

### 1<sup>st</sup> HYCON PhD School on Hybrid Systems



# Solution Concepts and Well-posedness of Hybrid Systems

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show

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Typrid systems combine continuous dynamic (differential or difference equations) typical of physical plants and discrete dynamics (automata and logical conditions) typical of control logic. By combining disciplines of computer science and systems and control theor research on hybrid systems provide a solid theory and computational tools for the analysis, simulation, verification, and control design of ded systems", and are used in a large ariety of applications (automotive systems, air traffic management biological systems process industries, and many others

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nocht a conportagen fin Siena, July 19-22, 2005 - Rectorate of the University of Siena



### Solution Concepts and Well-posedness of Hybrid Systems

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HYCON Summer School on Hybrid Systems



#### Key issues:

- Solution concepts
- Well-posedness: existence & uniqueness of solutions given an initial condition

#### **Outline lecture**

- Smooth systems: differential equations
- Switched systems: Discontinuous differential equations: "classics"
- Hybrid automata
- Zenoness: importance of choice of solution concept
- Some piecewise linear, linear relay and complementarity systems
- Summary



#### Solution concept

#### Description format / syntax / model ↓ solutions / trajectories / executions/ semantics/ behavior



**Well-posedness:** given initial condition does there **exists** a solution and is it **unique**?

Let's start simple ...



#### Smooth differential equations

**Example**  $\dot{x} = f(t, x)$   $x(t_0) = x_0$ .

A solution trajectory is a function  $x : [t_0, t_1] \mapsto \mathbb{R}^n$  that is continuous, differentiable and satisfies  $x(t_0) = x_0$  and

$$\dot{x}(t) = f(t, x(t))$$
 for all  $t \in (t_0, t_1)$ 

**Well-posedness:** given initial condition does there **exists** a solution and is it **unique**?



#### Well-posedness

**Example** 
$$\dot{x} = 2\sqrt{x}$$
,  $x(0) = 0$ . Solutions:  $x(t) = 0$  and  $x(t) = t^2$ .

Local existence and uniqueness of solutions given an initial condition:

**Theorem 1** Let f(t, x) be piecewise continuous in t and satisfy the following Lipschitz condition: there exist an L > 0 and r > 0 such that

$$||f(t,x) - f(t,y)|| \le L||x - y||$$

and all x and y in a neighborhood  $B := \{x \in \mathbb{R}^n \mid ||x - x_0|| < r\}$  of  $x_0$ and for all  $t \in [t_0, t_1]$ .

 $\Downarrow$ 

There is a  $\delta > 0$  s.t. a unique solution exists on  $[t_0, t_0 + \delta]$  starting in  $x_0$  at  $t_0$ .



#### Global well-posedness

**Example**  $\dot{x} = x^2 + 1$ , x(0) = 0. Solution:  $x(t) = \tan t$ . Local on  $[0, \pi/2)$ .

• Note that we have  $\lim_{t\uparrow \pi/2} x(t) = \infty$ . Finite escape time!

**Theorem 2 (Global Lipschitz condition)** Suppose f(t,x) is piecewise continuous in t and satisfies

$$\|f(t,x) - f(t,y)\| \le L \|x - y\|$$

for all x, y in  $\mathbb{R}^n$  and for all  $t \in [t_0, t_1]$ . Then, a unique solution exists on  $[t_0, t_1]$  for any initial state  $x_0$  at  $t_0$ .

- Not necessary:  $\dot{x} = -x^3$  not glob. Lipsch., but unique global solutions.
- Also in hybrid systems, but even more awkward stuff (Zeno)

### Discontinuous differential equations: a class of switched systems

$$\begin{array}{c} \mathbf{C}_{\star} \\ \mathbf{x}' = \mathbf{f}_{\star}(\mathbf{x}) \\ & & \\ \mathbf{x}' = \mathbf{f}_{\star}(\mathbf{x}) \end{array} \qquad \qquad \mathbf{C}_{\star} \\ & & \\ \mathbf{x}' = \mathbf{f}_{\star}(\mathbf{x}) \\ & & \\ \mathbf{x}' = \mathbf{f}_{\star}(\mathbf{x}) \\ & & \\ \mathbf{x}' = \left\{ \begin{array}{c} f_{+}(x) & , \text{ if } x \in C_{+} := \{x \in \mathbb{R}^{n} \mid \phi(x) > 0\} \\ f_{-}(x) & , \text{ if } x \in C_{-} := \{x \in \mathbb{R}^{n} \mid \phi(x) < 0\} \end{array} \right. \end{array}$$

