables

44 optimization variable(s) (28 continuous, 16 binary)
176 mixed-integer linear inequalities
sampling time = 1, MILP solver = 'glpk'

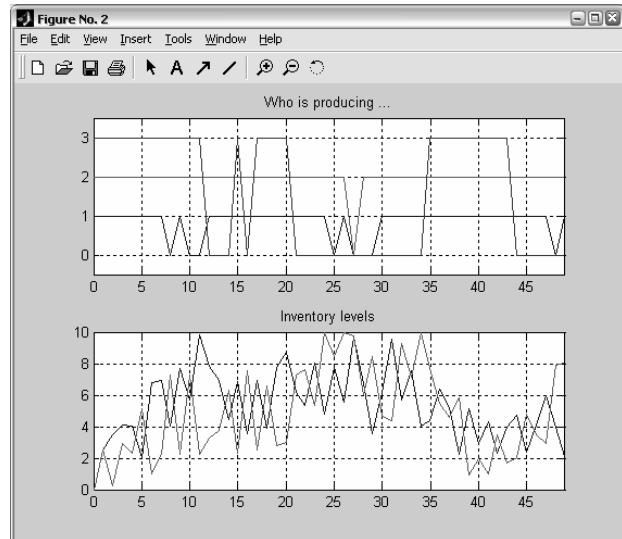
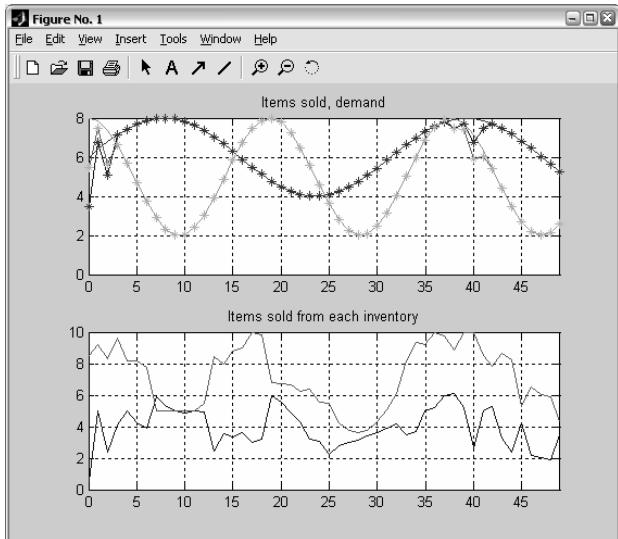
Type "struct(C)" for more details.

>>
```

134/150

Hybrid MPC - Example

```
>>x0=[0;0;0;0]; % Initial condition  
>>r.y=[6+2*sin((0:Tstop-1)'/5); % Reference trajectories  
      5+3*cos((0:Tstop-1)'/3)];  
  
>> [XX,UU,DD,ZZ,TT]=sim(C,S,r,x0,Tstop);
```



CPU time: $\approx 30\text{ms}$ per time step (using GLPK on this machine)

135/150



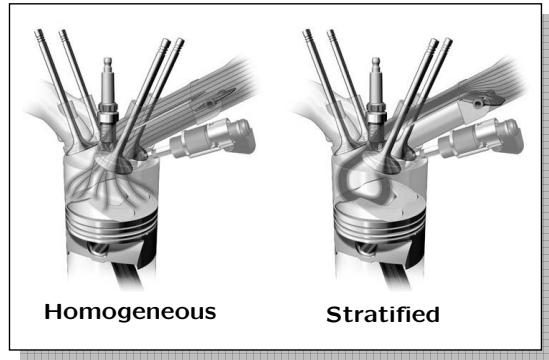
Hybrid Control of a DISC Engine

(Giorgetti, Ripaccioli, Bemporad, Kolmanovsky, Hrovat, IEEE Tr. Mechatronics, 2006)

DISC Engine

Two distinct regimes:

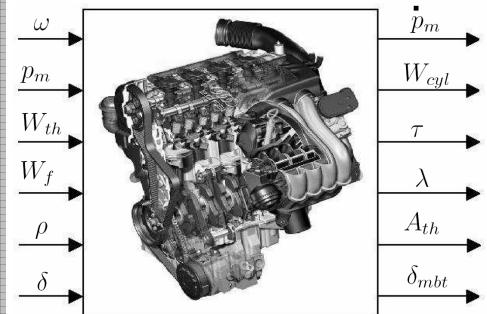
Regime	fuel injection	air-to-fuel ratio
Homogeneous combustion	intake stroke	$\lambda = 14.64$
Stratified combustion	compression stroke	$\lambda > 14.64$



Pro: reduce consumption up to 15%

Con: complex treatment of exhaust gas

- **States:** intake manifold pressure (p_m)
- **Outputs:** Air-to-fuel ratio (λ), torque (τ), max-brake-torque spark timing (δ_{mbt})
- **Inputs:** spark advance (δ), air flow (W_{th}), fuel flow (W_f), combustion regime (ρ);
- **Disturbance:** engine speed (ω) [measured]



137/150

DISC Engine – Control Objective

Objective: Design a controller for the engine that

