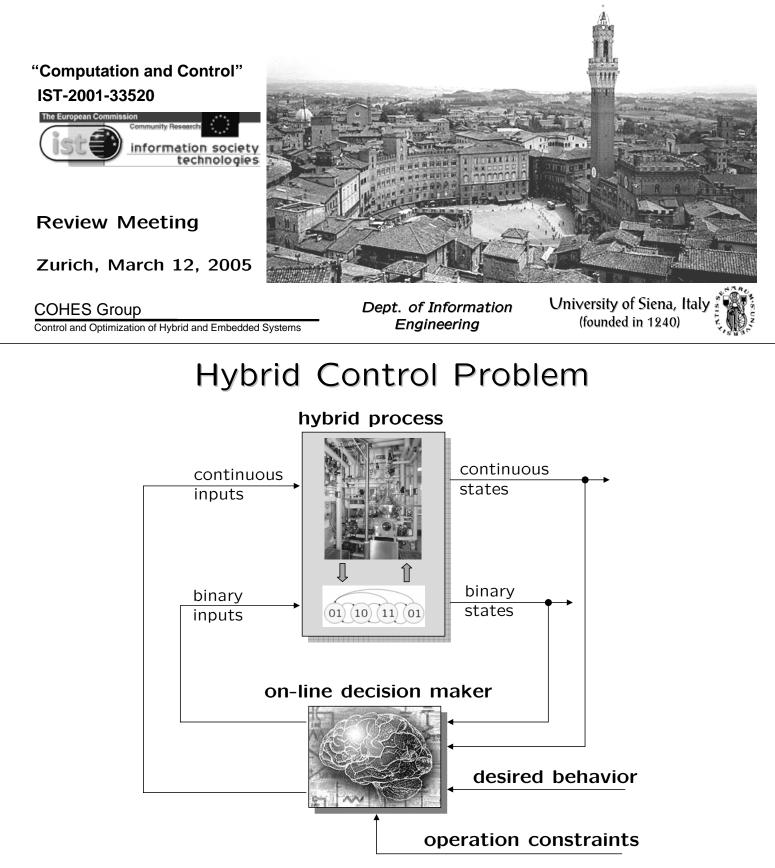
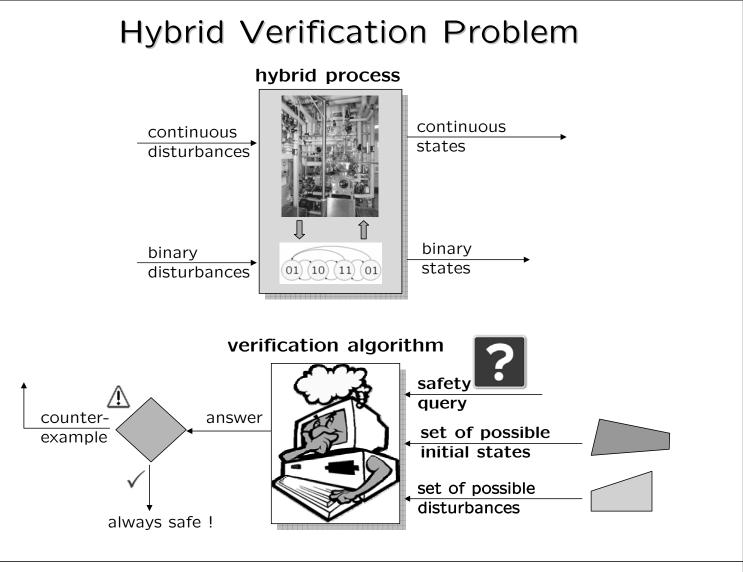
Optimization-Based Methods for Controller Synthesis and Verification