- x in interior of  $C_{-}$  or  $C_{+}$ : just follow!
- $f_{-}(x)$  and  $f_{+}(x)$  point in same direction: just follow!

$$n(x) = \frac{\nabla \phi(x)}{\|\nabla \phi(x)\|} \text{ then } (n(x)^T f_{-}(x)) \cdot (n(x)^T f_{+}(x)) > 0$$

•  $n(x)^T f_+(x) > 0$  ( $f_+(x)$  points towards  $C_+$ ) and  $n(x)^T f_-(x) < 0$  ( $f_-(x)$  points towards  $C_-$ ): At least two trajectories

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#### Sliding modes



 $n(x)^T f_+(x) < 0$  ( $f_+(x)$  points towards  $C_-$ ) and  $n(x)^T f_-(x) > 0$  ( $f_-(x)$  points towards  $C_+$ ).

#### No classical solution

- Relaxation: spatial (hysteresis)  $\Delta$ , time delay  $\tau$ , smoothing  $\varepsilon$
- Chattering / infinitely fast switching (limit case  $\Delta \downarrow 0$ ,  $\varepsilon \downarrow 0$ , and  $\tau \downarrow 0$ )

Filippov's convex definition: convex combination of both dynamics

 $\dot{x} = \lambda f_+(x) + (1 - \lambda) f_-(x)$  with  $0 \le \lambda \le 1$ 

such that x moves ("slides") along  $\phi(x) = 0$ . "Third mode ... "



#### **Differential inclusions**

$$\dot{x} = \begin{cases} f_{+}(x), & \text{if } \phi(x) > 0\\ \lambda f_{+}(x) + (1 - \lambda) f_{-}(x), & \text{if } \phi(x) = 0, \ 0 \le \lambda \le 1\\ f_{-}(x), & \text{if } \phi(x) < 0, \end{cases}$$

**Differential inclusion**  $\dot{x} \in F(x)$  with set-valued

$$F(x) = \begin{cases} \{f_+(x)\}, & \phi(x) > 0\\ \{\lambda f_+(x) + (1-\lambda)f_-(x) \mid \lambda \in [0,1]\}, & \phi(x) = 0\\ \{f_-(x)\}, & \phi(x) < 0 \end{cases}$$

**Definition 3** A function  $x : [a, b] \mapsto \mathbb{R}^n$  is a solution of  $\dot{x} \in F(x)$ , if x is absolutely continuous and satisfies  $\dot{x}(t) \in F(x(t))$  for almost all  $t \in [a, b]$ .



#### A well-posedness result



- $f_-$  and  $f_+$  are continuously differentiable ( $C^1$ )
- $\phi$  is  $C^2$
- $\bullet$  the discontinuity vector  $h(x):=f_+(x)-f_-(x)$  is  $C^1$

If for each point x with  $\phi(x) = 0$  at least one of the two inequalities  $n(x)^T f_+(x) < 0$  or  $n(x)^T f_-(x) > 0$  (for different points a different inequality may hold), then the Filippov solutions exist and are unique.



#### Alternative: Utkin's equivalent control definition

$$\dot{x} = f(x, u) \text{ with } u = \begin{cases} g_+(x), & \xi(x) > 0\\ g_-(x), & \xi(x) < 0 \end{cases}$$

• Sliding mode:  $f_+(x):=f(x,g_+(x))$  and  $f_-(x):=f(x,g_-(x))$  point outside  $C_+$  and  $C_-$  , resp.

$$u_{\text{equiv}} \in U(x) := \begin{cases} \{g_+(x)\}, & \text{if } \xi(x) > 0\\ \{\lambda g_+(x) + (1-\lambda)g_-(x) \mid \lambda \in [0,1]\}, & \text{if } \xi(x) = 0\\ \{g_-(x)\}, & \text{if } \xi(x) < 0 \end{cases}$$

Differential inclusion

$$\dot{x}\in F(x):=f(x,U(x))=\{f(x,u)\mid u\in U(x)\}$$

"Idealization" determines Filippov/ Utkin / different solution concept!