- Automatically chooses operating **mode** (homogeneous/stratified)
- Can cope with **nonlinear** dynamics
- Handles **constraints** (on A/F ratio, air-flow, spark)
- Achieves **optimal** performance (tracking of desired torque and A/F ratio)

138/150

DISC Engine - HYSDEL List

```

SYSTEM hysdisc{
  INTERFACE{
    STATE{
      REAL pm      [1, 101.325];
      REAL xtau     [-1e3, 1e3];
      REAL xlam     [-1e3, 1e3];
      REAL taud     [0, 100];
      REAL lamd     [10, 60];
    }
    OUTPUT{
      REAL lambda, tau, ddelta;
    }
    INPUT{
      REAL Wth      [0, 38.5218];
      REAL Wf       [0, 2];
      REAL delta    [0, 40];
      BOOL rho;
    }
    PARAMETER{
      REAL Ts, pm1, pm2;
      ...
    }
  }

  IMPLEMENTATION{
    AUX{
      REAL lam,taul,dmbtl,lmin,lmax;
    }
    DA{
      lam={IF rho THEN l11*pm+l12*Wth...
            +l13*Wf+l14*delta+l1c
        ELSE   101*pm+102*Wth+103*Wf...
            +104*delta+10c    };
      taul={IF rho THEN tau11*pm+...
            tau12*Wth+tau13*Wf+tau14*delta+tau1c
        ELSE   tau01*pm+tau02*Wth...
            +tau03*Wf+tau04*delta+tau0c  };
      dmbtl ={IF rho THEN dmbt11*pm+dmbt12*Wth...
            +dmbt13*Wf+dmbt14*delta+dmbt1c+7
        ELSE dmbt01*pm+dmbt02*Wth...
            +dmbt03*Wf+dmbt04*delta+dmbt0c-1};
      lmin ={IF rho THEN 13 ELSE 19};
      lmax ={IF rho THEN 21 ELSE 38};
    }
    CONTINUOUS{