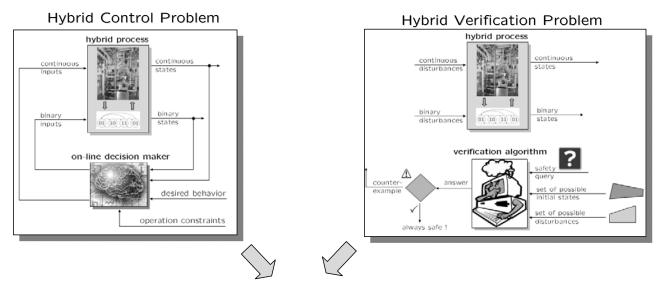
Alberto Bemporad

http://www.dii.unisi.it/hybrid



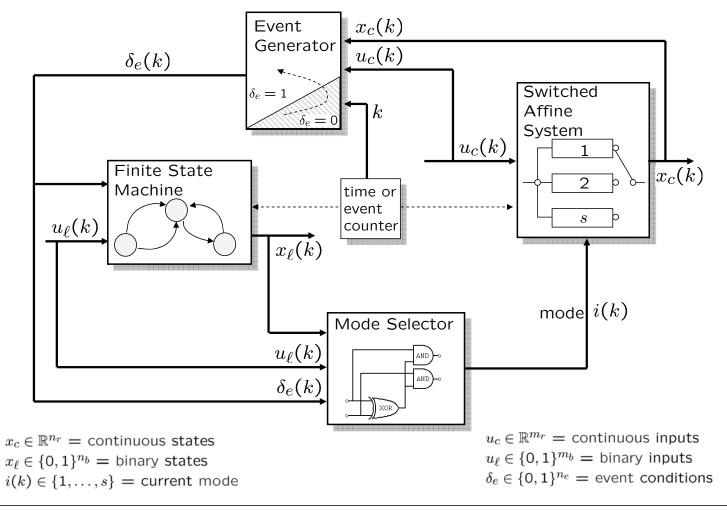


Model-based Optimization Approach

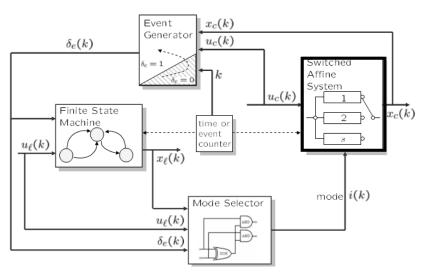


- Need for a <u>hybrid model</u> of the process reproducing the behavior of the process (simulation)
- A model suitable for controller synthesis and verification
- A model for which computational tools can be applied

Hybrid Model: Discrete Hybrid Automaton



Switched Affine System

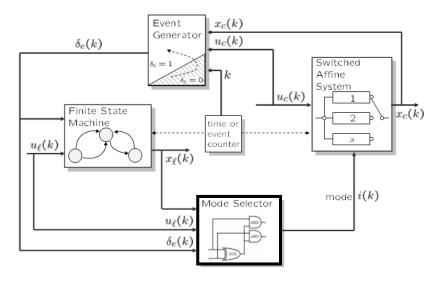


The affine dynamics depend on the current mode i(k):

 $x_c(k+1) = A_{i(k)}x_c(k) + B_{i(k)}u_c(k) + f_{i(k)}$

 $x_c \in \mathbb{R}^{n_c}, \ u_c \in \mathbb{R}^{m_c}$

Mode Selector



The active mode i(k) is selected by a Boolean function of the current binary states, binary inputs, and event variables:

$$i(k) = f_{\mathsf{M}}(x_{\ell}(k), u_{\ell}(k), \delta_{e}(k))$$
 $x_{\ell} \in \{0, 1\}^{n_{\ell}}, u_{\ell} \in \{0, 1\}^{m_{\ell}}, \delta_{e}\{0, 1\}^{n_{e}}$