#### Example

$$\begin{aligned} \dot{x}_1 &= -x_1 + x_2 - u \\ \dot{x}_2 &= 2x_2(u^2 - u - 1) \\ u &= \begin{cases} 1, & \text{if } x_1 > 0 \\ -1, & \text{if } x_1 < 0. \end{cases} \end{aligned}$$

Two "original" dynamics:

• 
$$C_+: x_1 > 0: \quad \dot{x} = f_+(x)$$
  
 $\dot{x}_1 = -x_1 + x_2 - 1$   
 $\dot{x}_2 = -2x_2$ 

• 
$$C_{-}: x_1 < 0: \quad \dot{x} = f_{-}(x)$$
  
 $\dot{x}_1 = -x_1 + x_2 + 1$   
 $\dot{x}_2 = 2x_2$ 



#### Vector fields



#### Vector fields: zoom







#### Sliding modes?

#### Two "original" dynamics:

- $C_+: x_1 > 0: \quad \dot{x} = f_+(x)$   $\dot{x}_1 = -x_1 + x_2 - 1$   $\dot{x}_2 = -2x_2$ •  $C_-: x_1 < 0: \quad \dot{x} = f_-(x)$   $\dot{x}_1 = -x_1 + x_2 + 1$   $\dot{x}_2 = 2x_2$ 
  - $n(x)^T f_+(x) = x_2 1 < 0 \longrightarrow x_2 < 1$
  - $n(x)^T f_-(x) = x_2 + 1 > 0 \longrightarrow x_2 > -1$
  - Sliding possible in  $x_1 = 0$  and  $x_2 \in [-1, 1]$ .



#### Filippov's solution concept

Two "original" dynamics:

•  $C_+: x_1 > 0: \quad \dot{x} = f_+(x)$   $\dot{x}_1 = -x_1 + x_2 - 1$   $\dot{x}_2 = -2x_2$ •  $C_-: x_1 < 0: \quad \dot{x} = f_-(x)$   $\dot{x}_1 = -x_1 + x_2 + 1$  $\dot{x}_2 = 2x_2$ 

- Filippov: Take convex combination of dynamics such that state slides on  $x_1 = 0$ : Hence,  $x_1 = \dot{x}_1 = 0$ .
- $\lambda(x_2 1) + (1 \lambda)(x_2 + 1) = 0$  implies  $\lambda = \frac{1}{2}(x_2 + 1)$
- Hence,  $\dot{x_2} = \lambda(-2x_2) + (1-\lambda)(2x_2) = -2x_2^2$
- 0 is unstable equilibrium.



#### Vector fields: Filippov's case





#### Utkin's solution concept

$$\begin{aligned} \dot{x}_1 &= -x_1 + x_2 - u \\ \dot{x}_2 &= 2x_2(u^2 - u - 1) \\ u &= \begin{cases} 1, & \text{if } x_1 > 0 \\ -1, & \text{if } x_1 < 0. \end{cases} \end{aligned}$$

• The equivalent control  $u_{\text{equiv}}$  is such that state slides along  $x_1 = 0$ . Hence,  $x_1 = \dot{x}_1 = 0$  and thus  $u_{\text{equiv}} = x_2$  and

$$\dot{x}_2 = 2x_2(x_2^2 - x_2 - 1)$$

• Equilibria: -0.618 (unstable) and 0 (stable)



#### Vector fields



## Embedded Systems

#### Solution trajectories





#### Two relaxations

- Smoothing  $u(t) = \tanh(x_1/\varepsilon)$
- $\bullet$  hysteresis type of switching <code>parameter</code>  $\Delta$



#### Solution trajectories: Filippov's case + hysteresis





#### Solution trajectories: Utkin's case + smoothing





#### Conclusions on discontinuous dynamical systems

- Two mathematical solutions concepts: Filippov + Utkin
- Both limit cases ("idealizations") of very fast switching
- Which one you use depends on non-ideal cases (regularizations)
- Sliding mode might be seen as third mode in hybrid automaton. Some subtleties in HA solution concept!