      pm=pm1*pm+pm2*Wth;
      xtau=xtau+Ts*(taud-taul);
      xlam=xlam+Ts*(lamd-lam);
      taud=taud; lamd=lamd;
    }
    OUTPUT{
      lambda=lam-lamd;
      tau=taul-taud;
      ddelta=dmbtl-delta;
    }
    MUST{
      lmin-lam      <=0;
      lam-lmax      <=0;
      delta-dmbtl  <=0;
    }
  }
}

```

139/150

MPC of DISC Engine

$$\min \sum_{k=0}^{N-1} (u_k - u_r)' R (u_k - u_r) + (y_k - y_r)' Q (y_k - y_r) \\ + (x_{k+1} - x_r)' S (x_{k+1} - x_r)$$

subj. to $\begin{cases} x_0 = x(t), \\ \text{hybrid model} \end{cases}$

$$u(t) = [W_{th}(t), W_f(t), \delta(t), \rho(t)]$$

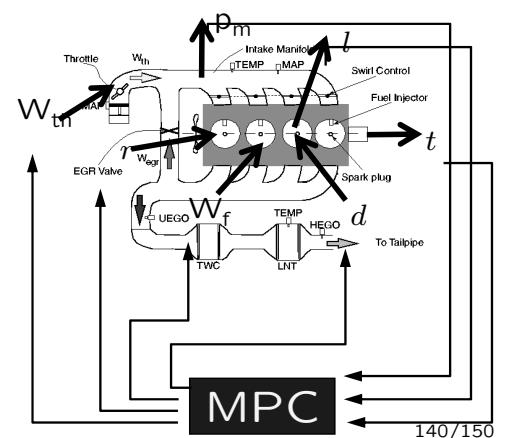
Weights:

$$R = \begin{pmatrix} 0.01 & 0 & 0 & 0 \\ 0 & 0.001 & 0 & 0 \\ 0 & 0 & 0 & 0 \\ 0 & 0 & 0 & 1 \end{pmatrix} \quad r_p \quad (\text{prevents unneeded chattering})$$

$$Q = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 0.001 & 0 \\ 0 & 0 & 0.01 \end{pmatrix} \quad S = \begin{pmatrix} 0.04 & 0 & 0 \\ 0 & 1500 & 0 \\ 0 & 0 & 0.01 \end{pmatrix} \quad s_{\varepsilon_\tau} \quad s_{\varepsilon_\lambda}$$

main emphasis on torque

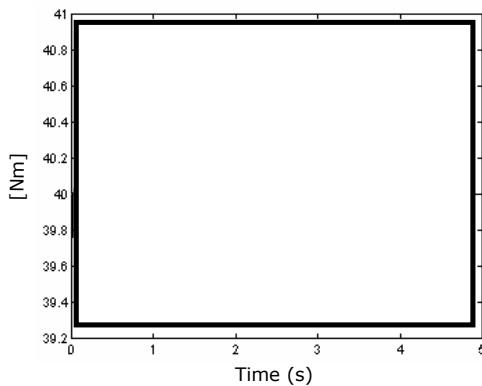
Solve
MIQP problem
(mixed-integer quadratic program)
to compute $u(t)$



140/150

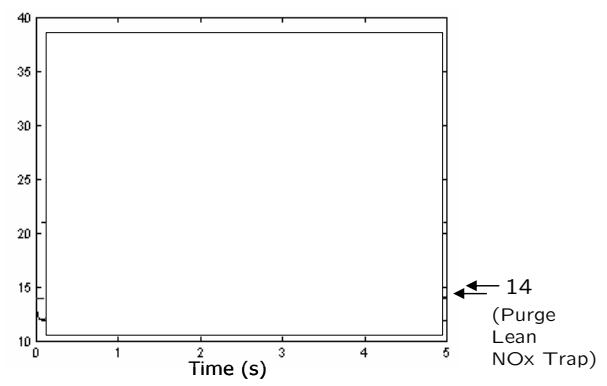
Simulation Results (nominal engine speed)

Engine Brake Torque

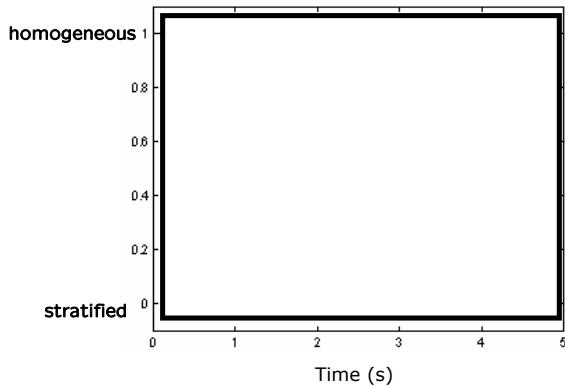


$\omega = 2000$ rpm

Air-to-Fuel Ratio



Combustion mode