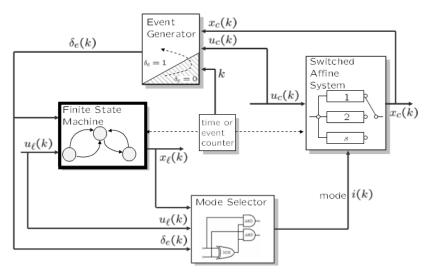
Example:

$$i(k) = \begin{bmatrix} \neg u_{\ell}(k) \lor x_{\ell}(k) \\ u_{\ell}(k) \land x_{\ell}(k) \end{bmatrix}$$

$$\longrightarrow \frac{\frac{u_{\ell}/x_{\ell}}{0} \quad 0 \quad 1}{\frac{1}{1} \quad i = \begin{bmatrix} 1 \\ 0 \end{bmatrix}} \quad i = \begin{bmatrix} 1 \\ 1 \end{bmatrix}^{-1}$$

the system has 3 modes

Finite State Machine

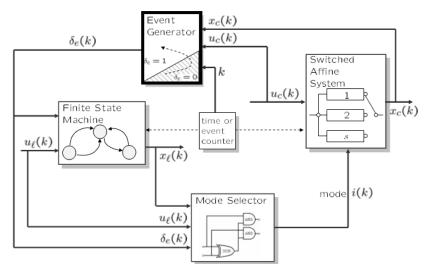


The binary state of the finite state machine evolves according to a Boolean state update function:

 $x_{\ell}(k+1) = f_{\mathsf{B}}(x_{\ell}(k), u_{\ell}(k), \delta_{e}(k)) \qquad x_{\ell} \in \{0, 1\}^{n_{\ell}}, \ u_{\ell} \in \{0, 1\}^{m_{\ell}}, \ \delta_{e} \in \{0, 1\}^{n_{e}}$

Example: $x_{\ell}(k+1) = \neg \delta_e(k) \lor (x_{\ell}(k) \land u_{\ell}(k))$

Event Generator



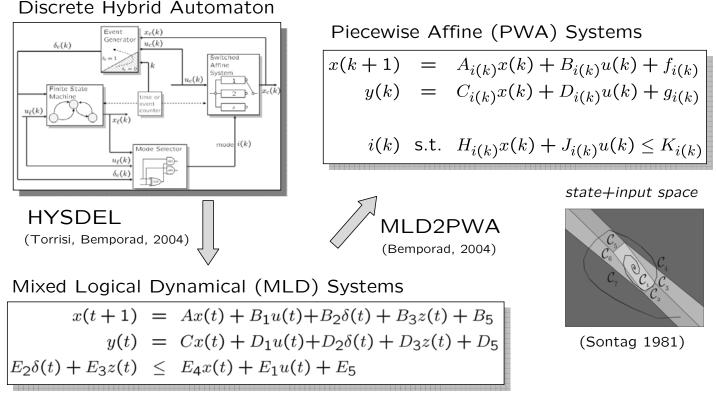
Event variables are generated by linear threshold conditions over continuous states, continuous inputs, and time:

$$[\delta_e^i(k) = 1] \leftrightarrow [H^i x_c(k) + K^i u_c(k) \le W^i]$$

 $x_c \in \mathbb{R}^{n_c}, \ u_c \in \mathbb{R}^{m_c}, \ \delta_e \in \{0, 1\}^{n_e}$

Example: $[\delta=1] \leftrightarrow [x_c(k) \ge 0]$

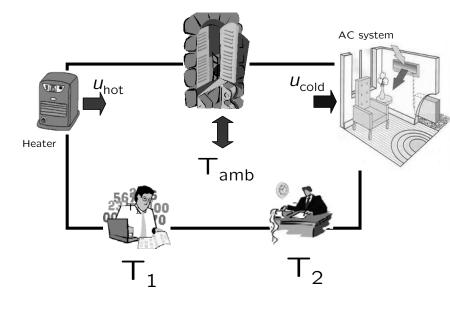
Computational Hybrid Models



(Bemporad, Morari 1999)

The translation from DHA to MLD/PWA is done automatically (using symbolic/mathematical programming tools)

Example: Room Temperature



Hybrid Dynamics

- #1 turns the heater (air conditioning) on whenever he is cold (hot)
- If #2 is cold he turns the heater on, unless #1 is hot
- If #2 is hot he turns the air conditioning on, unless #2 is cold
- Otherwise, heater and air conditioning are off
- $\dot{T}_1 = -\alpha_1(T_1 T_{amb}) + k_1(u_{hot} u_{cold})$
- $\dot{T}_2 = -\alpha_2(T_2 T_{amb}) + k_2(u_{hot} u_{cold})$
- (body temperature dynamics of #1)
 - (body temperature dynamics of #2)