#### From classical to modern solution concepts



#### Hybrid Systems

- Smooth phases (governed by differential equations)
- Discrete events and actions

Smooth phases separated by event times ...



#### **Event times**



$\dot{x}_1(t)$	=	$x_3(t)$
$\dot{x}_2(t)$	=	$x_4(t)$
$\dot{x}_3(t)$	=	$-2x_1(t) + x_2(t) + z(t)$
$\dot{x}_4(t)$	=	$x_1(t) - x_2(t)$
w(t)	=	$x_1(t)$
w(t)	$\geq$	$0, \ z(t) \ge 0, \ \{w(t) = 0 \text{ or } z(t) = 0\}$

unconstrained	<u>constrained</u>
$\dot{x}_1(t) = x_3(t)$	$\dot{x}_1(t) = x_3(t)$
$\dot{x}_2(t) = x_4(t)$	$\dot{x}_2(t) = x_4(t)$
$\dot{x}_3(t) = -2x_1(t) + x_2(t)$	$\dot{x}_3(t) = -2x_1(t) + x_2(t) + z(t)$
$\dot{x}_4(t) = x_1(t) + x_2(t)$	$\dot{x}_4(t) = x_1(t) + x_2(t)$
z(t) = 0	$w(t) = x_1(t) = 0.$



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• Event times set  $\mathcal{E}$  is  $\{0, 1, 1+\frac{\pi}{2}\}$ 

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#### Example: Bouncing ball



- Reset  $x_2(\tau+) := -cx_2(\tau-)$  when  $x_1(\tau-) = 0$  and  $x_2(\tau-) \le 0$
- The event times:  $\tau_{i+1} = \tau_i + \frac{2c^i x_2(0)}{g}$  when  $x_1(0) = 0$  and  $x_2(0) > 0$ .

• 
$$\lim_{i\to\infty} \tau_i = \tau^* = \frac{2x_2(0)}{g-gc} < \infty$$



#### Zeno of Elea and one of his paradoxes



Distance Travelled (m) by Achilles

I 0.5 0.25 0.125 0.0625 0.03125 0.015625 0.0078125 0.00390625 0.001953125



Event times of A reaching previous T position

I I.5 I.75 I.875 I.9375 I.96875 I.984375 I.9921875 I.99209375 I.99609375 I.998046875



**Definition 4** A set  $\mathcal{E} \subset \mathbb{R}_+$  is called an *admissible event times set*, if it is closed and countable, and  $0 \in \mathcal{E}$ . E.g.  $\mathcal{E} = \{\tau_0, \tau_1, \tau_2, \ldots\}$ .

- An element t of a set E is said to be a *left accumulation point* of E, if for all t' > t (t, t') ∩ E is not empty.
- It is called a *right accumulation point*, if for all t' < t  $(t', t) \cap \mathcal{E}$  is not empty

**Definition 5** An admissible event times set  $\mathcal{E}$  (or the corresponding solution) is said to be *left (right) Zeno free*, if it does not contain any left (right) accumulation points.

 $\bullet$  Bouncing ball  $\rightarrow$  right accumulation point ...



• Time-reversed bouncing ball:



#### Two-tank system and Zeno behavior



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#### A simulation

$$h_1 = h_2 = 1$$
,  $q_1 = 2$ ,  $q_2 = 3$ ,  $q_{\rm in} = 4$ ,  $x_1(0) = x_2(0) = 2$ ,  $q(0) = v_1$ 





#### Two-tank system and Zeno behavior

- Assume total outflow  $q_1 + q_2 > q_{in}$
- Control objective cannot be met and tanks will be empty in finite time
- Infinitely many switchings in finite time (right accumulation point)  $\rightarrow$  right Zeno behavior

Using a non-Zeno solution concept: analysis will show that tanks do not get empty! Analysis depends crucially on solution concept!