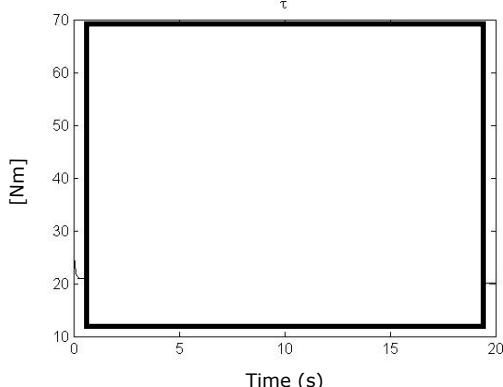


- Control horizon $N=1$;
 - Sampling time $T_s=10$ ms;
 - PC Xeon 2.8 GHz + Cplex 9.1
- \downarrow
- ≈ 3 ms per time step

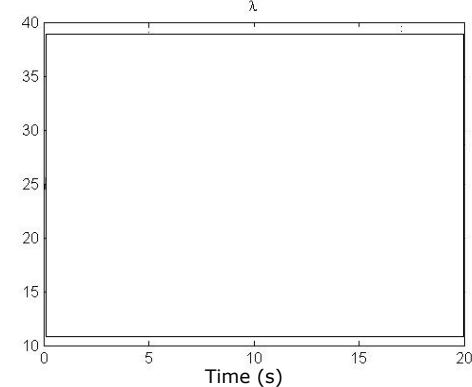
141/150

Simulation Results (varying engine speed)

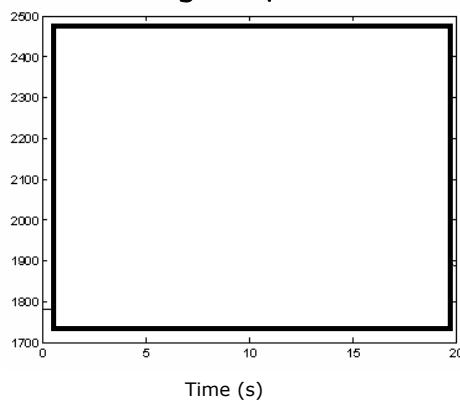
Engine Brake Torque



Air-to-Fuel Ratio



Engine speed

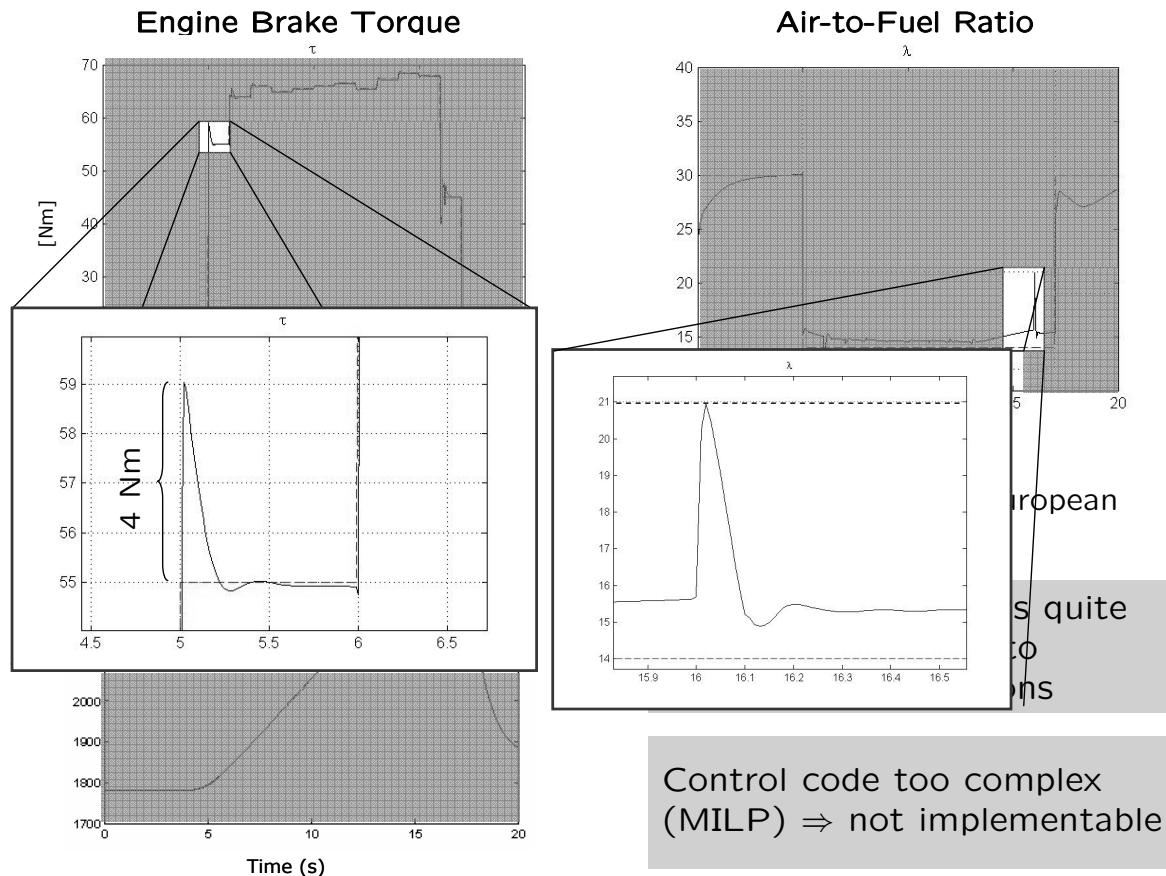


20 s segment of the European drive cycle (NEDC)

Hybrid MPC design is quite robust with respect to engine speed variations

142/150

Simulation Results (varying engine speed)



143/150

Explicit MPC Controller

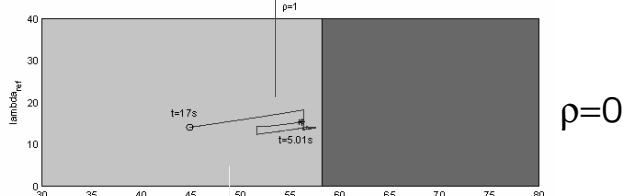
Explicit control law:
$$u(t) = f(\theta(t))$$

where: $u = [W_{th} \ W_f \ \delta \ \rho]'$ ➡
 $\theta = [p_m \ \epsilon_\tau \ \epsilon_\lambda \ \tau_{ref} \ \lambda_{ref}]'$
 $p_{m,ref} \ W_{th,ref} \ W_{f,ref} \ \delta_{ref}$

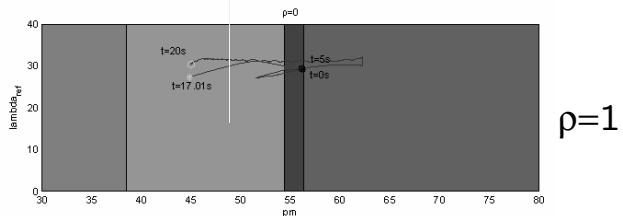
$N=1$ (control horizon)

42 partitions

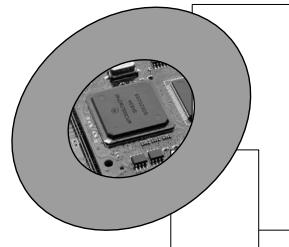
Cross-section by the $\tau_{ref}-\lambda_{ref}$ plane



- Time to compute explicit MPC: $\approx 3s$;
 - Sampling time $T_s=10$ ms;
 - PC Xeon 2.8 GHz + Cplex 9.1
- $\approx 8 \mu s$ per time step



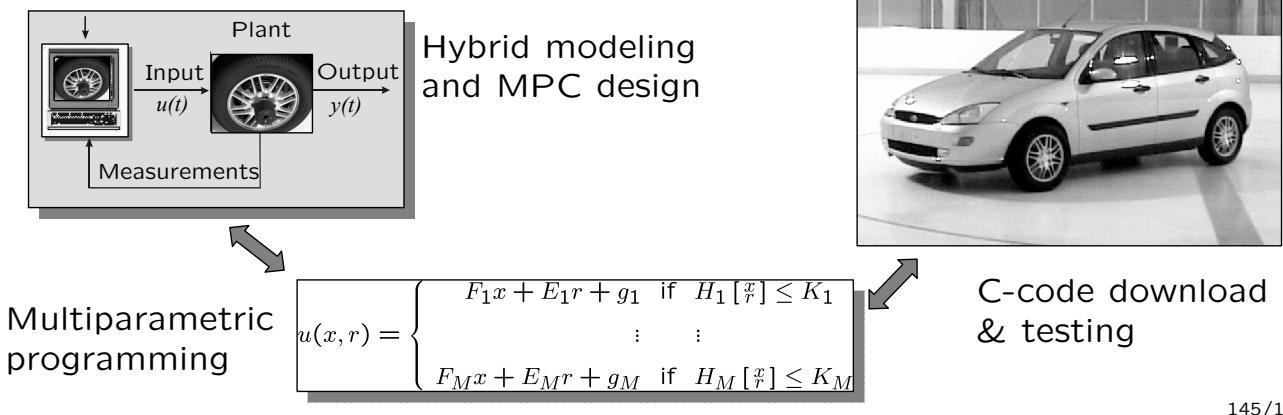
$\approx 3ms$ on
μcontroller
Motorola
MPC 555
43kb RAM
(custom made for Ford)



144/150

Conclusions

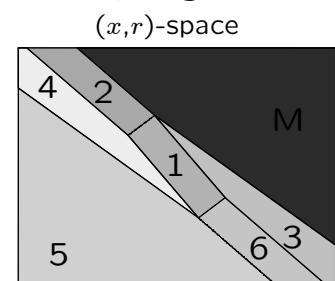