HYSDEL Model

```
SYSTEM heatcool (
INTERFACE {
    STATE { REAL T1 [-10,50];
             REAL T2 [-10,50];
    INPUT ( REAL Tamb [-10,50];
        - }
    PARAMETER {
        REAL Ts, alpha1, alpha2, k1, k2;
        REAL Thot1, Tcold1, Thot2, Tcold2, Uc, Uh;
IMPLEMENTATION {
        AUX { REAL uhot, ucold;
               BOOL hot1, hot2, cold1, cold2;
        AD { hot1 = T1>=Thot1;
hot2 = T2>=Thot2;
               cold1 = T1<=Tcold1;
cold2 = T2<=Tcold2;</pre>
        - }
        DA { uhot = {IF cold1 | (cold2 & ~hot1) THEN Uh ELSE 0};
               ucold = {IF hot1 | (hot2 & ~cold1) THEN UC ELSE 0};
        }
        CONTINUOUS { T1 = T1+Ts*(-alpha1*(T1-Tamb)+k1*(uhot-ucold));
                       T2 = T2 + Ts * (-alpha2 * (T2 - Tamb) + k2 * (uhot - ucold));
        }
```







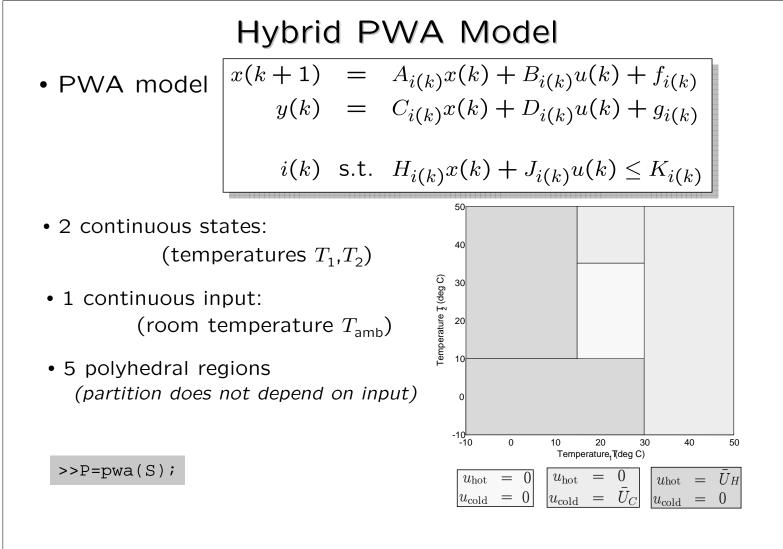
http://www.dii.unisi.it/hybrid/toolbox

>>S=mld('heatcoolmodel',Ts)

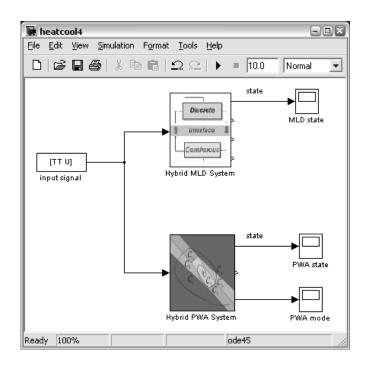
get the MLD model in Matlab

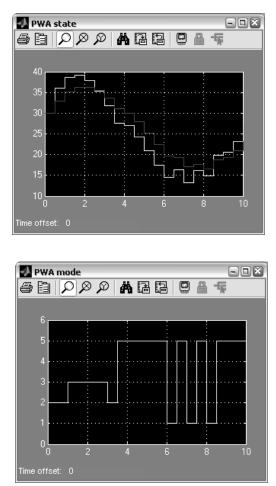
>>[XX,TT]=sim(S,x0,U);

simulate the MLD model



Simulation in Simulink





Verification of DHA/MLD/PWA

Verification Algorithm

- **QUERY**: Is the target set X_f reachable after N steps from some initial state $x_0 \in X_0$ for some input profile $u \in U$?
- Computation: Solve the mixed-integer linear program (MILP)

with respect to $u(0),\delta(0),z(0),..., u(N-1),\delta(N-1),z(N-1),x(0)$

- Alternative solutions:
 - Exploit the special structure of the problem and use polyhedral (Torrisi, 2003)
 - Use abstractions (LPs) + SAT solvers (Giorgetti, Pappas, 2005)