#### Hybrid automaton

Hybrid automaton H is collection  $H=(Q,X,f,\mathrm{Init},\mathrm{Inv},E,G,R)$  with

- $Q = \{q_1, \ldots, q_N\}$  is finite set of discrete states or *modes*
- $X = \mathbb{R}^n$  is set of continuous states
- $f: Q \times X \to X$  is vector field
- Init  $\subseteq Q \times X$  is set of initial states
- Inv :  $Q \rightarrow P(X)$  describes the *invariants*
- $E \subseteq Q \times Q$  is set of edges or *transitions*
- $G: E \rightarrow P(X)$  is guard condition
- $R: E \to P(X \times X)$  is reset map



#### What is what?

Hybrid automaton H = (Q, X, f, Init, Inv, E, G, R)

- Hybrid state: (q, x)
- Evolution of continuous state in mode q:  $\dot{x} = f(q, x)$
- Invariant Inv: describes conditions that continuous state has to satisfy at given mode
- $\bullet$  Guard G: specifies subset of state space where certain transition is enabled
- Reset map R: specifies how new continuous states are related to previous continuous states

## Embedded Systems





#### Evolution of hybrid automaton

- Initial hybrid state  $(q_0, x_0) \in \text{Init}$
- $\bullet$  Continuous state x evolves according to

$$\dot{x} = f(q_0, x) \quad \text{with } x(0) = x_0$$

discrete state q remains constant:  $q(t) = q_0$ 

- Continuous evolution can go on as long as  $x \in Inv(q_0)$
- If at some point state x reaches guard  $G(q_0, q_1)$ , then
  - transition  $q_0 \rightarrow q_1$  is enabled
  - discrete state *may* change to  $q_1$ , continuous state then jumps from current value  $x^-$  to new value  $x^+$  with  $(x^-, x^+) \in R(q_0, q_1)$
- Next, continuous evolution resumes and whole process is repeated



#### Hybrid time trajectory

**Definition 6** A hybrid time trajectory  $\tau = \{I_i\}_{i=0}^N$  is a finite  $(N < \infty)$  or infinite  $(N = \infty)$  sequence of intervals of the real line, such that

- $I_i = [\tau_i, \tau'_i]$  with  $\tau_i \le \tau'_i = \tau_{i+1}$  for  $0 \le i < N$ ;
- if  $N < \infty$ , either  $I_N = [\tau_N, \tau'_N]$  or  $I_N = [\tau_N, \tau'_N)$  with  $\tau_N \le \tau'_N \le \infty$ .
- For instance,

$$\tau = \{[0, 2], [2, 3], \{3\}, \{3\}, [3, 4.5], \{4.5\}, [4.5, 6]\}$$
  
$$\tau = \{[0, 2], [2, 3], [3, 4.5], \{4.5\}, [4.5, 6], [6, \infty)\}$$
  
$$I_i = [1 - 2^i, 1 - 2^{i+1}]$$

- $\mathcal{E} = \{\tau_0, \tau_1, \tau_2, \ldots\}$
- No left-accumulations of event times ...



#### Execution of hybrid automaton

**Definition 7** An execution  $\chi$  of a HA consists of  $\chi = (\tau, q, x)$ 

- $\tau$  a hybrid time trajectory;
- $q = \{q_i\}_{i=0}^N$  with  $q_i : I_i \to Q$ ; and
- $x = \{x_i\}_{i=0}^N$  with  $x_i : I_i \to X$

Initial condition  $(q(\tau_0), x(\tau_0)) \in \text{Init};$ 

**Continuous evolution** for all i

- $q_i$  is constant, i.e.  $q_i(t) = q_i(\tau_i)$  for all  $t \in I_i$ ;
- $x_i$  is solution to  $\dot{x}(t) = f(q_i(t), x(t))$  on  $I_i$  with initial condition  $x_i(\tau_i)$  at  $\tau_i$ ;
- for all  $t \in [\tau_i, \tau'_i)$  it holds that  $x_i(t) \in \text{Inv}(q_i(t))$ .