- **Hybrid systems** as a framework for new applications, where both logic and continuous dynamics are relevant
- **Supervisory MPC controllers** schemes can be synthesized via on-line mixed-integer programming (MILP/MIQP)
- **Piecewise Linear MPC Controllers** can be synthesized off-line via multiparametric programming for fast-sampling applications



Conclusions

- **Hybrid systems** as a framework for new applications, where both logic and continuous dynamics are relevant
- **Supervisory MPC controllers** schemes can be synthesized via on-line mixed-integer programming (MILP/MIQP)
- **Piecewise Linear MPC Controllers** can be synthesized off-line via multiparametric programming for fast-sampling applications

$$u(x, r) = \begin{cases} F_1x + E_1r + g_1 & \text{if } H_1[\vec{x}] \leq K_1 \\ \vdots & \vdots \\ F_Mx + E_Mr + g_M & \text{if } H_M[\vec{x}] \leq K_M \end{cases}$$



- **Matlab tools** available to assist the whole design process (models, simulation, MPC design, code generation):
 - MPC Toolbox (linear), Hybrid Toolbox (hybrid, explicit), Multi-Parametric Toolbox (PWA, explicit)

References

General References

- [1] J.M. Maciejowski, *Predictive Control with Constraints*, Prentice Hall, Harlow, UK, 2002.
- [2] E.F. Camacho and C. Bordons, *Model Predictive Control*, Advanced Textbooks in Control and Signal Processing. Springer-Verlag, London, 2nd edition, 2004.