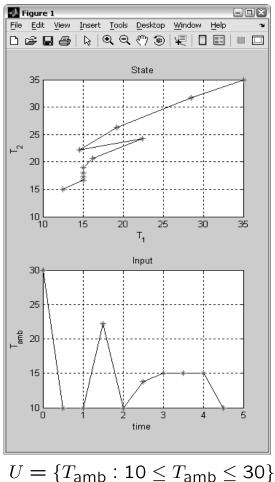
Verification Example

- MLD model: room temperature system
- $X_f = \{ \begin{bmatrix} T_1 \\ T_2 \end{bmatrix} : 10 \le T_1, T_2 \le 15 \}$ (set of unsafe states)
- $X_0 = \{ \begin{bmatrix} T_1 \\ T_2 \end{bmatrix} : 35 \le T_1, T_2 \le 40 \}$ (set of initial states)
- $U = \{T_{amb} : 10 \le T_{amb} \le 30\}$ (set of possible inputs)
- N=10 (time horizon)



>>[flag,x0,U]=reach(S,N,Xf,X0,umin,umax);

Verification Example

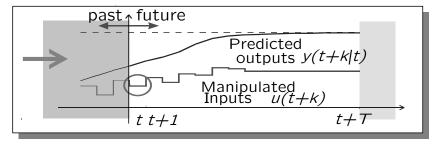


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 $U = \{T_{amb}: 20 \le T_{amb} \le 30\}$

Controller Synthesis





Model Predictive (MPC) Control

• At time *t* solve with respect to $U \triangleq \{u(t), u(t+1), \dots, u(t+T-1)\}$ the finite-horizon open-loop, optimal control problem:

$$\min_{\substack{u(t),...,u(t+T-1) \\ \text{subject to}}} \sum_{k=0}^{T-1} \|R(y(t+k|t) - r(t+k))\|_p + \|Qu(t+k)\|_p$$

$$\text{subject to} \begin{cases} \mathsf{MLD or PWA model} \\ x(t|t) = x(t) \end{cases}$$

$$p = 1, 2, \infty \quad \|v\|_2 = v'v \quad \|v\|_\infty = \max|v_i| \quad \|v\|_1 = \sum v_i$$

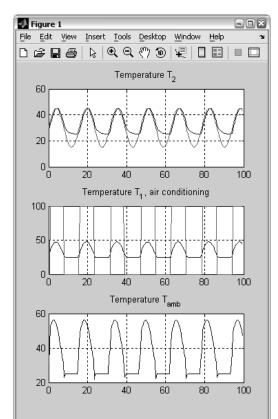
- Apply only $u(t) = u^*(t)$ (discard the remaining optimal inputs);
- Repeat the whole optimization at time t+1

Hybrid MPC - Example

>>refs.x=2; % just weight state #2
>>Q.x=1;
>>Q.rho=Inf; % hard constraints
>>Q.norm=2; % quadratic costs
>>N=2; % optimization horizon
>>limits.xmin=[25;-Inf];