Discrete evolution for all *i*,

- $e = (q_i(\tau'_i), q_{i+1}(\tau_{i+1})) \in E$ ,
- $x(\tau'_i) \in G(e)$ ;
- $(x_i(\tau'_i), x_{i+1}(\tau_{i+1})) \in R(e).$



#### Well-posedness for hybrid automata

- $\mathcal{H}^{\infty}_{(q_0,x_0)}$ : infinite executions:  $\tau$  is an infinite sequence or if  $\sum_i (\tau'_i \tau_i) = \infty$
- $\mathcal{H}^{M}_{(q_0,x_0)}$ : maximal executions:  $\tau$  is not a strict prefix of another one!
- A hybrid automaton is called *non-blocking*, if  $\mathcal{H}^{\infty}_{(q_0,x_0)}$  is non-empty for all  $(q_0, x_0) \in \text{Init.}$
- It is called *deterministic*, if  $\mathcal{H}^M_{(q_0,x_0)}$  contains at most one element for all  $(q_0, x_0) \in \text{Init.}$



#### Well-posedness for hybrid automata - continued

#### Assumption

- The vector field  $f(q, \cdot)$  is globally Lipschitz continuous for all  $q \in Q$ .
- The edge e = (q, q') is contained in E if and only if  $G(e) \neq \emptyset$  and  $x \in G(e)$  if and only if there is an  $x' \in X$  such that  $(x, x') \in R(e)$ .

A state  $(\hat{q}, \hat{x}) \in \text{Reach}$ , if there exists a finite execution  $(\tau, q, x)$  with  $\tau = \{[\tau_i, \tau'_i]\}_{i=0}^N$  and  $(q(\tau'_N), x(\tau'_N)) = (\hat{q}, \hat{x})$ .

The set of states from which continuous evolution is impossible :

$$\operatorname{Out} = \{ (q_0, x_0) \in Q \times X \mid \forall \varepsilon > 0 \exists t \in [0, \varepsilon) \; x_{q_0, x_0}(t) \notin \operatorname{Inv}(q_0) \}$$

in which  $x_{q_0,x_0}(\cdot)$  denotes the unique solution to  $\dot{x}=f(q_0,x)$  with  $x(0)=x_0.$ 



#### Well-posedness theorems

**Theorem** A hybrid automaton is non-blocking, if for all  $(q, x) \in \text{Reach} \cap$ Out, there exists  $e = (q, q') \in E$  with  $x \in G(e)$ . In case the automaton is deterministic, this condition is also necessary.

**Theorem** A hybrid automaton is deterministic, if and only if for all  $(q, x) \in \operatorname{Reach}$ 

- if  $x \in G((q,q'))$  for some  $(q,q') \in E$ , then  $(q,x) \in \text{Out}$ ;
- $\bullet$  if  $(q,q')\in E$  and  $(q,q'')\in E$  with  $q'\neq q''$  , then  $x\not\in G((q,q'))\cap G((q,q''));$  and
- if  $(q,q')\in E$  and  $x\in G((q,q')),$  then there is at most one  $x'\in X$  with  $(x,x')\in R((q,q')).$

 $\longrightarrow$  no explicit / algebraic conditions and not easily verifiable  $\rightarrow$  can we do more (like for DDE)?



#### Well-posedness issues

- Initial well-posedness: non-blocking + deterministic, i.e. absence of
  - **dead-lock**: no smooth continuation and no jump
  - splitting of trajectories

However, no statements by HA theory on existence beyond

- live-lock: an infinite number of jumps at one time instant, no solution on [0, ε) for some ε > 0.
- right-accumulations of event times to prevent global existence.

or absence of

• left-accumulations of event times preventing uniqueness:





#### Obstruction local existence

#### $\rightarrow$ Live-lock: Infinitely many jumps at one time instant



- ber of events
- $\longrightarrow$  Exclude live-lock or show convergence of state x for local existence
- Discrete mode is a function of continuous state! not for general HA!!!



#### **Obstruction global existence: Zenoness**

 $\rightarrow$  A right-accumulation of event times



- Exclude right-accumulations or show the existence of the left-limit  $\lim_{t\uparrow \tau^*} x(t)$  for global existence.
- Discrete mode is a function of continuous state! not for general HA!!!