- [3] A. Bemporad, M. Morari, and N. L. Ricker, *Model Predictive Control Toolbox for Matlab – User’s Guide*, The Mathworks, Inc., 2004,
<http://www.mathworks.com/access/helpdesk/help/toolbox/mpc/>.
- [4] A. Bemporad, *Hybrid Toolbox – User’s Guide*, Jan. 2004,
<http://www.dii.unisi.it/hybrid/toolbox>.
- [5] D.Q. Mayne, J.B. Rawlings, C.V. Rao, and P.O.M. Scokaert, “Constrained model predictive control: Stability and optimality,” *Automatica*, vol. 36, no. 6, pp. 789-814, June 2000.
- [6] J.B. Rawlings, “Tutorial overview of model predictive control,” *IEEE Control Systems Magazine*, pp. 38-52, June 2000.
- [7] A. Bemporad, “Model-based predictive control design: New trends and tools,” in *Proc. 45th IEEE Conf. on Decision and Control*, San Diego, CA, 2006.

147/150

References

MLD and HYSDEL Modeling

- [1] A. Bemporad, “Hybrid Toolbox – User’s Guide,” Dec. 2003,
<http://www.dii.unisi.it/hybrid/toolbox>
- [2] F.D. Torrisi and A. Bemporad, “HYSDEL - A tool for generating computational hybrid models,” *IEEE Transactions on Control Systems Technology*, vol. 12, no. 2, pp. 235-249, Mar. 2004
- [3] A. Bemporad and M. Morari, “Control of systems integrating logic, dynamics, and constraints,” *Automatica*, vol. 35, no. 3, pp. 407-427, Mar. 1999.
- [4] A. Bemporad, “Efficient conversion of mixed logical dynamical systems into an equivalent piecewise affine form,” *IEEE Trans. Automatic Control*, vol. 49, no. 5, pp. 832-838, 2004.