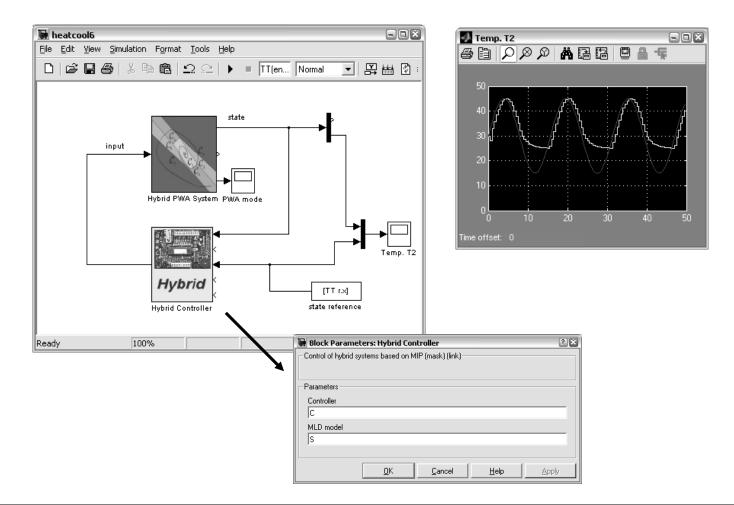
>>C=hybcon(S,Q,N,limits,refs);

>> C
Hybrid controller based on MLD model S <heatcoolmodel.hys>
2 state measurement(s)
0 output reference(s)
0 input reference(s)
1 state reference(s)
0 reference(s) on auxiliary continuous z-variables
20 optimization variable(s) (8 continuous, 12 binary)
46 mixed-integer linear inequalities
sampling time = 0.5, MILP solver = 'glpk'
Type "struct(C)" for more details.
>>



>>[XX,UU,DD,ZZ,TT]=sim(C,S,r,x0,Tstop);

Hybrid MPC - Example



On-Line vs Off-Line Optimization

 $\min_{U} J(U(x(t))) = \sum_{k=0}^{T-1} ||Rx(t+k|t)||_p + ||Qu(t+k)||_p$ subject to $\begin{cases} \mathsf{MLD} \mod k \\ x(t|t) = x(t) \end{cases}$

• On-line optimization: given x(t) solve the problem at each time step t.

Mixed-Integer Linear/Quadratic Program (MILP/MIQP)

- Good for large sampling times (e.g., 1 h) / expensive hardware but not for fast sampling (e.g. 10 ms) / cheap hardware !
- Off-line optimization: solve the MILP/MIQP for all x(t)

$$\min_{\zeta} J(\zeta, x(t)) = \begin{cases} f'\zeta & \infty/1\text{-norms} \\ \zeta'H\zeta + f'\zeta & \text{quadratic forms} \end{cases}$$

s.t. $G\zeta \leq W + Fx(t)$

multi-parametric programming

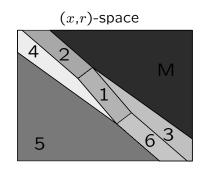
Explicit Hybrid MPC

$$\min_{U} J(U, x, r) = \sum_{k=0}^{T-1} ||R(y(k) - r)||_p + ||Qu(k)||_p$$

subject to
$$\begin{cases} \mathsf{PWA model} \\ x(0) = x \end{cases}$$

- Solution u(x,r) found via a combination of
 - Dynamic programming or enumeration of feasible mode sequences, multiparametric linear or quadratic programming, and polyhedral computation. (Borrelli, Baotic, Bemporad, Morari, 2003)
 - (Mayne, ECC 2001) (Alessio, Bemporad, 2005)
- The MPC controller is piecewise affine in x,r

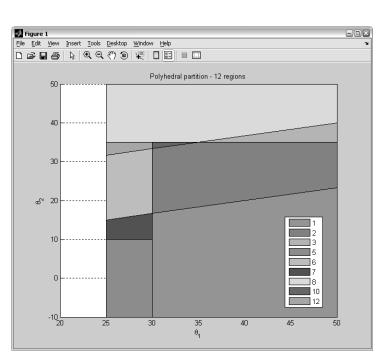
$$u(x,r) = \begin{cases} F_1 x + E_1 r + g_1 & \text{if } H_1[x] \le K_1 \\ \vdots & \vdots \\ F_M x + E_M r + g_M & \text{if } H_M[x] \le K_M \end{cases}$$