#### Obstructions local uniqueness: Filippov's example



$$\dot{x}_1 = \operatorname{sgn}(x_1) - 2\operatorname{sgn}(x_2)$$
  
 $\dot{x}_2 = 2\operatorname{sgn}(x_1) + \operatorname{sgn}(x_2),$ 

Left accumulation point ...  $\mathcal{E}$  is not left Zeno free!

Well-posedness:

- Due to left-accumulations non-uniqueness in origin
- Using HA framework: non-blocking and deterministic
- Using Filippov's solution: non-uniqueness!



#### Well-posedness

- Initially solvable from each initial state there exists a state jump or a continuous hybrid solution on  $[0, \varepsilon)$  (non-blocking)
- Initially unique from each initial state the jump/hybrid solution is unique (deterministic)
- Local well-posedness from each initial state there exists an  $\varepsilon > 0$  and a hybrid solution on  $[0, \varepsilon)$ .
- Global well-posedness ... on  $[0,\infty)$ .

Embedded Systems

#### Piecewise linear systems

$$\begin{aligned} \operatorname{SAT}(A,B,C,D) & \dot{x}(t) = Ax(t) + Bu(t) & e_2^i - e_1^i > 0 \text{ and } f_1^i \geq f_2^i \\ y(t) = Cx(t) + Du(t) \\ (u(t),y(t)) \in \operatorname{saturation}_i \end{aligned}$$



Note that if  $f_2^i = f_1^i$ , then relay-type of nonlinearity.



#### Example of linear relay system: non-uniqueness

 $\begin{array}{lll} \dot{x} &=& x-u \\ y &=& x \\ u &\in& -\mathrm{sgn}(y) \end{array}$ 



• 
$$x(t) = e^t - 1$$
,  $(y(t) = x(t) \ge 0)$ 

• 
$$x(t) = -e^t + 1$$
,  $(y(t) = x(t) \le 0)$ 

• 
$$x(t) = 0$$
,  $(y(t) = x(t) = 0)$ 





#### Example of linear relay system: uniqueness

$$\begin{aligned} \dot{x} &= x + u \\ y &= x \\ u &\in -\mathrm{sgn}(y) \end{aligned}$$

x(0) = 0:

• 
$$x(t) = 0$$
,  $(y(t) = x(t) = 0)$ 





Piecewise linear systems



Consider SAT(A, B, C, D).

- Let R and S be the diagonal matrices with  $e_2^i e_1^i$  and  $f_2^i f_1^i$ , resp.
- $G(s) = C(sI A)^{-1}B + D$

Suppose that  $G(\sigma)R - S$  is a P-matrix for all sufficiently large  $\sigma$ . Then, there exists a unique (left Zeno free) hybrid execution of SAT(A, B, C, D) for all initial states.

•  $M \in \mathbb{R}^{m \times m}$  is a *P*-matrix, if det $M_{II} > 0$  for all  $I \subseteq \{1, \ldots, m\}$ .



Linear relay systems and Filippov's solution concept: left accumulations

 $\dot{x}(t) = Ax(t) + Bu(t); \quad y(t) = Cx(t); \quad u(t) \in -\operatorname{sgn}(y(t))$ 

Previous result: If  $G(\sigma) = CB\sigma^{-1} + CAB\sigma^{-2} + ... > 0$  for sufficiently large  $\sigma$ , then existence and uniqueness of (left-Zeno free) executions.



#### Other solution concept ...?

Filippov's solutions include left-accumulations and satisfy  $\dot{x} \in F(x)$  almost everywhere, with

- $F(x) = \{Ax + B\}$  for Cx < 0
- $F(x) = \{Ax B\}$  for Cx > 0
- $F(x) = \{Ax + B\overline{u} \mid \overline{u} \in [-1, 1]\}$  when Cx = 0

In case of relative degree 1 (CB > 0) and relative degree 2 (and order 2) sufficient for Filippov uniqueness.