- [5] A. Bemporad, G. Ferrari-Trecate, and M. Morari, “Observability and controllability of piecewise affine and hybrid systems,” *IEEE Trans. Automatic Control*, vol. 45, no. 10, pp. 1864-1876, 2000.
- [6] W.P.H.M Heemels, B. de Schutter, and A. Bemporad, “Equivalence of hybrid dynamical models,” *Automatica*, vol. 37, no. 7, pp. 1085-1091, July 2001
- [7] A. Bemporad, W.P.M.H. Heemels, and B. De Schutter, “On hybrid systems and closed-loop MPC systems,” *IEEE Trans. Automatic Control*, vol. 47, no. 5, pp. 863-869, May 2002.

148/150

References

- [8] A. Bemporad, A. Garulli, S. Paoletti, and A. Vicino, "A bounded-error approach to piecewise affine system identification," *IEEE Trans. Automatic Control*, vol. 50, no. 10, pp. 1567-1580, Oct. 2005 .
- [9] J. Roll, A. Bemporad, and L. Ljung, "Identification of piecewise affine systems via mixed-integer programming," *Automatica*, vol. 40, no. 1, pp. 37-50, 2004
- [10] G. Ferrari-Trecate, M. Muselli, D. Liberati, and M. Morari, "A clustering technique for the identification of piecewise affine systems," *Automatica*, vol. 39, no. 2, pp. 205-217, Feb. 2003.

Model predictive control

- [11] F. Borrelli, M. Baotic, A. Bemporad, and M. Morari, "Dynamic programming for constrained optimal control of discrete-time linear hybrid systems," *Automatica*, vol. 41, no. 10, Oct. 2005
- [12] A. Bemporad, M. Morari, V. Dua, and E.N. Pistikopoulos, "The explicit linear quadratic regulator for constrained systems," *Automatica*, vol. 38, no. 1, pp. 3-20, 2002.
- [13] A. Bemporad, F. Borrelli, and M. Morari, "Piecewise linear optimal controllers for hybrid systems," In *Proc. American Control Conference*, 2000, pp. 1190-1194.
- [14] A. Bemporad and N. Giorgetti, "Logic-based methods for optimal control of hybrid systems," *IEEE Trans. Autom. Control*, vol. 51, no. 6, pp. 963-976, 2006.
- [15] M. Lazar, M. Heemels, S. Weiland, A. Bemporad, "Stability of Hybrid Model Predictive Control," *IEEE Trans. Automatic Control*, vol. 51, no. 11, pp. 1813-1818, 2006.

149/150

Reachability and observability

References

- [16] A. Bemporad, G. Ferrari-Trecate, and M. Morari, "Observability and controllability of piecewise affine and hybrid systems," *IEEE TAC*, vol. 45, no. 10, pp. 1864-1876, 2000.
- [17] A. Bemporad, D. Mignone, and M. Morari, "Moving horizon estimation for hybrid systems and fault detection," in *Proc. American Control Conf.*, 1999, Chicago, IL, pp. 2471-2475.
- [18] G. Ferrari-Trecate, D. Mignone, and M. Morari, "Moving horizon estimation for hybrid systems," *IEEE TAC*, vol. 47, no. 10, pp. 1663-1676, 2002.

Applications

- [19] F. Borrelli, A. Bemporad, M. Fodor, and D. Hrovat, "An MPC/Hybrid System Approach to Traction Control," *IEEE Control Syst. Tech.*, vol. 14, n. 3, pp. 541-552, 2006.
- [20] N. Giorgetti, G. Ripaccioli, A. Bemporad, I.V. Kolmanovsky, and D. Hrovat, "Hybrid model predictive control of direct injection stratified charge engines," *IEEE/ASME Transactions on Mechatronics*, vol. 11, no. 5, pp. 499-506, Aug. 2006.
- [21] N. Giorgetti, A. Bemporad, H. E. Tseng, and D. Hrovat, "Hybrid model predictive control application towards optimal semi-active suspension," *International Journal of Control*, vol. 79, no. 5, pp. 521-533, 2006.
- [22] A. Bemporad, P. Borodani, and M. Mannelli, "Hybrid control of an automotive robotized gearbox for reduction of consumptions and emissions," LNCS 2623, pp. 81-96, Springer-Verlag, 2003.

more on <http://www.dii.unisi.it/~bemporad/publications>

150/150