Note: in the quadratic case the partition may not be fully polyhedral

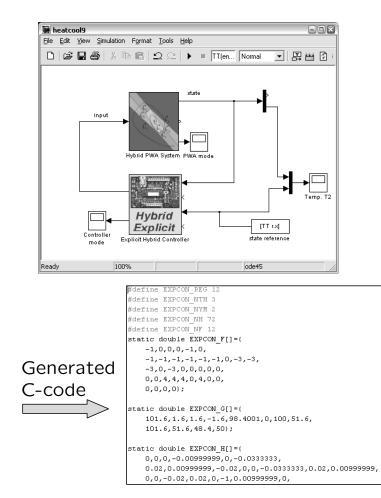
Explicit Hybrid MPC - Example

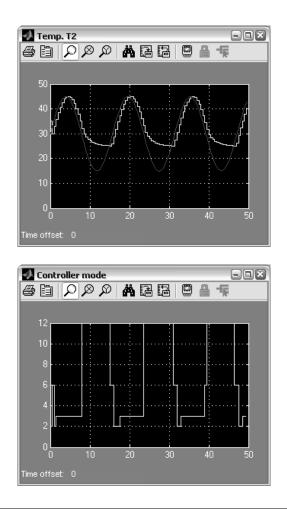
>>E=expcon(C,range,options);					
>> E					
<pre>Explicit controller (based on hybrid controller C) 3 parameter(s) 1 input(s) 11 partition(s) sampling time = 0.5</pre>					
The controller is for hybrid systems (tracking) This is a state-feedback controller.					
Type "struct(E)" for more details.					