So, (other) example of HA uniqueness (deterministic), but non-uniqueness in "Filippov"



#### Linear complementarity systems





 $\{z_i(t) = 0 \text{ and } w_i(t) \ge 0\} \text{ or } \{w_i(t) = 0 \text{ and } z_i(t) \ge 0\}$ 

• modes parameterized by  $I \subseteq \{1, \ldots, k\}$  such that

$$\begin{aligned} \dot{x}(t) &= Ax(t) + Bz(t) \\ w(t) &= Cx(t) + Dz(t) \\ w_i &= 0, \ i \in I \ \text{ and } z_i = 0, i \notin I \end{aligned}$$
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#### Example 1





• 
$$z = 0$$
:  $\dot{x} = x$ ,  $w = x \ge 0$ 

• 
$$w = 0$$
:  $\dot{x} = 2x$ ,  $z = x \ge 0$ 

Hence, x(0) = 1 two solutions and x(0) = -1 no solution trajectory!



#### Example 2

$$\dot{x} = x + z$$
  

$$w = x + z$$
  

$$0 \le w \perp z \ge 0$$
  
•  $z = 0: \dot{x} = x, w = x \ge 0$ 

• 
$$w = 0$$
:  $\dot{x} = 0$ ,  $z = -x > 0$ 

Existence and uniqueness!

Model test ...





#### Well-posedness including jumps

- Initially solvable from each initial state there exists a state jump or a continuous hybrid solution on  $[0, \varepsilon)$  (non-blocking)
- Initially unique from each initial state the jump/hybrid solution is unique (deterministic)
- Local well-posedness from each initial state there exists an  $\varepsilon > 0$  and a hybrid solution on  $[0, \varepsilon)$ .
- Global well-posedness ... on  $[0,\infty)$ .

#### Local well-posedness (including jumps)

$$\label{eq:constraint} \begin{split} \dot{x}(t) &= Ax(t) + Bz(t), \quad w(t) = Cx(t) + Dz(t), \ 0 \leq z(t) \perp w(t) \geq 0 \\ \text{Markov parameters:} \ H^0 &= D \text{ and } H^i = CA^{i-1}B \text{, } i = 1,2,\ldots \end{split}$$

$$\eta_j = \inf\{i \mid H^i_{\bullet j} \neq 0\}, \ \rho_j = \inf\{i \mid H^i_{j\bullet} \neq 0\},\$$

The leading row and column coefficient matrices  $\mathcal M$  and  $\mathcal N$ 

$$\mathcal{M} := \begin{pmatrix} H_{1\bullet}^{\rho_1} \\ \vdots \\ H_{k\bullet}^{\rho_k} \end{pmatrix} \text{ and } \mathcal{N} := (H_{\bullet 1}^{\eta_1} \dots H_{\bullet k}^{\eta_k})$$

•  $M \in \mathbb{R}^{m \times m}$  is a *P*-matrix, if det $M_{II} > 0$  for all  $I \subseteq \{1, \ldots, m\}$ .

If  $\mathcal{N}$  and  $\mathcal{M}$  are defined and P-matrices, then LCS(A, B, C, D) has for all  $x_0$  a unique left Zeno free execution on an interval of the form  $[0, \varepsilon)$  for some  $\varepsilon > 0$ .

- Moreover, live-lock does not occur: at most one jump
- Necessary and sufficient for **global** well-posedness for **bimodal** LCS

Embedded Systems



#### Summary

- Smooth differential equations
  - Solution concept straightforward
  - Lipschitz continuity sufficient for well-posedness
  - absence Lipschitz: possibly non-uniqueness
  - absence global Lipschitz finite escape times and no global existence
- Switched systems (discontinuous differential equations)
  - Sliding modes (Filippov's convex or Utkin's equivalent control definition)
  - Solution concept from differential inclusions
  - Well-posedness: directions of vector field at switching plane

"No events"



#### Summary - continued

- Hybrid systems:
  - Complications due to Zeno
  - Relation between solution concept and well-posedness and analysis
    - \* Tanks stay full along non-Zeno solutions!!!
    - \* Filippov's example has unique non-Zeno solutions, but nonunique Zeno solutions
  - Well-posedness
    - \* Initial well-posedness (non-blocking and deterministic)
    - \* Local well-posedness:  $[0,\varepsilon)$  (live-lock)
    - $\ast$  Global well-posedness:  $[0,\infty)$  (right-accumulations)
  - Conditions for hybrid automata: implicit!
  - Algebraic conditions for certain classes with more structure!



#### Selected Literature

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