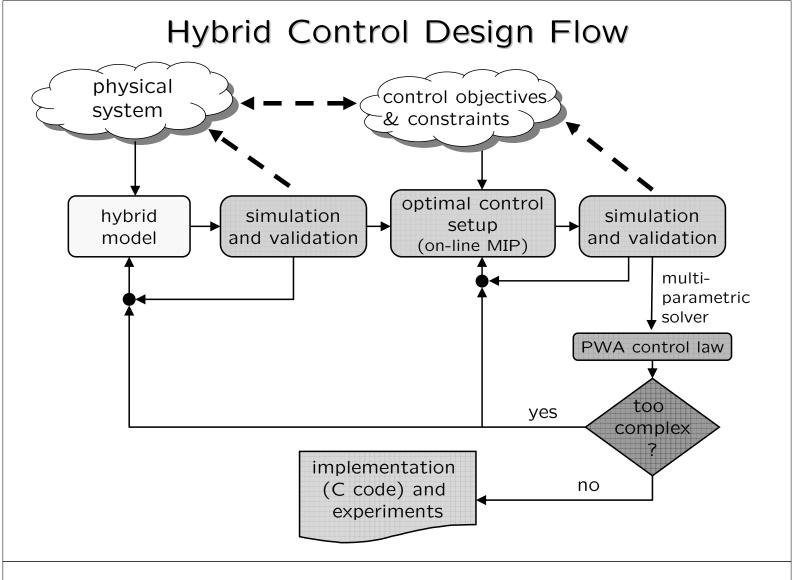


Section in the (T_1, T_2) -space for $T_{ref} = 30$

Explicit Hybrid MPC - Example







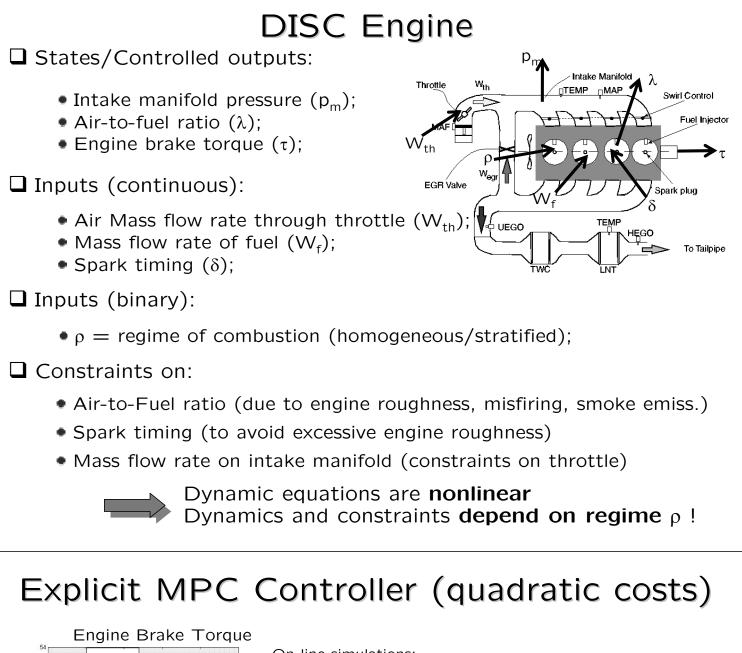


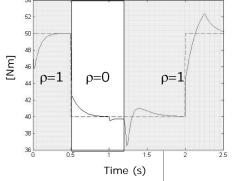
(Photo: Courtesy Mitsubishi)

Application of Explicit Hybrid MPC to DISC Engine Control

(joint work with N. Giorgetti, I. Kolmanovsky, D. Hrovat)







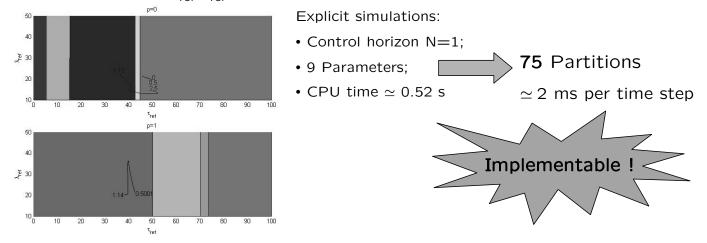
On-line simulations:

- Control horizon N=1; $\simeq 32$
- Simulation time T=2.5 s;
- Sampling time Ts=10 ms;
- CPU time $\simeq 8$ s; (on an Intel Centrino 1.2 GHz, 640 Mb RAM with Cplex 9.01)

 \simeq 32 ms per time step

Not directly implementable !

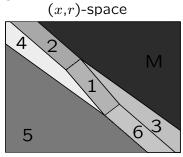
Cross-section in the $\tau_{ref} - \lambda_{ref}$ plane



Conclusions

- Discrete hybrid automata are simple yet versatile models of hybrid systems, and lead immediately computationally-useful models
- Optimization-based control handles performance specs and constraints in a natural and direct way. Quite complex hybrid systems can be controlled using on-line mixed-integer programming
- Piecewise affine MPC controllers can be synthesized, off-line, and implemented as look-up tables of linear gains

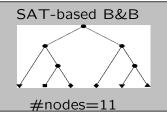
$$u(x,r) = \begin{cases} F_1 x + E_1 r + g_1 & \text{if } H_1[\frac{x}{r}] \le K_1 \\ \vdots & \vdots \\ F_M x + E_M r + g_M & \text{if } H_M[\frac{x}{r}] \le K_M \end{cases}$$

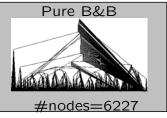


• Hybrid Toolbox for Matlab available to assist controller design: modeling, simulation, verification, MPC, code generation

Ongoing Research

• Logic-based methods: use combined SAT+LP to largely improve computation performance of optimal control/verification





(N. Giorgetti)

- Stochastic hybrid models: (robust) optimal control can be solved using mixed-integer programming (Di Cairano, HSCC05)
- Event-based optimal control of hybrid systems: index k denotes events occurring in piecewise-affine continuous-time systems, rather than discrete time

	N	t_s	t_{cpu} (sec.)	t_{viol} (time units)	$t_{viol}^{(sup)}$ (time units)
optimal	5	20.00	0.45	3.82	40.00
control	6	16.66	1.46	21.83	33.32
train-gate	10	10.00	8.58	9.82	20.00
system	15	6.66	66.77	6.55	13.32
	20	5.00	747.14	4.82	10.00
		cont.time	1.09		

(Di Cairano, Julvez)



